

5. The decoration system at Level 1

5.1 Overview

An implementation makes the decoration system available by providing:

- a decorated version of each interval extension of an arithmetic operation, of each interval constructor, and of some other operations; and
- various auxiliary functions, e.g., to extract the interval and decoration parts, and to apply a standard initial decoration to an interval.

The decoration system is specified here at a mathematical level, with the finite-precision aspects presented in Clause 6. Subclauses 5.2, 5.3, and 5.4 give the basic concepts; 5.5 and 5.6 define how intervals are given an initial decoration, and how decorations are bound to library interval arithmetic operations to give correct propagation through expressions; 5.7 is about non-arithmetic operations.

5.2 Definitions and properties

A decoration d is a property (that is, a boolean-valued function) $p_d(f, \mathbf{x})$ of pairs (f, \mathbf{x}) , where f is a real-valued function with domain $\text{Dom}(f) \subseteq \mathbb{R}^n$ for some $n \geq 0$ and $\mathbf{x} \in \overline{\mathbb{IR}}^n$ is an n -dimensional box, regarded as a subset of \mathbb{R}^n .

The notation (f, \mathbf{x}) unless said otherwise denotes such a pair, for arbitrary n , f and \mathbf{x} . Equivalently, d is identified with the set of pairs for which the property holds:

$$d = \{ (f, \mathbf{x}) \mid p_d(f, \mathbf{x}) \text{ is true} \}. \quad (4)$$

The set \mathbb{D} of decorations has five members:

Value	Short description	Property	Definition
com	common	$p_{\text{com}}(f, \mathbf{x})$	\mathbf{x} is a bounded, nonempty subset of $\text{Dom}(f)$; f is continuous at each point of \mathbf{x} ; and the computed interval $f(\mathbf{x})$ is bounded.
dac	defined & continuous	$p_{\text{dac}}(f, \mathbf{x})$	\mathbf{x} is a nonempty subset of $\text{Dom}(f)$, and the restriction of f to \mathbf{x} is continuous.
def	defined	$p_{\text{def}}(f, \mathbf{x})$	\mathbf{x} is a nonempty subset of $\text{Dom}(f)$.
trv	trivial	$p_{\text{trv}}(f, \mathbf{x})$	always true (so gives no information).
ill	ill-formed	$p_{\text{ill}}(f, \mathbf{x})$	Not an Interval; formally $\text{Dom}(f) = \emptyset$, see 5.3.

These are listed according to the propagation order (10), which may also be thought of as a quality-order of (f, \mathbf{x}) pairs—decorations above **trv** are “good” and **ill** is “bad”.

A **decorated interval** is a pair, written interchangeably as (\mathbf{u}, d) or \mathbf{u}_d , where $\mathbf{u} \in \overline{\mathbb{IR}}$ is a real interval and $d \in \mathbb{D}$ is a decoration. (\mathbf{u}, d) may also denote a decorated box $((\mathbf{u}_1, d_1), \dots, (\mathbf{u}_n, d_n))$, where \mathbf{u} and d are the vectors of interval parts \mathbf{u}_i and decoration parts d_i , respectively.

The set of decorated intervals is denoted by $\overline{\mathbb{DIR}}$, and the set of decorated boxes with n components is denoted by $\overline{\mathbb{DIR}}^n$.

When several named intervals are involved, the decorations attached to $\mathbf{u}, \mathbf{v}, \dots$ are often named du, dv, \dots for readability, for instance (\mathbf{u}, du) or \mathbf{u}_{du} , etc.

An interval may be called a **bare** interval to emphasize that it is not a decorated interval.

Treating the decorations as sets as in (4), **trv** is the set of all (f, \mathbf{x}) pairs, and the others are nonempty subsets of **trv**. By design, they satisfy the **exclusivity rule**

$$\text{For any two decorations, either one contains the other or they are disjoint.} \quad (6)$$

Namely, the definitions (5) give:

$$\text{com} \subset \text{dac} \subset \text{def} \subset \text{trv} \supset \text{ill}, \quad \text{note the change from } \subset \text{ to } \supset \quad (7)$$

$$\text{com}, \text{dac} \text{ and } \text{def} \text{ are disjoint from } \text{ill}. \quad (8)$$

Property (6) implies that for any (f, \mathbf{x}) there is a unique tightest (in the containment order (7)) decoration, such that $p_d(f, \mathbf{x})$ is true, called the **strongest decoration of (f, \mathbf{x})** , or of f over \mathbf{x} , and written $\text{Dec}(f | \mathbf{x})$. That is,

$$\text{Dec}(f | \mathbf{x}) = d \iff p_d(f, \mathbf{x}) \text{ holds, but } p_e(f, \mathbf{x}) \text{ fails for all } e \subset d.$$

- 1 NOTE—Like the exact range $\text{Rge}(f | \mathbf{x})$, the strongest decoration is theoretically well-defined, but its value for a
2 particular f and \mathbf{x} may be impractically expensive to compute, or even undecidable.

3 5.3 The ill-formed interval

4 An ill-formed decorated interval is also called NaI, **Not an Interval**. Conceptually, there shall be only one
5 NaI. Its interval part has no value at Level 1.

6 The **ill** decoration results from invalid constructions and propagates unconditionally through arithmetic
7 expressions. Namely, the **ill** decoration arises as a return value of

- 8 – a constructor when it cannot construct a valid decorated interval, or of
- 9 – a library arithmetic operation if and only if one of its inputs is ill-formed.

10 Formally, **ill** may be identified with the property $\text{Dom}(f) = \emptyset$ of (f, \mathbf{x}) pairs.

11 [Example. The constructor call `numsToInterval(2, 1)` is invalid, so its decorated version returns NaI.]

12 Information may be stored in a NaI in an implementation-defined way (like the payload of an IEEE 754
13 floating-point NaN), and functions may be provided for a user to set and read this for diagnostic purposes.

14 An implementation may provide means for an exception to be signaled when a NaI is produced.

15 5.4 Permitted combinations

16 A decorated interval \mathbf{y}_{dy} shall always be such that

$$\mathbf{y} \supseteq \text{Rge}(f | \mathbf{x}) \text{ and } p_{dy}(f, \mathbf{x}) \text{ holds, for some } (f, \mathbf{x}).$$

17 If $dy = \text{dac}, \text{def}$ or com , then by definition \mathbf{x} is nonempty, and f is everywhere defined on it, so that $\text{Rge}(f | \mathbf{x})$
18 is nonempty, implying \mathbf{y} is nonempty. Hence the decorated intervals

$$\emptyset_{\text{dac}}, \emptyset_{\text{def}}, \text{ and } \mathbf{x}_{\text{com}} \text{ if } \mathbf{x} \text{ is empty or unbounded}$$

19 are contradictory—implementations shall not produce them.

20 No other combinations are forbidden.

21 5.5 Operations on/with decorations

22 5.5.1 Initializing

A bare interval is initialized with a decoration by the operation

$$\text{newDec}(\mathbf{x}) = \mathbf{x}_d \quad \text{where} \quad d = \begin{cases} \text{com} & \text{if } \mathbf{x} \text{ is nonempty and bounded,} \\ \text{dac} & \text{if } \mathbf{x} \text{ is unbounded, and} \\ \text{trv} & \text{if } \mathbf{x} \text{ is empty.} \end{cases}$$

23

5.5.2 Disassembling and assembling

For a decorated interval x_{dx} , the operations `intervalPart(x_{dx})` and `decorationPart(x_{dx})` shall be provided, with value x and dx , respectively. For the case of NaI, `decorationPart(NaI)` has the value `ill`, but `intervalPart(NaI)` has no value at Level 1.

Given an interval x and a decoration dx , the operation `setDec(x, dx)` returns the decorated interval x_{dx} if this is an allowed combination. The cases of forbidden combinations are as follows:

- `setDec(\emptyset, dx)`, where dx is one of `def`, `dac` or `com`, returns \emptyset_{trv} ;
- `setDec(x, com)`, for any unbounded x , returns x_{dac} ; and
- `setDec(x, ill)` for any x , whether empty or not, returns NaI.

5.5.3 Comparisons

For decorations, comparison operations for equality `=` and its negation `≠` shall be provided, as well as comparisons `>`, `<`, `≥`, `≤` with respect to the propagation order (10).

5.6 Decorations and arithmetic operations

Given a scalar point function φ of k variables, a **decorated interval extension** of φ —denoted here by the same name φ —adds a decoration component to a bare interval extension of φ . It has the form $w_{dw} = \varphi(v_{dv})$, where $v_{dv} = (v, dv)$ is a k -component decorated box $((v_1, dv_1), \dots, (v_k, dv_k))$. By the definition of a bare interval extension, the interval part w depends only on the input intervals v ; the decoration part dw generally depends on both v and dv . In this context, NaI is regarded as being \emptyset_{ill} .

The definition of a bare interval extension implies

$$w \supseteq \text{Rge}(\varphi | v) \quad (\text{enclosure}).$$

The decorated interval extension of φ determines a dv_0 such that

$$p_{dv_0}(\varphi, v) \text{ holds} \quad (\text{a “local decoration”}). \quad (9)$$

It then evaluates the output decoration dw by

$$dw = \min\{dv_0, dv_1, \dots, dv_k\}, \quad (\text{the “min-rule”}),$$

where the minimum is taken with respect to the **propagation order**:

$$\text{com} > \text{dac} > \text{def} > \text{trv} > \text{ill}. \quad (10)$$

5.7 Decoration of non-arithmetic operations

5.7.1 Interval-valued operations

This give interval results but are not interval extensions of point functions.

- the cancellative operations `cancelPlus(x, y)` and `cancelMinus(x, y)` of 4.5.3;
- the set-oriented operations `intersection(x, y)` and `convexHull(x, y)` of 4.5.4

No one way of decorating these operations gives useful information in all contexts. Therefore, a *trivial* decorated interval version is provided as follows. If any input is NaI, the result is NaI; otherwise the corresponding operation is applied to the interval parts of the inputs, and its result decorated with `trv`. The user may replace this by an appropriate nontrivial decoration via `setDec()`, see 5.5, where this can be deduced in a given application.

1 **5.7.2 Non-interval-valued operations**

2 These give non-interval results:

3 – the numeric functions of 4.5.6 and

4 – the boolean-valued functions of 4.5.7.

5 For each such operation, if any input is NaI, the result has no value at Level 1. Otherwise, the opera-
6 tion acts on decorated intervals by discarding the decoration and applying the corresponding bare interval
7 operation.