## 12. Input and output (I/O) of intervals

12.1. Overview. This standard specifies conversion from a text string that holds an interval literal to an interval interval to a program (input), and the reverse (output). The method by which strings are read from, or written to, a character stream is language- or implementation-defined.

Containment shall hold on input and output so that, when a program computes an enclosure of some quantity given an enclosure of the data, a correctly written program can ensure this holds all the way from text data to text results.

In addition to normal I/O, the standard requires each interval type  $\mathbb{T}$  to have an *exact representation*. This has two parts: operations to convert any internal  $\mathbb{T}$ -interval  $\boldsymbol{x}$  to a string  $\boldsymbol{s}$ , and back again to recover  $\boldsymbol{x}$  exactly; and documentation of the mathematical relation between  $\boldsymbol{s}$  and the Level 1 interval value of  $\boldsymbol{x}$ .

12.2. Input. Input is provided for each supported bare or decorated interval type  $\mathbb{T}$  by the  $\mathbb{T}$ -version of text2interval(s), where s is a string, as specified in §11.11.8. It accepts an arbitrary interval literal s and returns a  $\mathbb{T}$ -interval enclosing the Level 1 value of s.

[Note. This provides the basis for free-format input of interval literals from a text stream, as might be provided by overloading the >> operator in C++.]

12.3. Output. An implementation shall provide an operation

## $\texttt{interval2text}(\boldsymbol{X}, cs)$

where cs is optional. X is a bare or decorated interval datum of any supported interval type  $\mathbb{T}$ , and cs is a string, the conversion specifier. The operation converts X to a valid interval literal string s, see§11.11.1, which shall be related to X as follows, where Y is the Level 1 value of s.

- (i) Let  $\mathbb{T}$  be a bare type. Then Y shall contain X, and shall be empty if X is empty.
- (ii) Let  $\mathbb{T}$  be a decorated type. If X is NaI then Y shall be NaI. Otherwise, write  $X = x_{dx}$ ,  $Y = y_{dy}$ . Then
  - y shall contain x, and shall be empty if x is empty.
  - dy shall equal dx, except in the case that dx = com and overflow occurred, that is, x is bounded and y is unbounded. Then dy shall equal dac.

[Note. Y being a Level 1 value is significant. E.g., for a bare type  $\mathbb{T}$ , it is not allowed to convert  $X = \emptyset$  to the string garbage, even though converting garbage back to a bare interval at Level 2 by  $\mathbb{T}$ -text2interval gives  $\emptyset$ , because garbage has no Level 1 value as a bare interval literal.]

The tightness of enclosure of X by the value Y of s is language- or implementation-defined. If present, cs controls the layout of the string s and its syntactic parts, see§11.11.1, for instance

l, u, m and r where relevant. The standard does not specify cs but among the user-controllable elements should be the following:

- (i) The overall field width (the length of s), and whether output is in inf-sup or uncertain form.
- (ii) How Empty, Entire and NaI are output, e.g., whether lower or upper case, and whether Entire becomes [Entire] or [-Inf, Inf].
- (iii) For l, u, m and r, the field width and the number of places after the point or the number of significant figures. There should be a choice of radix.

If cs is absent, output should be in a general-purpose layout (analogous, e.g., to the g specifier of **fprintf** in C). There should be a value of cs that selects this layout explicitly.

[Note. This provides the basis for free-format output of intervals to a text stream, as might be provided by overloading the << operator in C++.]

If  $\mathbb{T}$  is a 754-conforming bare type, there shall be a value of cs that produces behavior identical with that of interval2exact, below. That is, the output is an interval literal that, when read back by text2interval, recovers the original datum exactly.

12.4. Exact representation. For any supported bare interval type  $\mathbb{T}$  an implementation shall provide operations interval2exact and exact2interval. Their purpose is to provide a portable text representation of every bare interval datum. They are characterized by the recovery requirement:

For any T-datum x, the value s = interval2exact(x) is a string, such that y = exact2interval(s) is a T-datum equal to x.

[Note. This is equality at Level 2. x and y may differ at Level 3, e.g., a zero endpoint might be stored as -0 in one and +0 in the other.]

If  $\mathbb{T}$  is a 754-conforming type, the string s shall be an interval literal. The operation exact2interval shall be the same as text2interval for this type. If x is nonempty, s shall be of inf-sup form [l, u] with l and u represented exactly. If  $\mathbb{T}$  is a binary type, l and u if finite shall be represented in the hexadecimal-significant form of 754§5.12.3. If  $\mathbb{T}$  is a decimal type, they shall be represented as decimal numbers.

If  $\mathbb{T}$  is not 754-conforming, there are no restrictions on the form of the string s apart from the above recovery requirement. However, the algorithm by which interval2exact converts x to s is regarded as part of the definition of the type and shall be documented by the implementation.

For any type  $\mathbb{T}$ , since exact2interval inputs a string s and outputs a bare interval, it is a constructor similar to text2interval, having a set V of valid strings that it regards as denoting an interval. This set includes by definition the set  $V_0$  of all strings that can be output by interval2exact. For a 754-conforming type, V is the set of valid interval literals, and is larger than  $V_0$  since any such literal can be written in many equivalent ways. For other types  $\mathbb{T}$ , the documentation shall specify the set V, which might or might not be strictly larger than  $V_0$ .

As with text2interval (see §11.11.8), exact2interval shall return  $\emptyset$  and raise a languageor implementation-defined exception when its input is invalid.

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