Congestion Management Capabilities of Various Fabrics

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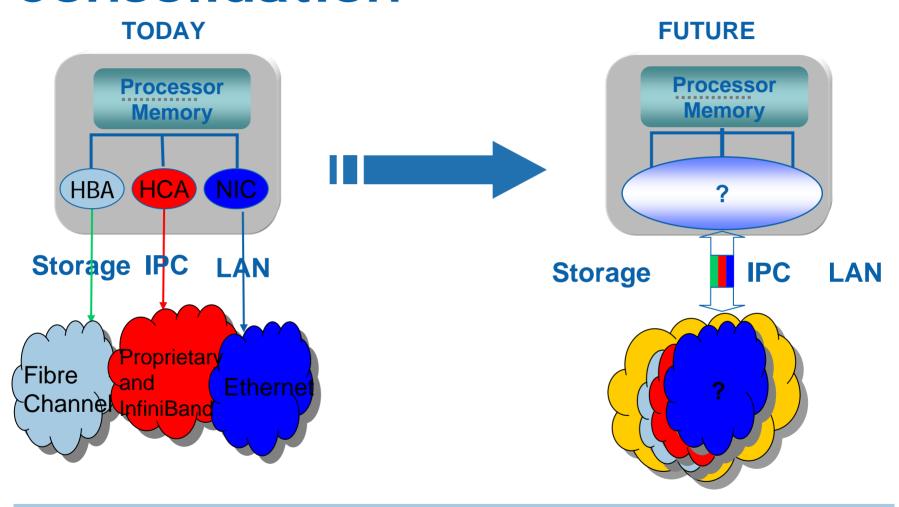
IEEE 802.1 Congestion Management Interim

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Some Fabric Interconnects in Today's Data Centers

- Ethernet
 - Default fabric for networking traffic
- InfiniBand
 - Emerging fabric for inter-process communication (IPC) traffic
- Fiber Channel
 - Preferred fabric for storage traffic
- Advanced Switching Interconnect (ASI)
 - Aimed at inter-board, backplane, and multi-rack systems typically used in Telecom environments
- Proprietary (Myrinet, Quadrics, ...)
 - Preferred fabric for low latency traffic

Consolidation



Trend towards consolidation due to higher cost and complexity of multiple fabrics – Can Ethernet be the consolidated fabric?

- Link level flow control
- End-to-End flow control
- Static Rate Control
- End-to-End injection rate control
- VL arbitration

Link level flow control

- Granular link level credit based flow control
 - Per Virtual Link (VL) flow control (upto 15 data VLs)
 - Per VL flow control eliminates HOL blocking between VLs
 - Latency sensitive traffic using one set of VLs can be unaffected by congestion of best effort traffic in other VLs
- Effectively deals with transient congestion without dropping packets
 - Feedback-based mechanisms cannot deal with time constants of transient congestion

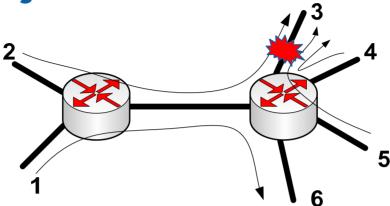
End-to-End flow control

- End-to-End (message level) credit based flow control
- Used by reliable connections to prevent receive queue overflow
- End-to-End credits are generated by a responder's receive queue and consumed by a requester's send queue
- Encoded credits are transported from the responder to the requester in an ACK message

InfiniBand Static Rate Control

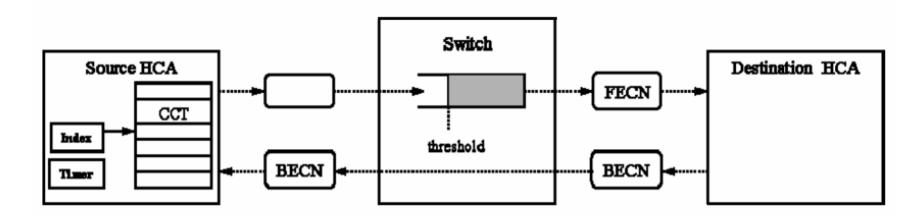
- Enables Fabric Manager to configure the network to avoid oversubscribed links in the fabric
- Effectively handle link speed mismatches
 - A path with 4X link feeding a 1X link can be constrained to a 1X rate

End-to-End injection rate control (1)



- Added in 2004 (Release 1.2) to address congestion spreading problem
 - Oversubscription to destination 3, blocks traffic on the inter-switch link
 - This blocks flow to destination 6 which is uncongested
- Uses Forward Explicit Congestion Notification (FECN) and Backward Explicit Congestion Notification (BECN)
- Pushes congestion to the source to avoid head-of-line blocking and congestion spreading

End-to-End injection rate control (2)



- Switches detect congestion based on threshold and mark packets with FECN bit
- Destination responds to source queue pair with BECN, either piggybacked on ACK or special Congestion Notification packet
- Source reduces the injection rate of flow by increasing index into Congestion Control Table (CCT) which stores inter-packet delay values
- Source recovers rate based on timer by decreasing index into CCT on timer expiry

VL arbitration

- Dual priority WRR arbitration scheme between VLs
- Allows configurable link sharing between VLs



Fibre Channel

Fibre Channel

Flow Control

- Buffer-to-Buffer credit based flow control
 - Destination port signals by sending a Receiver_Ready primitive signal to the transmitting port when it has free receive buffers
- End-to-End credit based flow control
 - Happens between initiator (source) and responder (destination) ports
 - ACK frames used to replenish credits

Fibre Channel

Class of Service

- Classes define communication strategy to use depending on the type of data to be transmitted
 - Determines the types of flow control used
- Class 1: Dedicated connection (reserves entire link on the path)
 - Only End-to-End flow control
- Class 2: Connectionless communication with end-to-end acknowledgements
 - Both buffer-to-buffer and end-to-end flow control
- Class 3: Connectionless with no end-to-end acknowledgements
 - Only buffer-to-buffer flow control
- Class 4: Similar to Class 1 except that only part of link bandwidth reserved for a VC
 - Only End-to-End flow control
- Class 5: Undefined
- Class 6: Multicast service
 - Only End-to-End flow control
- Intermix class: Allows Class 2 or Class 3 frames to be transmitted at times Class 1 frames are not being transmitted

Advanced Switching Interconnect (ASI)

- Link level flow control
- Endpoint source injection rate limiting
- Link bandwidth arbitration
- Reliable data transport
- Centrally managed fabric

Link level flow control

- Multiple "virtual channels" per link
 - Up to 16 unicast VCs: 8 are bypass-capable
 - Up to 4 multicast VCs
 - Separate flow control for each VC.
- Credit-based link-level flow control
 - Link VC receiver gives the VC transmitter a limit on how much data it can forward based on the currently free buffering in Rx.
 - The limit is updated as Rx buffering empties.
 - Credit update messages conveyed in a special, 4-byte packet.
- Optional Status-based flow control
 - Switches send output port status messages to their immediate upstream neighbors.
 - Provides one stage look-ahead view on congestion situation.
 - Upstream nodes temporarily stops data to congested downstream output ports.

Endpoint source injection rate limiting

- Optional
- Match bandwidth of
 - Slowest link along the path
 - Data rate of traffic's destination
- Source node maintains a connection queue (CQ) for each flow.
 - Per traffic class
 - Per path through the fabric
- Token bucket controls the flow of traffic through a CQ.
 - Shapes the transmission of the packets in time.
 - Allows limited burstiness
- Higher layer protocols determine the token bucket parameter values.
 - Protocols not specified in ASI spec., as far as I can see.

Link bandwidth arbitration

- Output port uses a "Minimum BW" scheduler to allocate link bandwidth to the active VCs.
 - An active VC has both waiting traffic and sufficient credits.
- Allocates available BW based on per-VC weighting parameter.
 - Allocates link BW on a data basis, not on a packet basis.
- Work conserving
 - If there is at least one packet ready to send, it will be sent.
 - Traffic is not held back for rate limiting/shaping at the switches.
- Specific algorithm not specified by ASI.

Reliable data transport

- Retransmission of lost or corrupted packets.
 - Minimum delay to detect the problem.
- Avoids the need to resend a lost packet end-to-end.
 - Reduces network BW needed to recover a lost packet.
 - Congestive collapse is almost impossible (?).

Centrally managed fabric

- Agents responsible for various aspects of fabric operation.
- Fabric-wide optimization and management
- Some fabrics have multiple paths between end nodes.
 - Important redundancy for failovers
 - For more BW or QoS than a single path can provide.
- Source routing of each packet allows traffic source to pick the route.
 - Potentially different for each traffic class * end-point.
 - Fabric manager typically discovers and defines the routes.
- Fabric management software may regulate access to the ASI fabric.
 - No new flows except when sufficient resources are available.

Ethernet

EthernetPriority and Flow Control

- Traffic Classes
 - Support for up to 8 priorities for traffic differentiation
- 802.3X flow control with link level granularity
 - Causes congestion spreading
 - Can not differentiate between latency sensitive and best effort traffic

Can Ethernet be the consolidated fabric?

- For Ethernet to be the consolidated fabric, it needs to have better congestion control capabilities including
 - End-to-End congestion control to push congestion to the sources
 - More granular link level flow control
 - Standard enhanced scheduling mechanisms

Backup