



# Scope Considerations for Link Layer in Resilient Packet Rings

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## **RPR Scopes & Issues**

- **Need to develop a forward-looking RPR standard which can scale to meet carrier & subscriber requirements for many years**
- **Focus on Factors that help RPR Acceptance in other Optical Networking Development Groups**
- **Control and Data Plane Separation to provide a Unified Data Transport**
- **Ring and Mesh Networks**

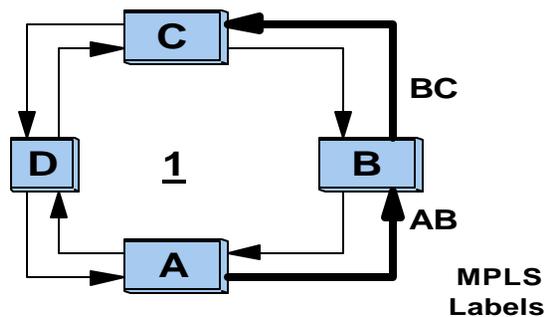
- **Multiple RPRs connected by common node(s) and point-to-point links**
- **Propose no L3 changes, instead provide L2 hooks**
- **Leverage MPLS as Foundation for Traffic Engineering & QoS operations - avoid re-inventions**
- **Different rates at different spans in an RPR Network**
- **Considerations for a Packet Delineation Mechanism**

- **Go top-down - important to first focus on scopes & requirements**
- **Then come up with a technology to achieve goals**
- **Suggest no bottoms-up approach - to force proprietary technologies to make them RPR**
- **A bottoms-up approach will limit our scope and functionality. High risk for RPR acceptance.**

- **Traffic Engineering, Load balancing, and QoS - leverage existing and continuing work in MPLS**
- **Packets are classified into different Flows at the Ingress Node - using MPLS, for instance**
- **No Need for Buffering Requirements in RPR Specification (creating two-priority buffers, for example, are way insufficient)**
- **Buffering and 'Arithmetic' Traffic Balancing at Intermediate Nodes Balances all Types of Traffic - not desirable**
- **Let Nodes participate in Bandwidth Allocation - and Leave Buffer & Priority Management to the Nodes (Nodes may have hundreds of queues & buffers to direct flows - it is a local decision). No other Protocols mandate buffering.**

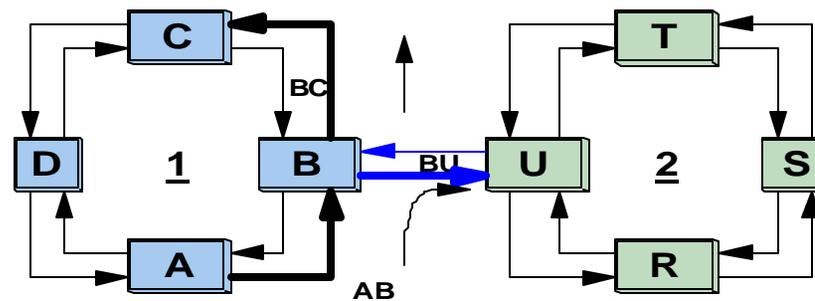
# NBMA Nature of Optical Networks

## SINGLE OPTICAL NETWORK



All Packets on A-B Link go to B

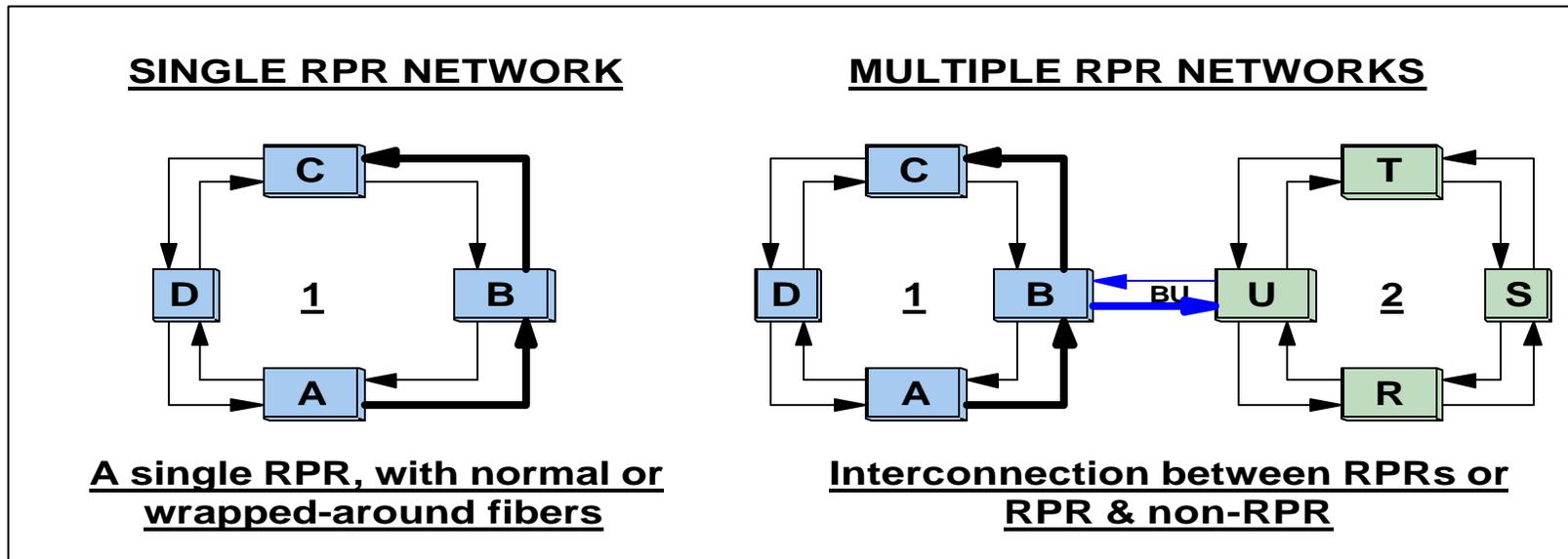
## MULTIPLE OPTICAL NETWORKS



All Packets on A-B Link go to B.  
B sends a packet to either U or C

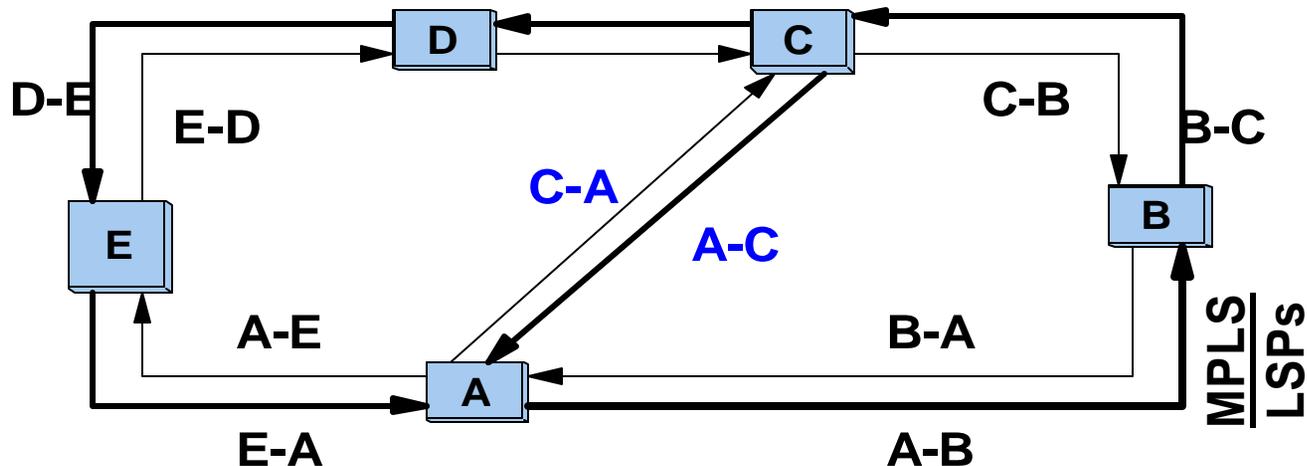
- Every packet passes through one and only one node at a time at the input.
- Since packets are Rx then Tx, a node ('B') is always a node in A-B-C, A-B-U-T Paths

# RPR MAC Address Resolution



- Since Packets go to nodes down the ring one node at a time, RPR still behaves in an NBMA mode, never actually as a 'Broadcast LAN'
- 'A' cannot resolve 'C's MAC addressing through ARP (unlike LAN)
- A single RPR may be shared by multiple router nodes, each belonging to a different IP domain
- Hence we must develop method(s) to resolve Remote Destination MAC Address (e.g. NHRP)

## RPR Links as (a set of) LSP(s)



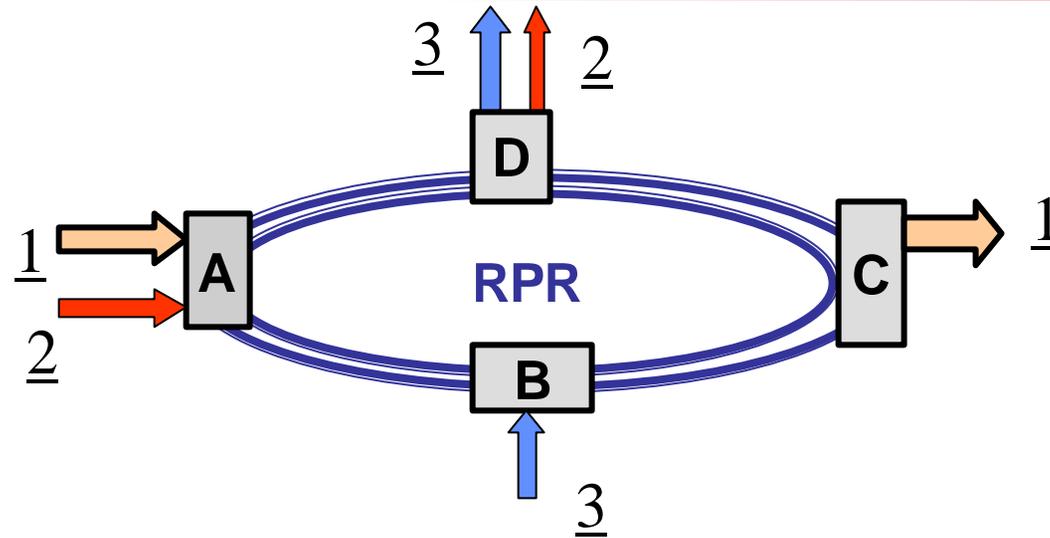
- An RPR Ring can be treated as one (or a collection of) LSPs.
- Any number of rings (not just 2) can be addressed using LSPs
- Mesh networks become just another type of LSP
- MPLS based fault protection and fault recovery MAY be used for RPR. Protection would easily work for 1:1 and 1:N, and any ring-mesh combination.
- Different traffic flows can easily be given different levels of protection



## **Control & Data Planes for RPR**

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- **Control plane provides all of RPR features: fault-recovery, traffic engineering, QoS, load balancing, etc.**
- **Control and Data Planes for RPR should be de-linked to simplify designs for Resilience Packet Rings.**
- **Use MAC addresses for Control Plane**
- **A unified control plane will allow traffic pooling for different data types, and minimize links and keep alive messages**
- **Allow Native Packets to travel on RPR. Multiservice transport becomes easy in RPR.**
- **Native packets from non-RPR networks can pass through RPR network.**



- Control plane may be used to set up MPLS path for a flow.
- Control Plane Packets may include TTL, etc. relevant for packet transport over rings
- Following header, normal Ethernet packet can be used
- MPLS LSPs are established for a flow.
- Link establishment, and fault detection



# RPR Header

RPR Core Header is 32 bits wide, with a 16 bits Header CRC



32 bits

16 bits





# RPR Control Packet

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**RPR control packets handle:**

- **Topology Discovery**
- **Fault Discovery, Isolation, Recovery, and Restoration**

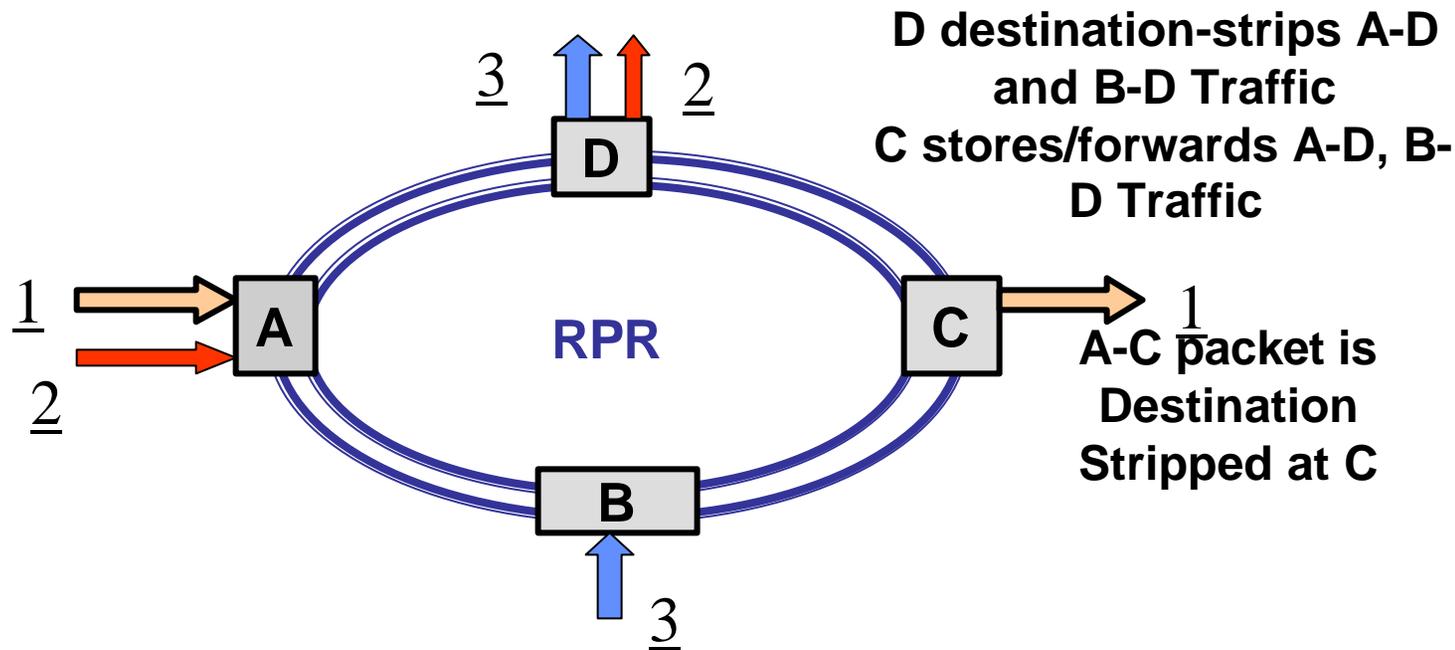


## **Data Plane: Addressing Methods**

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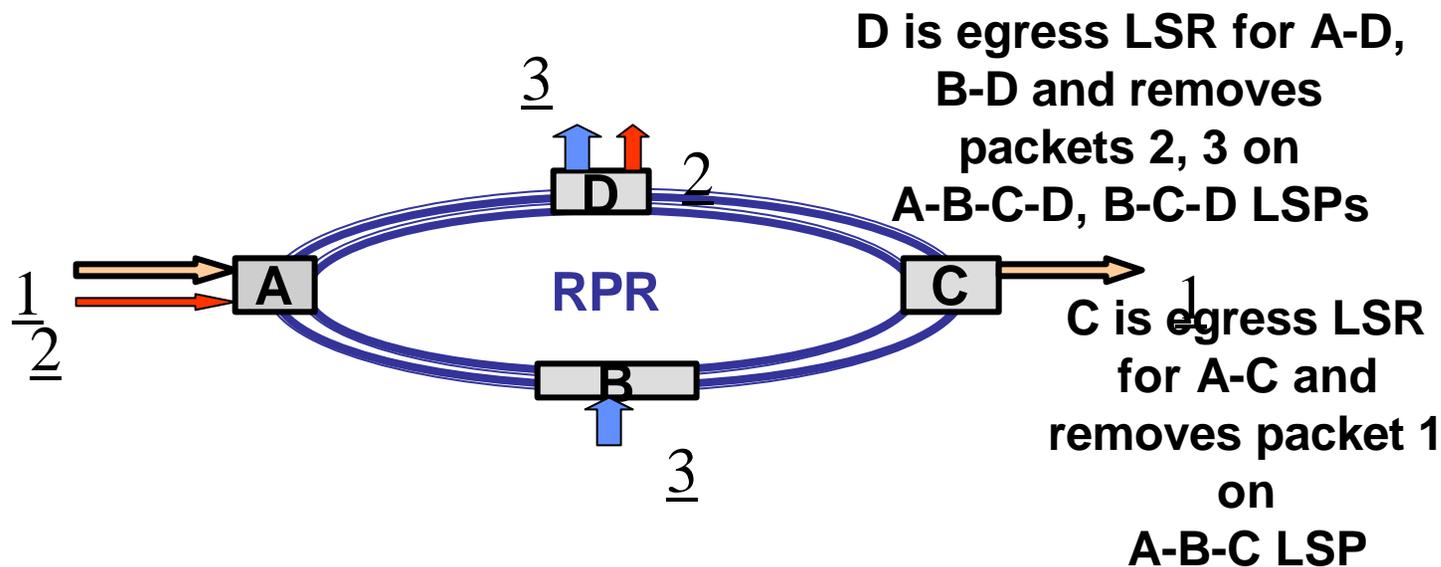
- **Once Source and Destination Nodes are identified, traffic flows can occur either using MAC address or MPLS labels**
- **MPLS Labels can be used as Data Link Layer Address on RPR for Data Packets**
- **Use of MAC addresses requires high-speed gigabit CAM at all nodes**
- **Since MPLS Labels have Local Significance only, simple Logic at Nodes to look up labels and classify packets into different queues and schedulers. No 48-bit lookups needed for high-speed data plane.**
- **Fine-grained traffic engineering and QoS without any extra protocol changes.**

## Data Plane: MAC-addressable Nodes

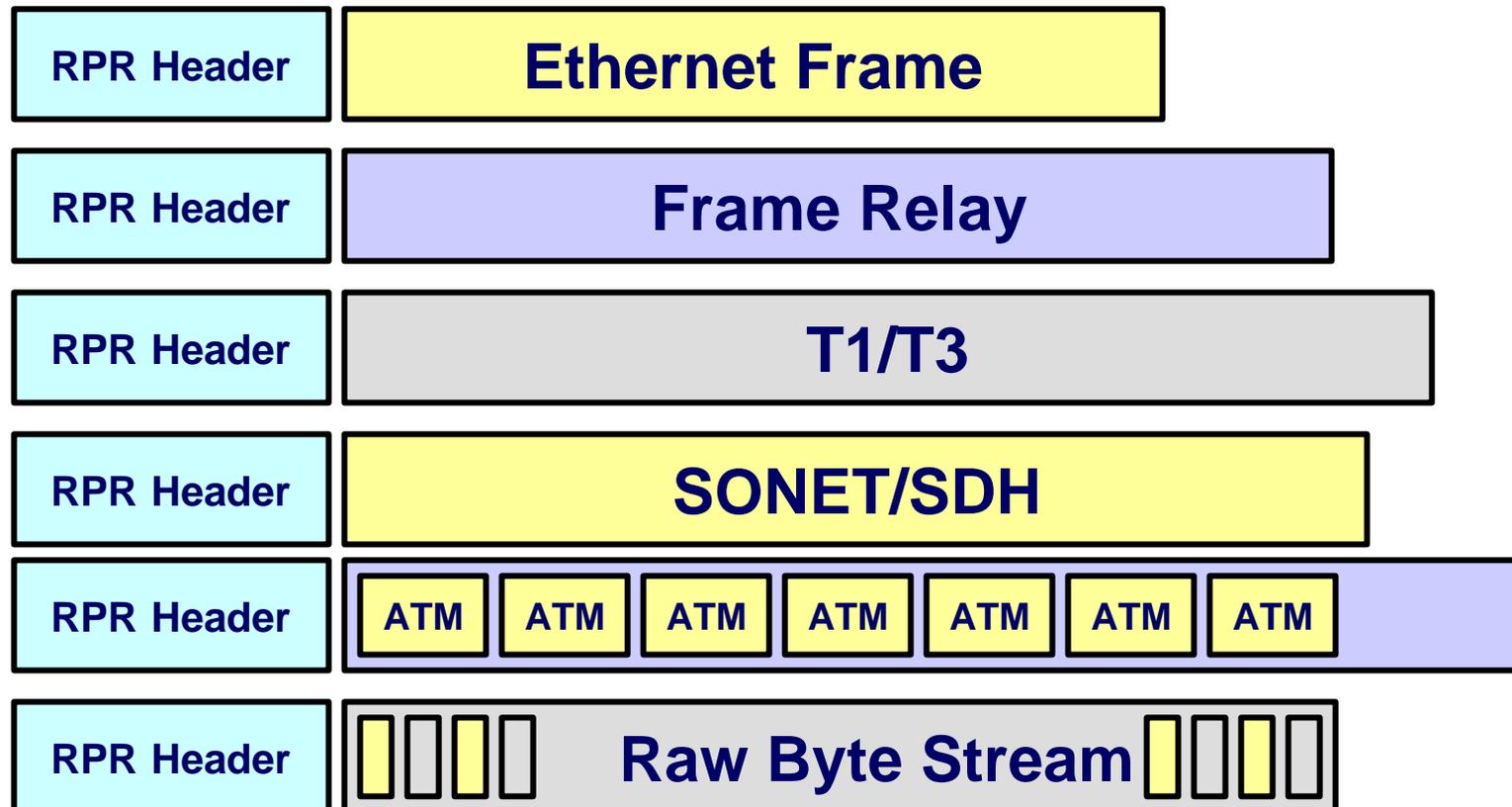


- For MAC addressing, { Dest MAC, Src MAC } used for Packet Transfer
- Mega/Gigabit CAM logic required at every intermediate node, since each node must sift through hundreds of passing destination MAC values to determine stripping

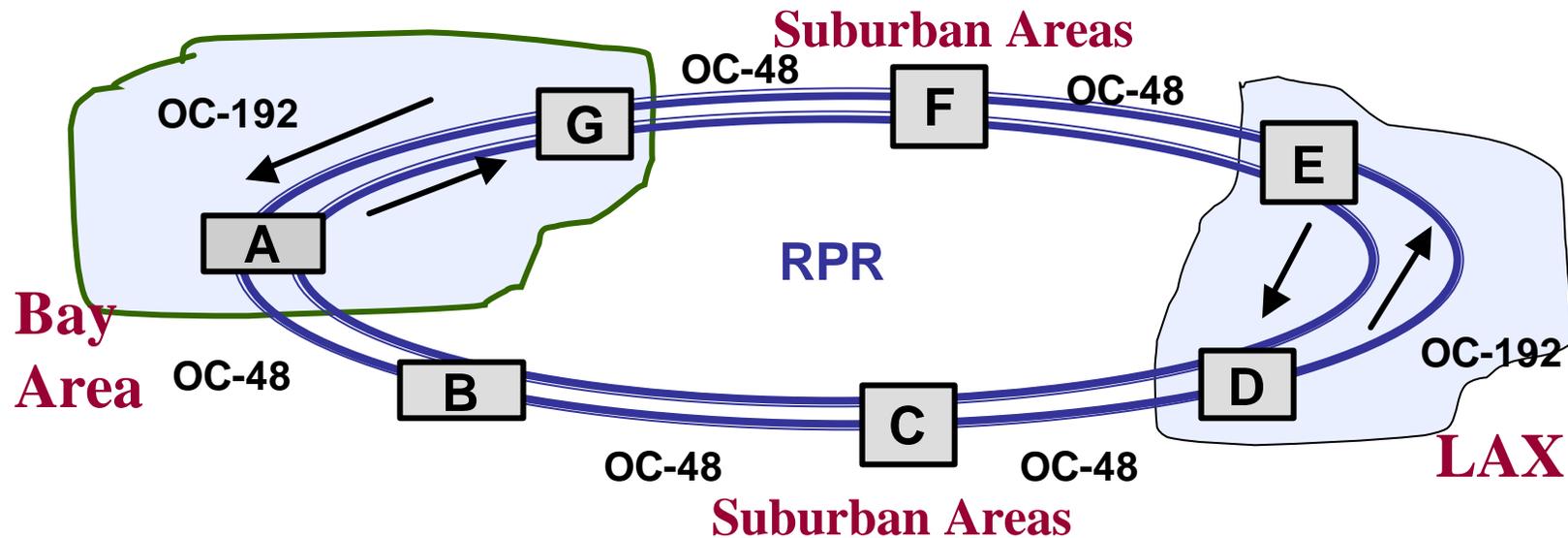
## Data Plane: RPR as MPLS LSP(s)



- MPLS Labels carry packet from node to node, until packet reaches egress
- MAC addresses not needed for data transport due to NBMA nature of Optical Networks.
- Both directions of fiber can be used for data traffic, if needed.
- Unidirectional nature of LSPs fits well with fiber ring spans



## Different Rates on RPR Spans



- Traditional SONET/SDH TDM networks required ALL nodes to be upgraded to move to a higher speed. Expensive & Time Consuming
- With packet rings, it should be possible to increase rate of a section while keeping older nodes intact, with no impact to rest of the ring.
- When one fiber breaks, traffic flows are adjusted to fit the single fiber traffic rate.
- MPLS LSPs are unidirectional in nature, like optical fiber. Easy to send data packets on different LSPs with different speeds.

- **RPR must choose a robust frame delineation method**
- **Frame Delineation must provide a high recovery rate in high BER situations**
- **Avoid having to look at and de-stuff every byte - this is particularly problematic (high-speed hardware needed) at 10G+ speeds**
- **RPR packets can be sent on SONET/SDH as well as direct fiber.**
- **Two considerations compared here are - HDLC and SDL**



- **Byte-by-byte destuffing required**
- **Poor processing at higher speeds**
- **Loss of a single 0x7E can become a hopeless situation, since same value used both for SOP and EOP**
- **Hard to allocate queues/buffers - length of packet not known in advance.**
- **Poor delineation**



# Robust Transport with SDL

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## Simple Data Link (SDL) Protocol (rfc2823)

- Robust framing for any packet data over a point-to-point or a ring network.
- Packets are delimited using length/CRC construct instead of an HDLC-type framing with 0x7E at both ends
- Optional 16/32-bit CRC at the end of the packet
- No need to perform byte-by-byte de-stuffing at high optical data transport speeds.

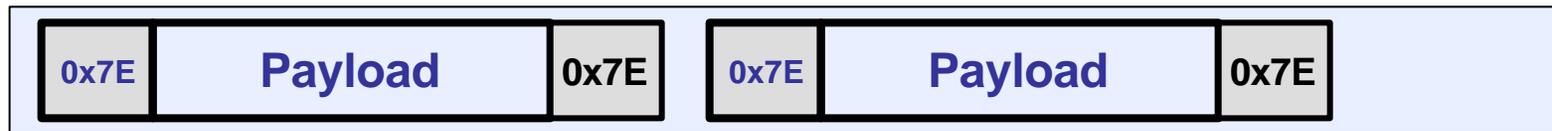




# Packet Delineation

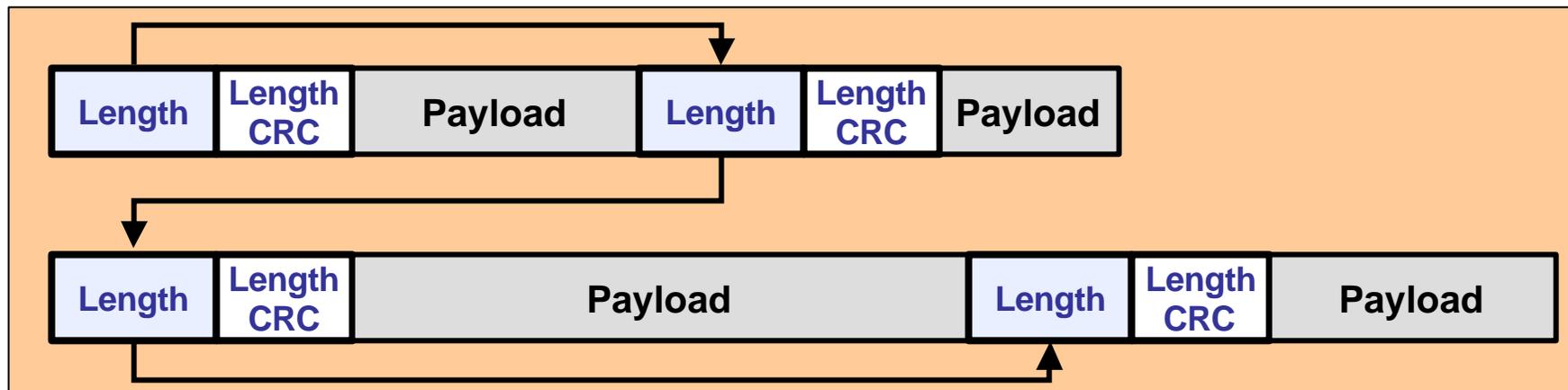
## SDL vs. HDLC

### HDLC



Byte-by-byte destuffing and checking for 0x7E and other control patterns  
Loss of any 0x7E becomes a hopeless situation for recovery

### SDL



Framers can jump Length bytes to get to next frame. No byte-by-byte lookup.

- **RPR has a perfect opportunity to provide a unified mechanism for multiservice data transport over optical ring networks**
- **RPR would work for ring & mesh networks**
- **Native mapping for IEEE802.3 packets**
- **RPR can become the first protocol to formulate and allow use of MPLS as a L2 protocol for optical networks**
- **Use all existing protocols developed for Traffic Engineering with no modifications**
- **Native Data Transport for any Type of Data**