

Why an additional Shaper?

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Structure of this Presentation

Recap of Shaper: CBS, TAS and BLS

Handling of Overload-Frames

Comparison: Time Aware Shaper (TAS) vs. Burst Limiting "Shaper" (BLS)
 Assumptions for the Comparison
 Optimization Goals for industrial/automotive control networks
 Previous Simulations
 Calculation of optimized TAS Windows's
 Simulation results
 Simplified TAS Windows calculation "algorithm"

Conclusion

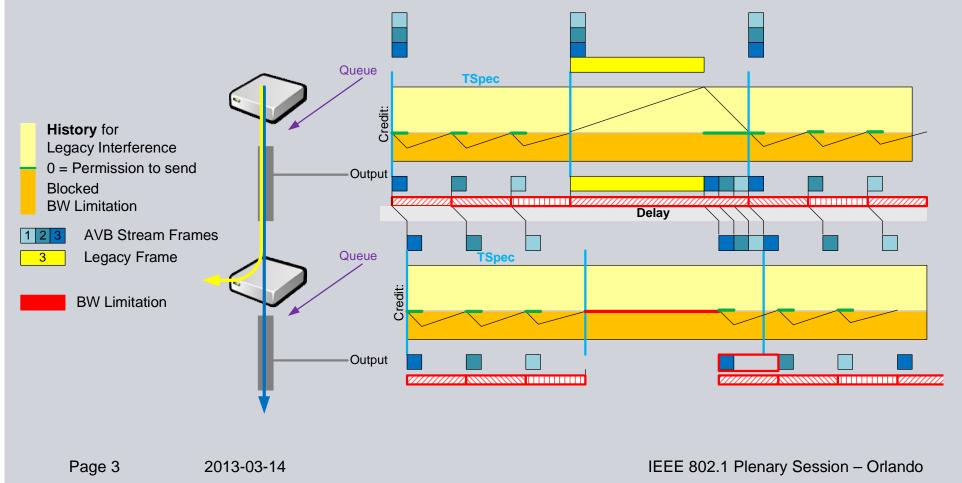
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Recap of CBS BW Limitation and Delay

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The Credit limits the bandwidth over a "frame" time

The Credit must be set to 0, if no AVB Frame is in the queue and stays zero!



TAS Shaper Pro- and Cons

The Time Aware Shaper (TAS):

http://www.ieee802.org/1/files/public/docs2012/bv-boiger-time-aware-shaper-0912-v02.pdf

For TAS all network devices must be synchronized

- Synchronized End Stations for scheduled sending times
- Synchronized Switches with defined forwarding times

Pros:

- Best possible latency (immediately forwarding)
- No Jitter in arrival time (no congestion)

Cons:

- Reservation of Bandwidth must be exclusive (not useable for other Traffic)
- Effort for Calculation of scheduling

Performance of TAS compared to Preemption with Bandwidth Limitation

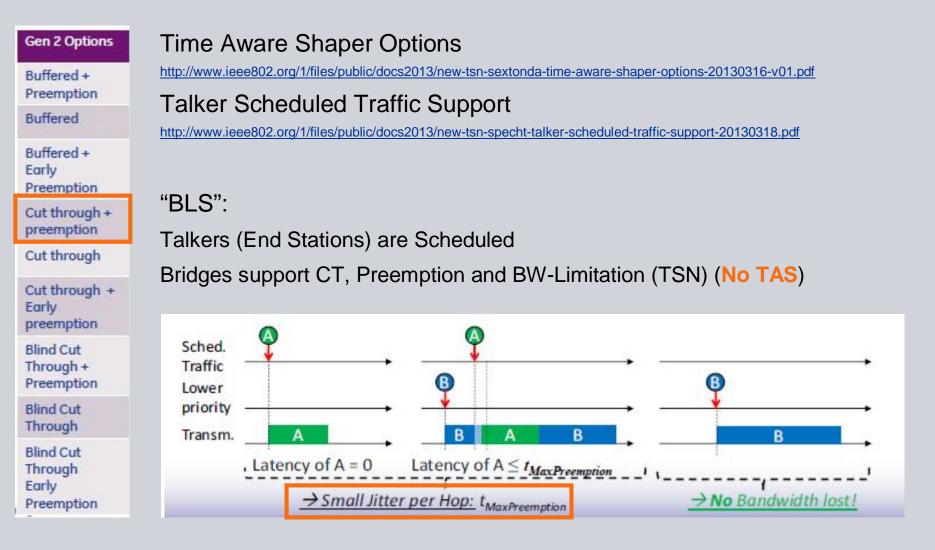


The Burst Limiting "Shaper" (BLS):

http://www.ieee802.org/1/files/public/docs2013/new-goetz-TSN-4-Industrial-Networks-20130115-v1.pdf

- limits the bandwidth usage (ensures reservation of bandwidth and resources – AVB core feature)
- streams use highest priority
 - (~strict priority with highest priority transmitted immediately if frames available)
- is used with preemption
- (avoids waiting time due to congestion)

Other Names of these Shapers



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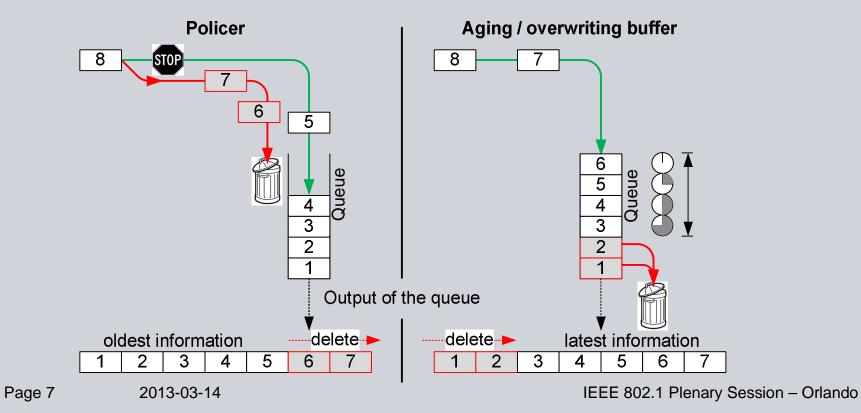
Bandwidth Limitation Policer against aging



The limitation of the bandwidth can be done in different ways:

- use a policer to prevent frames getting in the queue
- "aging" / overwriting of frames inside the queue

But should only effect the transmission in case of errors (e.g. wrong configuration)



Performance of TAS compared to Preemption with Bandwidth Limitation

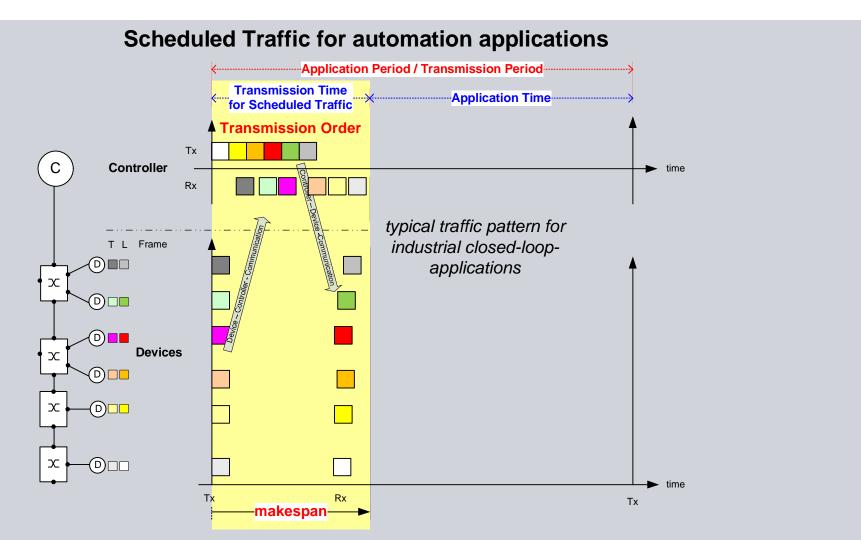
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ASSUMPTIONS for the comparison of the two Shaper:

- End Stations are always synchronized to get repeatable result (Not needed for BLS – only for comparison of achievable performance)
- Switches are only synchronized when using the TAS Shaper (TAS only works with accurate synchronized End station and Switches)
- Goal is a minimal time for transmitting all information
- Additional Transmission Delay for "Legacy" Traffic should be low
- Waste of Bandwidth should be minimal (e.g. Waste using Guard Time)
- Best achievable Performance is compared so that L2 must be used

Recap: Low Latency is required to minimize Transmission Time for Scheduled Traffic





Transmission of Scheduled Traffic within transmission time for Scheduled Traffic

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Previous Simulations

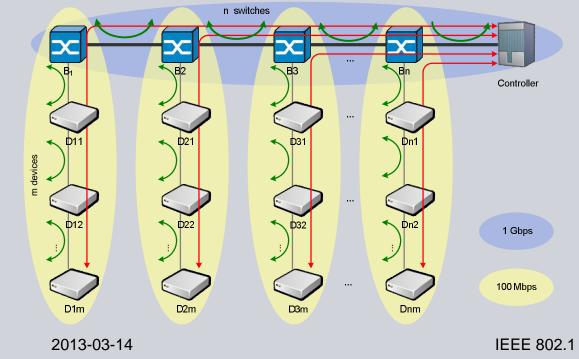
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Performance without optimized TAS windows against BLS:

(http://www.ieee802.org/1/files/public/docs2013/new-goetz-TSN-4-Industrial-Networks-20130115-v1.pdf)

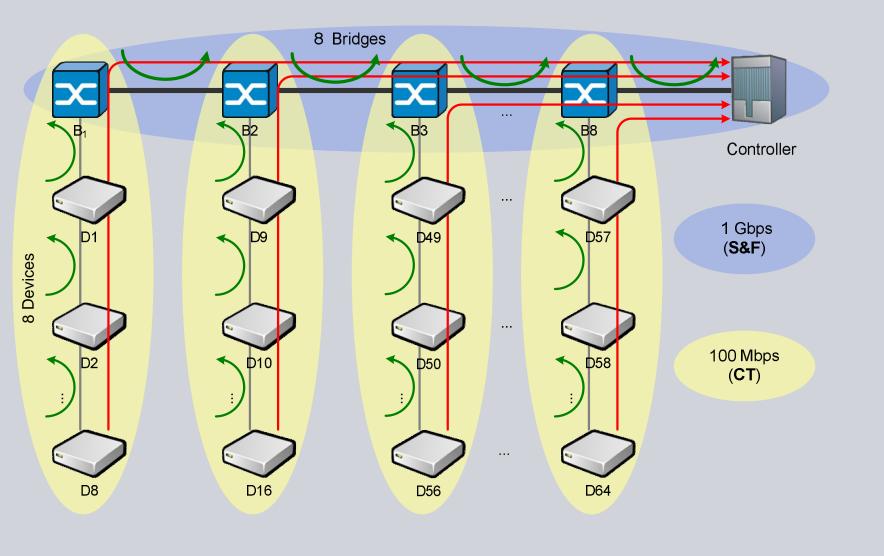
Simulated Use Case: Time aware Shaper (TAS) <-> Burst limiting Shaper (BLS)

a) Low Latency for Scheduled Traffic with constant Frame Size
b) Low Latency for Scheduled Traffic with random Frame Size
c) Low Latency for Scheduled Traffic with random Frame Size
and optimized Window Size for TAS



Simulation - Latency for Scheduled Traffic TAS w/ multiple windows in GE for D->C





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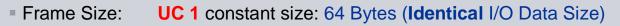
Latency for Scheduled Traffic Time aware Shaper (TAS)<-> Burst limiting Shaper (BLS)

General Settings (1):

Network: 8 bridges, 8x8 devices (bridged end stations)

Real time application (synchronized)

- Transmission order (C->D): farthest first, nearest last
- Traffic load for Scheduled Traffic < 50%</p>



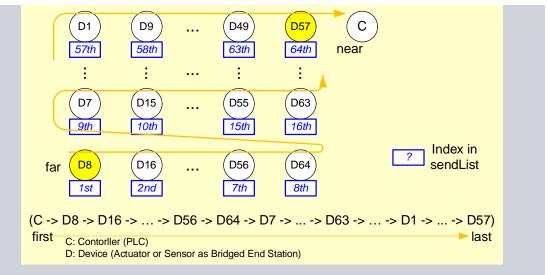
UC 2 random size: 10% 64 Bytes / 10% 512 Bytes / 80% between 128~384 Bytes (Random I/O Data Size)

Best effort traffic:

- Traffic load < 30%</p>
- Frame size: 25% max_size, 25% min_size, 50% between 250~1250 Bytes

25% burst (frames in chain), 75% non-burst

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Latency for Scheduled Traffic Time aware Shaper (TAS)<-> Burst Limiting Shaper (BLS)

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General Settings (2)

BLS & TAS

 Transmission period: 250 us for Scheduled Traffic with constant frame size, 500 us for Scheduled Traffic with random frame size,

•Window size (only for TAS):

- **UC 1** Window size is 72 us for Scheduled Traffic with constant frame size
- **UC 2** Window size is 400 us for Scheduled Traffic with random frame size
- **Optimized UC1 and UC 2** Window size optimized
 - Window start time always at the beginning of cycle
 - Window close time varies for different location of station
 - Window close time right after the station has transmitted the last Scheduled frame
- Cut-through only for Scheduled Traffic: 48 bytes Delay

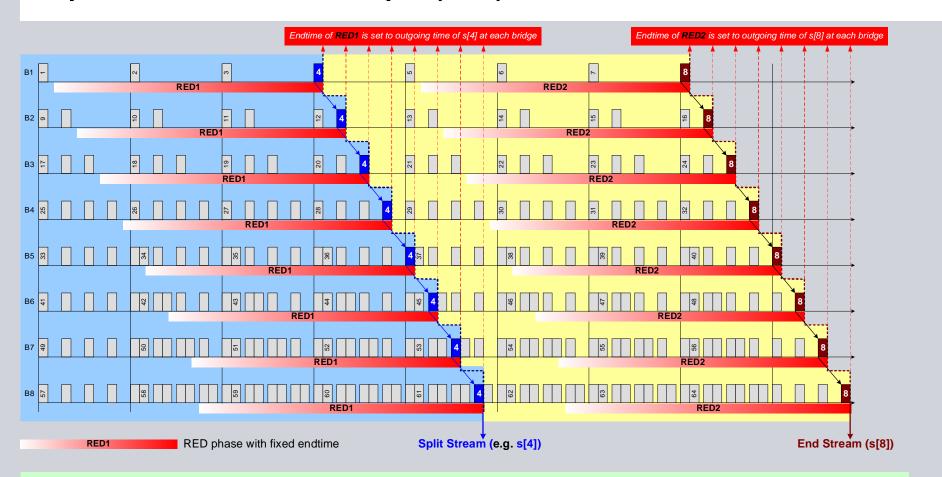
Bridging delay: 500 ns; cable + PHY delay: 750 ns

Pre-emption in combination with TAS or BLS



Goal: find a scheduling scheme splitting the long RED phase into e.g. two short ones for each bridge, in order to reduce bandwidth waste, while keeping worst-case latency (e.g. at stream[8]) unchanged.

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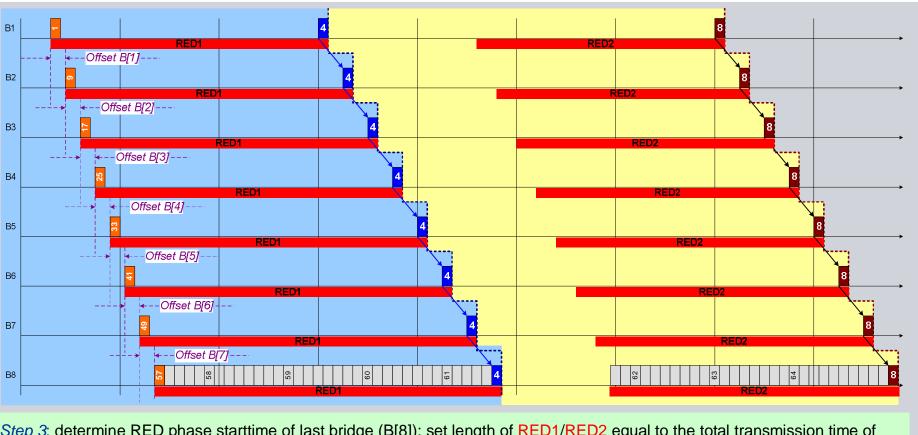


<u>Step 1:</u> locate a split stream (from farthest branch) and divide the RT transmissions into two segments (of similar loads) <u>Step 2:</u> align the **EndTime** of RED1 and RED2 with the outgoing time of split stream and end stream at each bridge

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Step 3: determine RED phase starttime of last bridge (B[8]): set length of RED1/RED2 equal to the total transmission time of

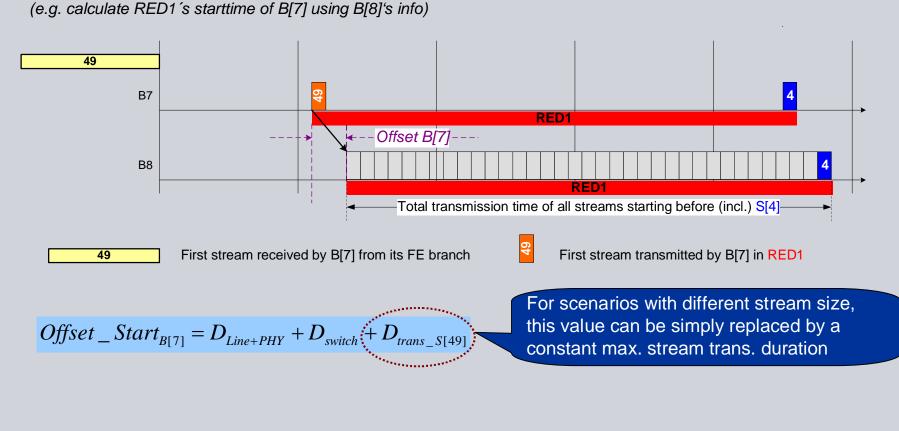
all streams that start before (incl.) split stream (S[4])/end stream (S[8]) How to determine this value? Step 4: calculate RED phase starttime of rest bridges recursively in reverse order using $StartTime_{RED_B[i]} = StartTime_{RED_B[i+1]} Offset_{B[i]}$

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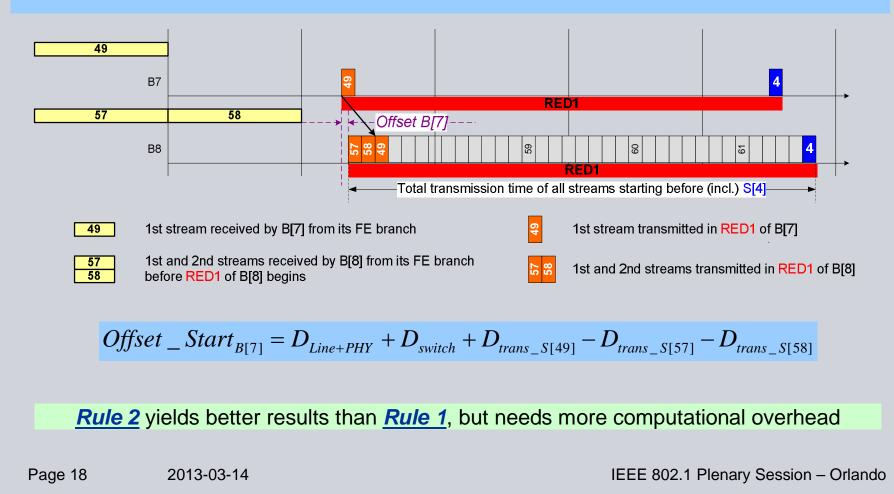


Calculation of Offset_{B[i]} <u>Rule 1</u>: Red phase of B[i] must start at an earlier time than that of B[i+1], so that the first outgoing stream of B[i] is ready for transmission at the beginning of Red phase of B[i+1]



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Calculation of Offset_{B[i]} <u>*Rule 2*</u>: length of Red phase of B[i] calculated by <u>*Rule 1*</u> can be further shortened by considering the RT transmissions that are already scheduled at the beginning of B[i+1]'s RED phase



UC 1 Simulation - Low Latency for Scheduled Traffic **SIEMENS** with constant Frame Size of 64Bytes, CT@GE

-Cycle = 250 us-B1 72 us 178 us 72 us B2 178 us TAS w/o opt. 72 us 178 us B3 Β4 72 us 178 us B5 72 us 178 us 72 us B6 178 us **B**7 72 us 178 us B8 72 us 178 us B1 193 32 us 56.68 us TAS w/ endtime opt. B2 191 40 us 58 60 us B3 189 47 us 60.53 us B4 62.45 us 187 55 us 64.37 us 185 63 us B5 B6 66.29 us 183 71 us 181 79 us B7 68.21 us 179 86 us B8 70.14 us TAS w/ 2 red phases* B1 15.3 us B2 16.64 us B3 17.99 us Β4 19.33 us 13.94 us B5 20.67 us 15.28 us B6 22.02 us 16.63 us B7 23.36 us 17.97 us B8 25 36 us 18 64 us

* Red phases calculated with Rule 2

(Drawn to scale)

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UC 1 Simulation - Low Latency for Scheduled Traffic **SIEMENS** with constant Frame Size of 64Bytes, CT@GE

ldentical 64bytes		Bridg	Cycle e length (us)	Red Phase Length (us)	Green Phase Length (us)	Max. RT trans. delay (us)	NRT-GE E2E delay (us) (mean, max)	Max RT queue size (at GE bridges)	Gained bandwidth for NRT vs. w/o opt (relative to cycle len.)	Consumed bandwidth for RT trans. (us) per cycle	Bandwidth waste in RED phase (relative to cycle len.)
	BLS		250	n/a	n/a	73.03	15.50 (151.59)	3 (at B[7,8])	n/a	n/a	n/a
TAS		1			178	70.89	37.00 (327.27)	1 (at all bridges)	0.00%	5.38	26.65%
		2								10.75	24.50%
		3								16.13	22.35%
		4	250	72						21.50	20.20%
	w/o op	5	250	12						26.88	18.05%
		6								32.26	15.90%
		7								37.63	13.75%
		8								43.01	11.60%
		1		56.68	193.32		31.73 (268.71)	1 (at all bridges)	6.13%	5.38	20.52%
		2		58.60	191.40	70.89			5.36%	10.75	19.14%
	Opt. c	^{of} 3		60.53	189.47				4.59%	16.13	17.76%
	Red Phase	4	250	62.45	187.55				3.82%	21.50	16.38%
	end	5	200	64.37	185.63				3.05%	26.88	15.00%
	time	6		66.29	183.71				2.28%	32.26	13.61%
		7		68.21	181.79				1.51%	37.63	12.23%
		8		70.14	179.86				0.75%	43.01	10.85%
		1		22.23	227.77	70.89	16.01 (174.32)	3 (at B[7,8])	19.91%	5.38	6.74%
		2		25.58	224.42				18.57%	10.75	5.93%
	ant w/	, 3		28.94	221.06				17.22%	16.13	5.12%
	opt w two Re	4	250	32.29	217.71				15.88%	21.50	4.32%
	phase	1 6	250	34.98	215.02				14.81%	26.88	3.24%
	phase	6		37.66	212.34				13.74%	32.26	2.16%
		7		40.35	209.65				12.66%	37.63	1.09%
		8		43.02	206.98				11.59%	43.01	0.00%

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UC 1 Simulation - Low Latency for Scheduled Traffic **SIEMENS** with random Frame Size of 64Bytes, CT@GE



* Red phases calculated with Rule 2

(Drawn to scale)

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UC 1 Simulation - Low Latency for Scheduled Traffic **SIEMENS** with random Frame Size of 64Bytes, CT@GE

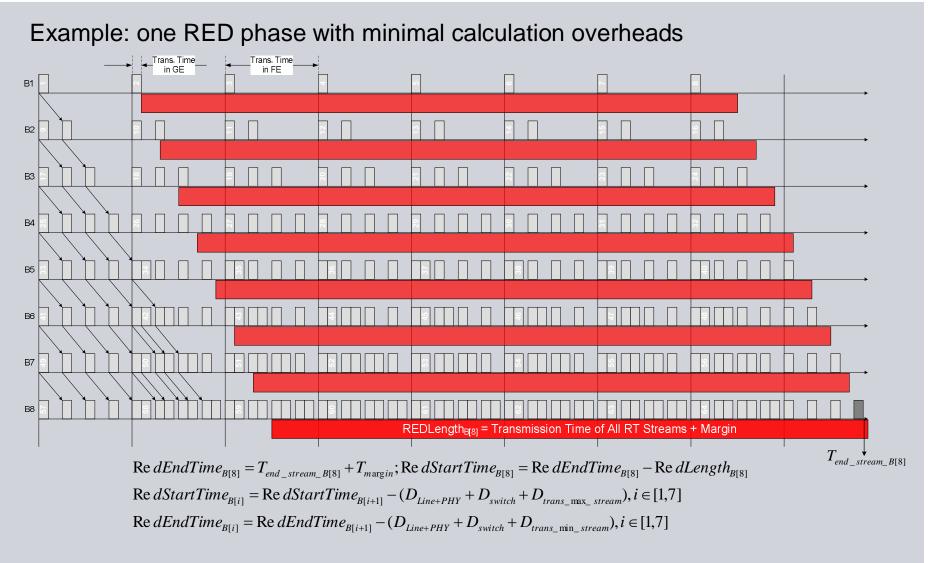
Random Size (seed = 0)		Bridge	Cycle length (us)	Red Phase Length (us)	Green Phase Length (us)	Max. RT trans. delay (us)	NRT-GE E2E delay (us) (mean, max)	Max RT queue size (at GE bridges)	Gained bandwidth for NRT vs. <mark>w/o opt</mark> (relative to cycle len.)	Consumed bandwidth for RT trans. (us) per cycle	Bandwidth waste in RED phase (relative to cycle len.)
В	BLS		500	n/a	n/a	239.36	23.49 (331.95)	4 (at B[8])	n/a	n/a	n/a
	w/o opt.	1	500	238	262	237.2	109.12 (602.03)	4 (at B[8])	0.00%	20.82	43.44%
		2								35.58	40.48%
		3								52.90	37.02%
		4								72.27	33.15%
		5			202	251.2				88.62	29.88%
		6								106.87	26.23%
		7								129.31	21.74%
		8								147.25	18.15%
	Opt. of Red Phase end time	1	500	213.06	286.94		99.29 (577.98)	4 (at B[8])	4.99%	20.82	38.45%
TAS		2		216.46	283.54	237.2			4.31%	35.58	36.18%
		3		219.87	280.13				3.63%	52.90	33.39%
		4		223.27	276.73				2.95%	72.27	30.20%
		5		226.67	273.33				2.27%	88.62	27.61%
		6		230.07	269.93				1.59%	106.87	24.64%
		7		233.47	266.53				0.91%	129.31	20.83%
		8		236.96	263.04				0.21%	147.25	17.94%
	opt w/ two Red phases	1	500	94.78	405.22	237.2	31.92 (252.93)	4 (at B[2,5,6,7,8])	28.64%	20.82	14.79%
		2		102.12	397.88				27.18%	35.58	13.31%
		3		108.46	391.54				25.91%	52.90	11.11%
		4		114.73	385.27				24.65%	72.27	8.49%
		5		121.27	378.73				23.35%	88.62	6.53%
		6		129.80	370.20				21.64%	106.87	4.59%
		7		140.21	359.79				19.56%	129.31	2.18%
		8		147.26	352.74				18.15%	147.25	0.00%

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Calculation of TAS Window Size



TAS Optimization: Simplified Algorithms for Engineering Systems

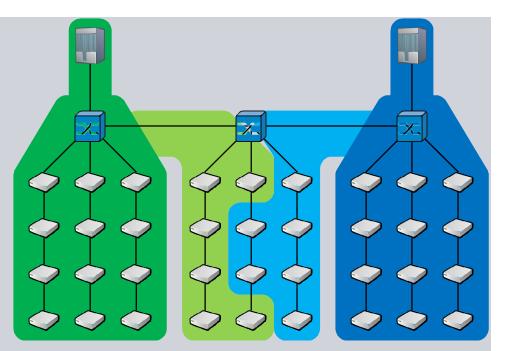


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Conclusion

TAS

- Leads to optimal latency
- Window Size Optimizations lead to better legacy performance
- Easy to calculate TAS windows in small topologies (line, star)
- Huge Effort to calculate TAS windows in **complex** Topologies



BLS

- can lead with an optimal Scheduling to nearly same performance (Preemption)
- legacy performance is good without optimizations
- Doesn't require time aware scheduling inside the bridges
- Work's also without scheduled send times in end devices (lower performance)

Unique Names are needed for the Traffic Class / Shaper combination

Thank you for your attention!



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