

# Tutorial: The Time-Synchronization Standard from the AVB/TSN suite

## IEEE Std 802.1AS™-2011 (and following)

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# Abstract

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**This presentation provides an overview of clock and time synchronization across heterogeneous networks using the published standards:**

**IEEE Std 802.1AS™-2011 and**

**IEEE Std 802.11™-2012 Timing Measurement**

**With special focus on wireless / 802.11 links**

# Agenda

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- **802.1AS goals / applications**
- Overview
- Grand Master selection
- Time propagation
  - Media-independent
  - Media-dependent
- Time **\*inside\*** the system

# 802.1AS goals

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## Distribute a single, accurate, time reference that is optimized for audio/video synchronization

- **Accurate**
  - 802.1AS time in participating LAN-attached stations is accurate to within 1us
    - If one station thinks it's 09:57:55 AM PST, the others agree within specified # of network hops
    - Assuming +/-100PPM crystal, 1PPM / second maximum frequency drift
    - Accuracy has been physically demonstrated in even near-worst-case LAN environments
- **One time reference for the entire LAN/Subnet**
  - LAN-agnostic architecture
    - Can be applied to any 802-compatible LAN
      - Is explicitly defined for 802.3 (Ethernet), 802.11 (Wi-Fi), MoCA, ITU-T G.hn, and EPON.
    - Contains a Profile of IEEE Std.1588™ -2008 (PTP, or Precision Time Protocol)
      - 1588-2008 reserves a special value for the 802.1AS profile (see Transport Specific in IEEE Std. 1588™-2008)
  - LAN-specific measurements
    - 802.3 measurements follow generic 1588
    - 802.11 measurements leverages IEEE Std. 802.11™ -2012 “Timing Measurement”
      - Formerly known as IEEE Std. 802.11v™ -2011
    - Coordinated Shared Network (CSN) measurements defined also
- **Plug and play**
  - Grand Master (GM) clock is selected automatically
    - Time stabilizes in a fraction of a second
  - Clock tree reconfigures automatically if GM is lost

# These are **NOT** explicit 802.1AS goals/characteristics

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## **NO!:** Improve the latency of packets/frames

Time *can* be used to do this, but 802.1AS is silent on the topic (see 802.1 TSN)

## **NO!:** Improve the delivery jitter of packets/frames

Time *can* be used to do this, but 802.1AS is silent on the topic (see 802.1 TSN)

## **NO!:** Constrain delivery jitter of 802.1AS packets

The protocol is effectively immune to jitter, since each delay is explicitly measured

## **NO!:** Force every media clock in the [W]LAN to be equal

...though many times it's desirable. Related media clocks are typically described indirectly by cross-referencing 802.1AS and the samples using presentation times

## **NO!:** Work over Ethernet 802.3 only (or 802.11 only)

The application should not care what the physical network is

## **NO!:** Limited to 7 network hops

Networks of extremely long chains of bridges/relays have been deployed

# How 802.1AS is used with audio/video

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- **First, 802.1AS time-sync is NOT needed where:**
  - *Stored A/V* content is streamed to *one* renderer
    - Renderer can pull content from storage as it is needed
  - *Self contained A/V* system (network isn't in the path)
- **Shared time IS needed:**
  - When there are multiple renderers or sources
  - Where the renderer cannot tell the source to slow down or speed up, yet latency must be very small
- **The “Presentation Time” uses 802.1AS time to port the media rate**
  - “Render audio sample #896 when 802.1AS-time is 8:11:35.023712”
  - Every “presentation time” provides render time of specific sample
  - Every pair of “presentation times” communicates the actual “rate”
    - Nominal media rate (e.g., 44.1KHz audio, or 165MHz for HDMI) is assumed known
    - But the actual media rate always differs by a non-zero PPM, and changes over time due to thermal and other factors

# How 802.1AS is used in industrial/IOT

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- **First, 802.1AS time-sync is NOT needed where:**
  - Coordination BETWEEN networked devices is unnecessary
  - *All the data comes from and is acted upon by a single device*
- **Shared time IS needed:**
  - In cyber-physical systems (CPS)
    - For coordination of distributed inputs, outputs, computation, and communication
  - In sensor networks where big-data analytics are applied to the aggregate
    - ...without losing temporal information
    - Especially for
      - Low-power, infrequently-communicating devices
      - Cooperative Diversity in wireless communications
      - Devices responsible cooperatively for beamforming or array processing
  - Where communication+computation latency exceeds coordination requirement

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# High level architecture

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- **Grand Master selection**

- GM-capable stations advertise themselves via ANNOUNCE messages
- If a station hears from station with “better” clock, it does not send ANNOUNCE
  - Configurable “Priority” field can override clock quality
  - MAC address is tie breaker
- Time relays drop all inferior ANNOUNCE messages
  - Forward only the best
- Last one standing is Grand Master for the domain
  - GM is the root of the 802.1AS timing tree
  - GM periodically sends the current time

- **Propagation of time**

- Non-leaf devices in the tree propagate time toward the leaves
  - Taking queuing delay into account (aka “Residence Time”)

# 802.1AS is not alone

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- **802.1AS is one of three core 802.1 AVB standards**
  - The others define a reservation protocol and traffic shaping algorithm
- **802.1AS over Ethernet (802.3) qualifies as a Profile of IEEE 1588-2008**
  - 802.1AS simplifies 1588, adds other media (e.g. 802.11)
  - Also other optimizations
    - Better stability over large networks (no cascaded PLLs)
    - Faster GM failover

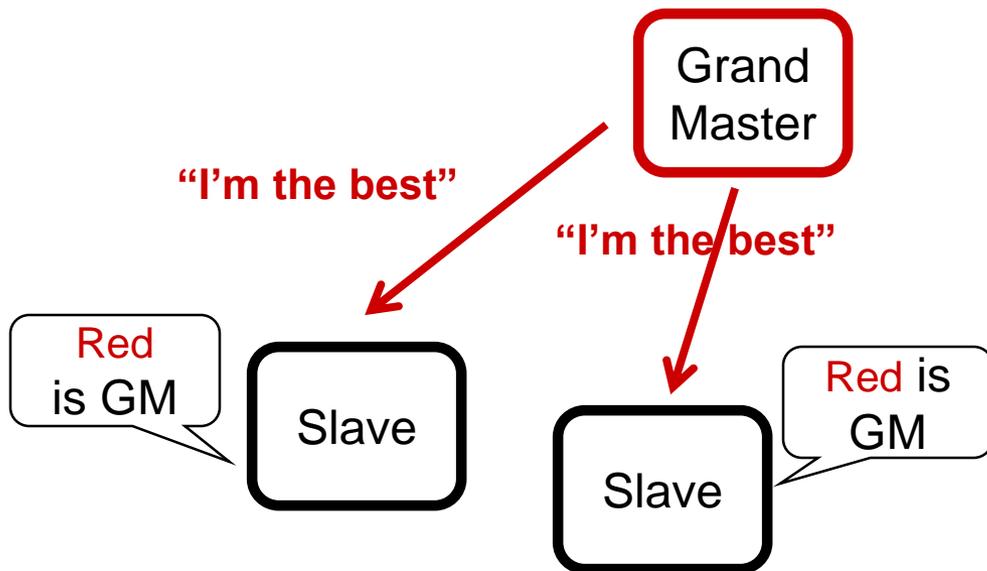
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# Grand Master selection – Steady state

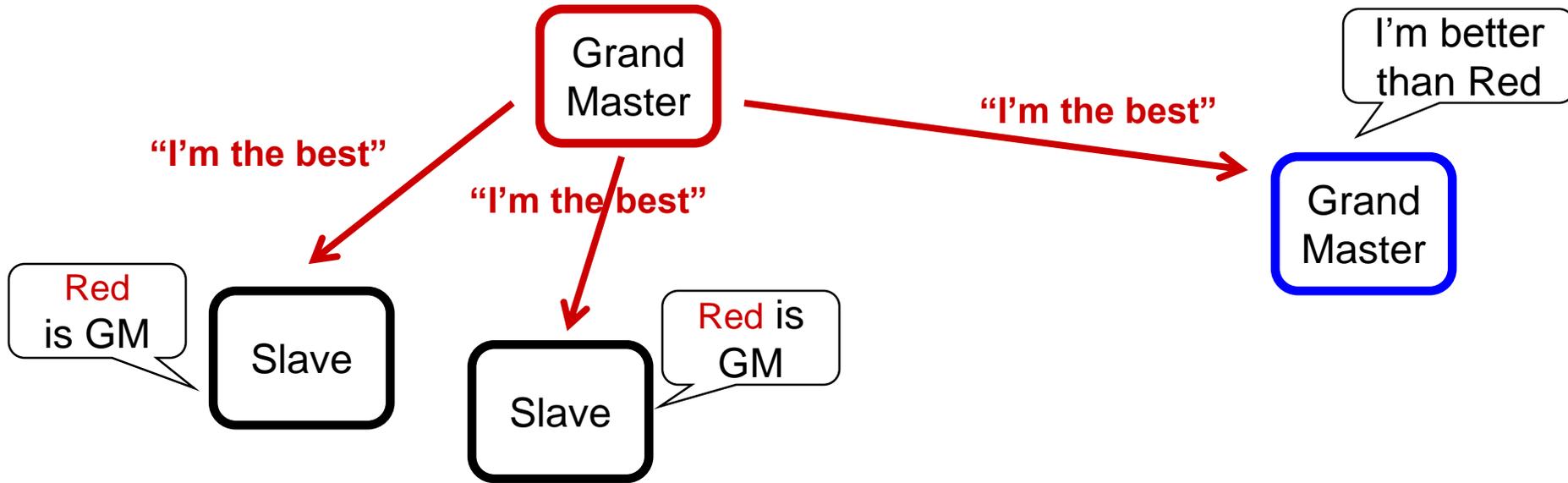
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## Steady state:

- **The GM sends ANNOUNCE**
- **Slaves acknowledge the GM is best (including self)**
  - ...using simple bitwise compare
  - And do not send ANNOUNCE messages
- **Everyone knows their role**
- **Life is good**

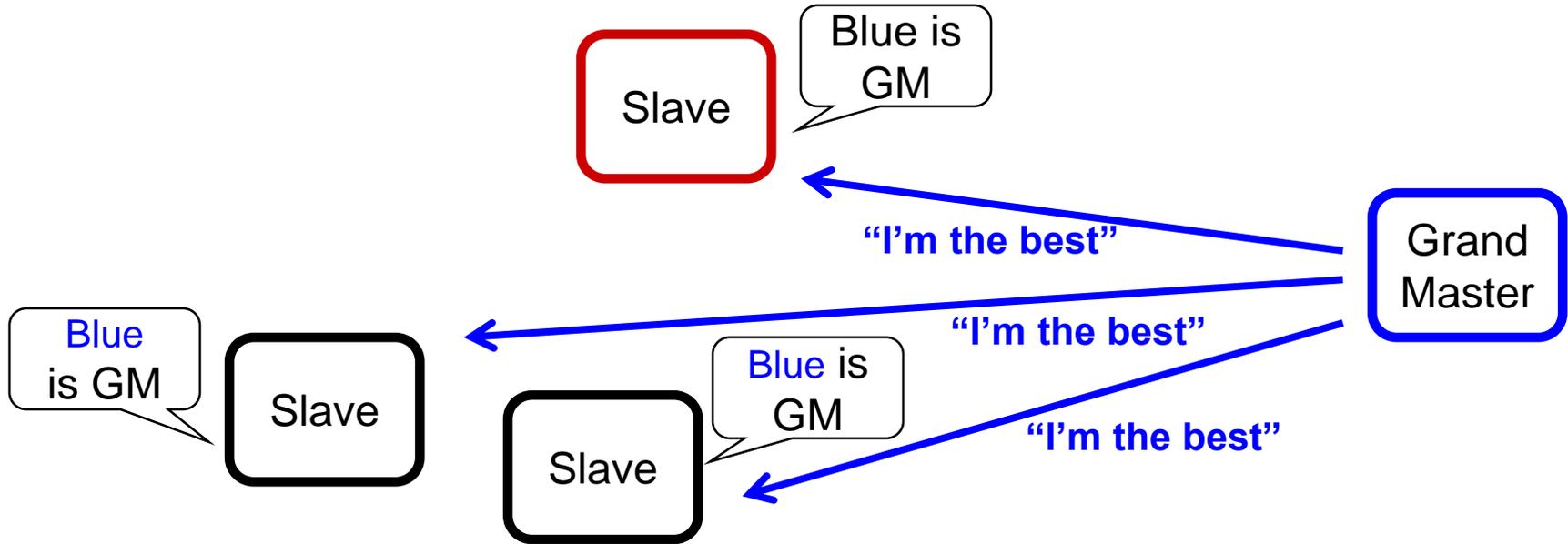
# Grand Master selection – New, better GM (1)



- **Blue station with better clock appears**

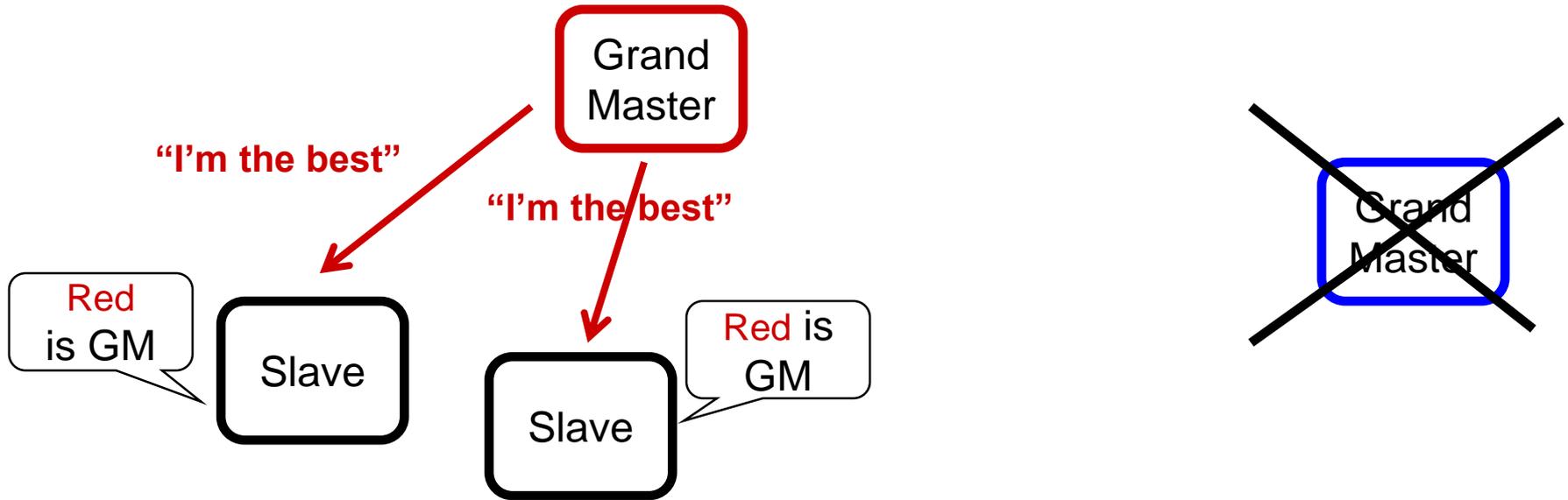


# Grand Master selection – New, better GM (3)



- **Blue station with better clock appears**
- **Blue sends ANNOUNCE**
- **Stations all realize blue is superior**
- **Red stops sending ANNOUNCE**
- **Blue is quickly the undisputed GM**

# Grand Master selection – Lost GM



- If Blue disappears, all GM-capable stations send **ANNOUNCE**
- Eventually, only Red sends **ANNOUNCE**

# ANNOUNCE comparison

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- **The credentials passed in ANNOUNCE messages are compared (in order of decreasing importance):**
  - Priority (settable by management)
  - Multiple “**quality of my clock**” fields
  - MAC address

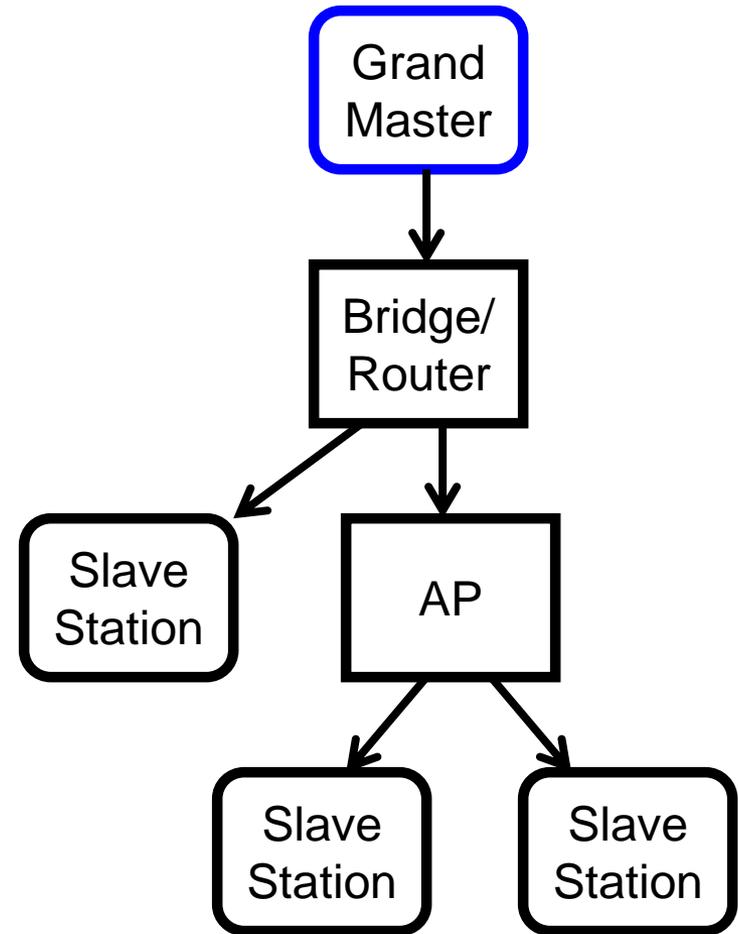
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# Time domain is propagated to the “subnet”

- **Time-aware Relays**
  - The best ANNOUNCE
  - The GM’s time
- **Relays can be a**
  - Bridge
  - Router
  - Linux computer with 2 Ethernet, a wireless and MoCA NIC(s)
- **Links can be**
  - 802.3 Ethernet
  - 802.11 WiFi
  - Almost anything compatible with IEEE 802
- **Let’s look at one relay**

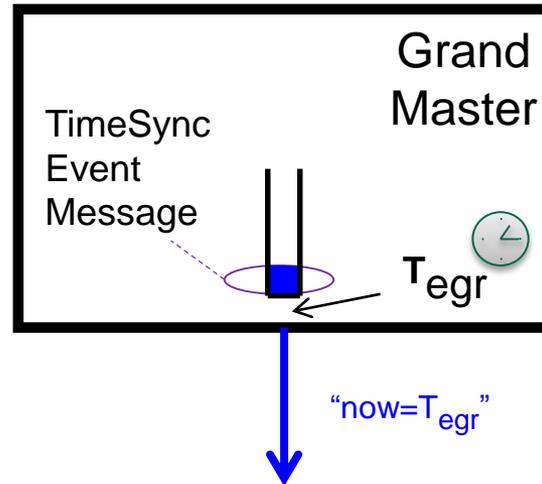


# Grand Master selection – Relays help decide



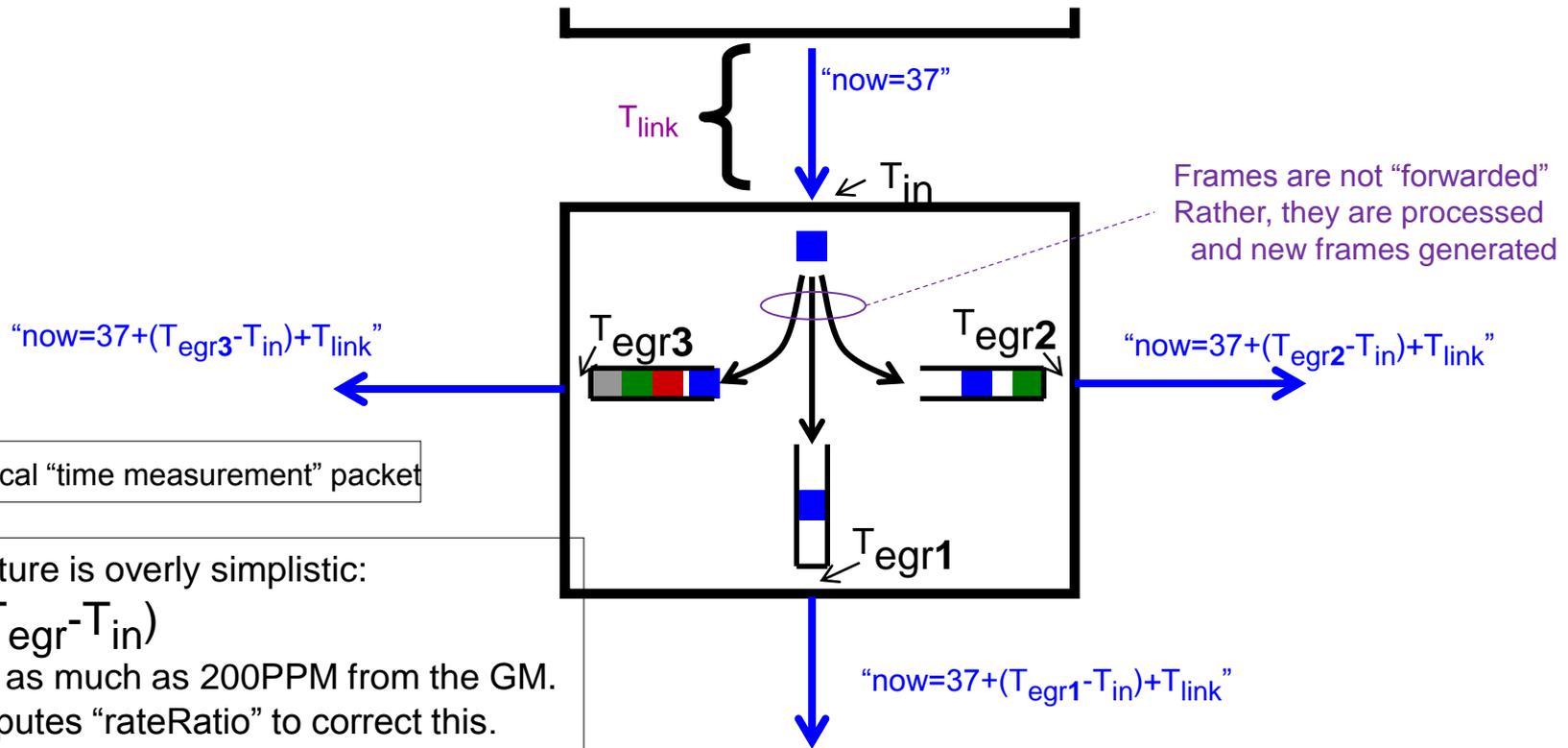
- **Relays filter inferior ANNOUNCE**
  - Relays forward best ANNOUNCE messages
  - Relays drop inferior ANNOUNCE messages
  - Benefits:
    - Reduction in network traffic
    - Faster GM selection
- **Announce messages establish the Clock Tree**

# Grand Master responsibilities



- **Transmit an Event message each Sync Interval**
- **Hardware captures the egress, or Tx time ( $T_{egr}$ )**
- **Pass  $T_{egr}$  downstream**
  - Typically in a Follow-Up message
- **GM time need not be synchronized to atomic time**
  - ...for many applications, but UTC offset can also be provided

# Transferring time through a network device

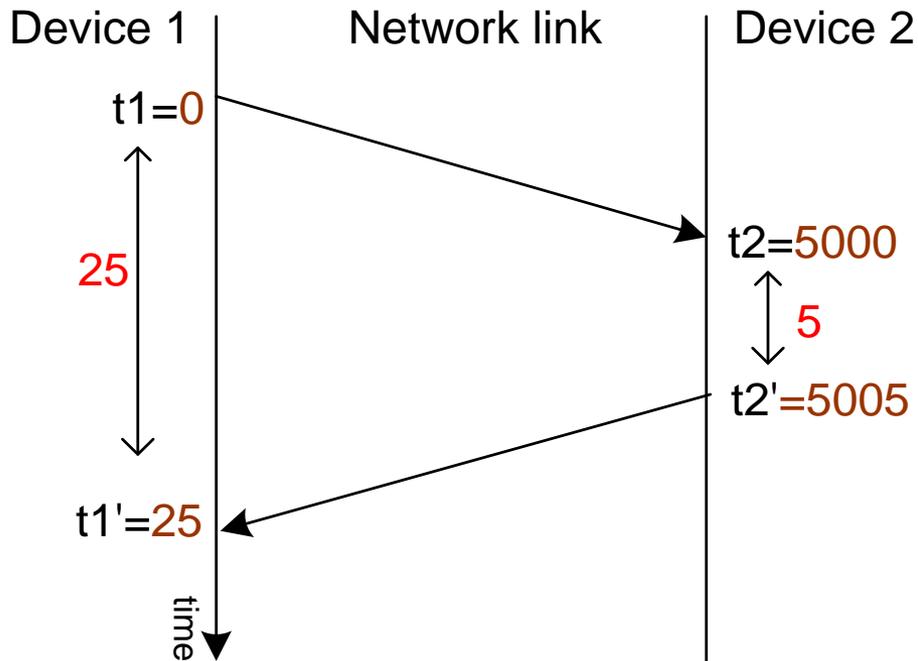


■ is a hypothetical "time measurement" packet

Note: This picture is overly simplistic:  
 $(T_{egr}-T_{in})$   
 can be off by as much as 200PPM from the GM.  
 802.1AS computes "rateRatio" to correct this.

- **Time is sent from master over the link with frames that are LAN/media-specific**
  - Illustrated abstractly here by a Blue frame
- **Accurately measure how the long Blue frame is delayed in processing/queueing**
  - Called "Residence Time"
- **Also compensate for link delay ( $T_{link}$ ) and rateRatio**

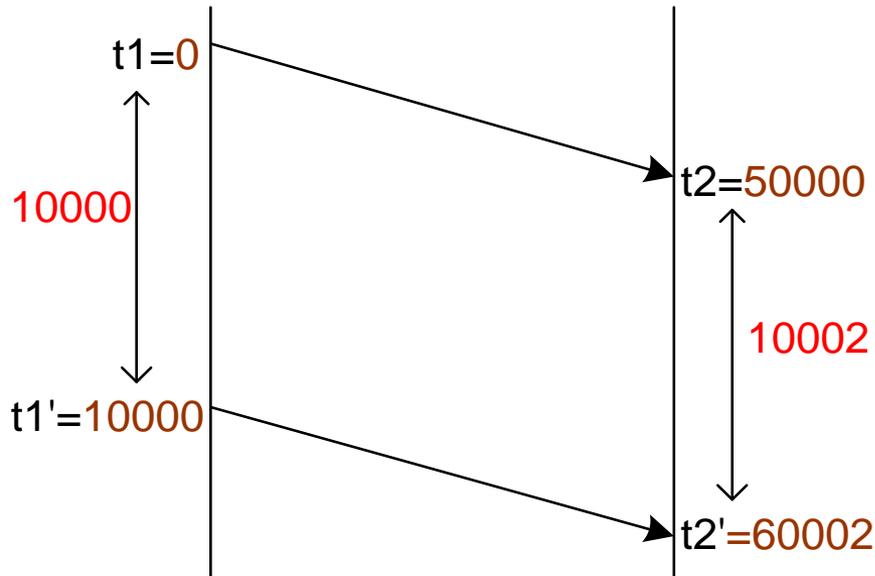
# Example linkDelay measurement



$$\begin{aligned} \text{linkDelay} &= \left( \frac{(t1'-t1) - (t2'-t2)}{2} \right) \\ &= \left( \frac{(25-0) - (5005-5000)}{2} \right) \\ &= 10 \end{aligned}$$

**If link delay is symmetrical, link delay is 10**

# Example neighborRateRatio measurement



$$\begin{aligned}\text{neighborRateRatio} &= \frac{(t_1' - t_1)}{(t_2' - t_2)} \\ &= \frac{(10000 - 0)}{(60002 - 50000)} \\ &\approx 0.998 \\ &\rightarrow \text{neighbor is} \\ &\quad \text{slow by 2000 PPM}\end{aligned}$$

**If link delay is constant, station on left is running 2000 PPM (0.2%) slower than station on the right**

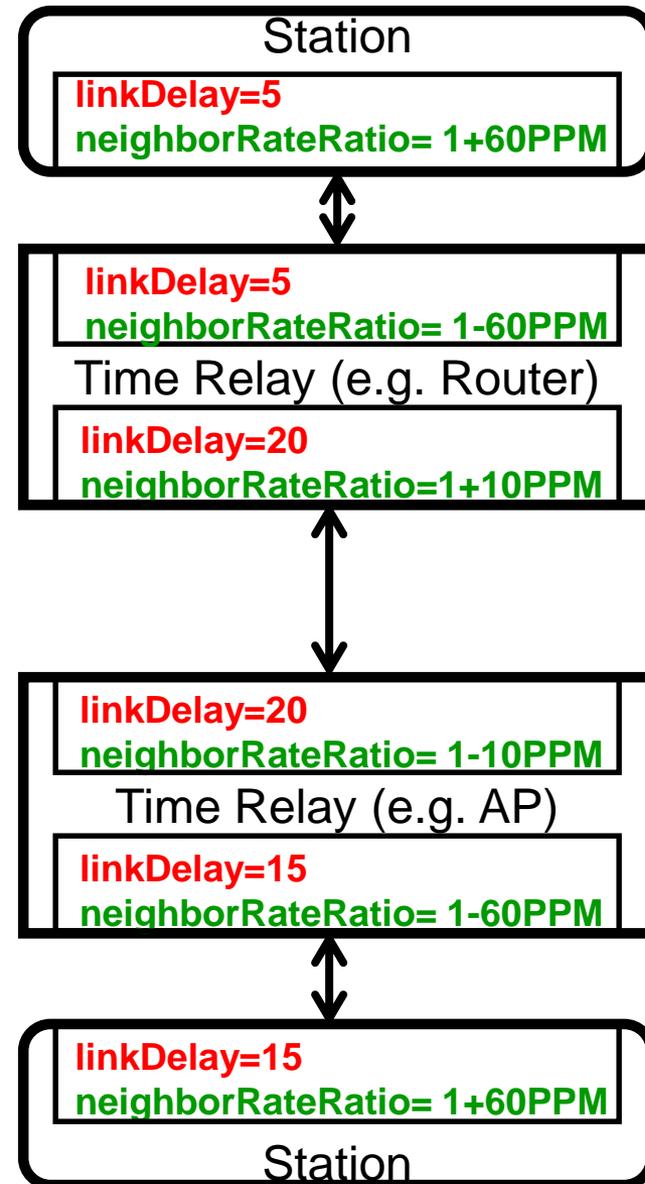
**Note: one *could* measure neighborRateRatio and linkDelay using the same packets (802.11 does, 802.3 does not)**

# Link Delay and Neighbor-Rate-Ratio

## All ports measure

- linkDelay to neighbor
- neighborRateRatio

## Example values are shown



Note: No GrandMaster is needed here (yet)

# End-to-end Rate Ratio

As time is propagated downward,  
neighborRateRatio is  
accumulated at each hop

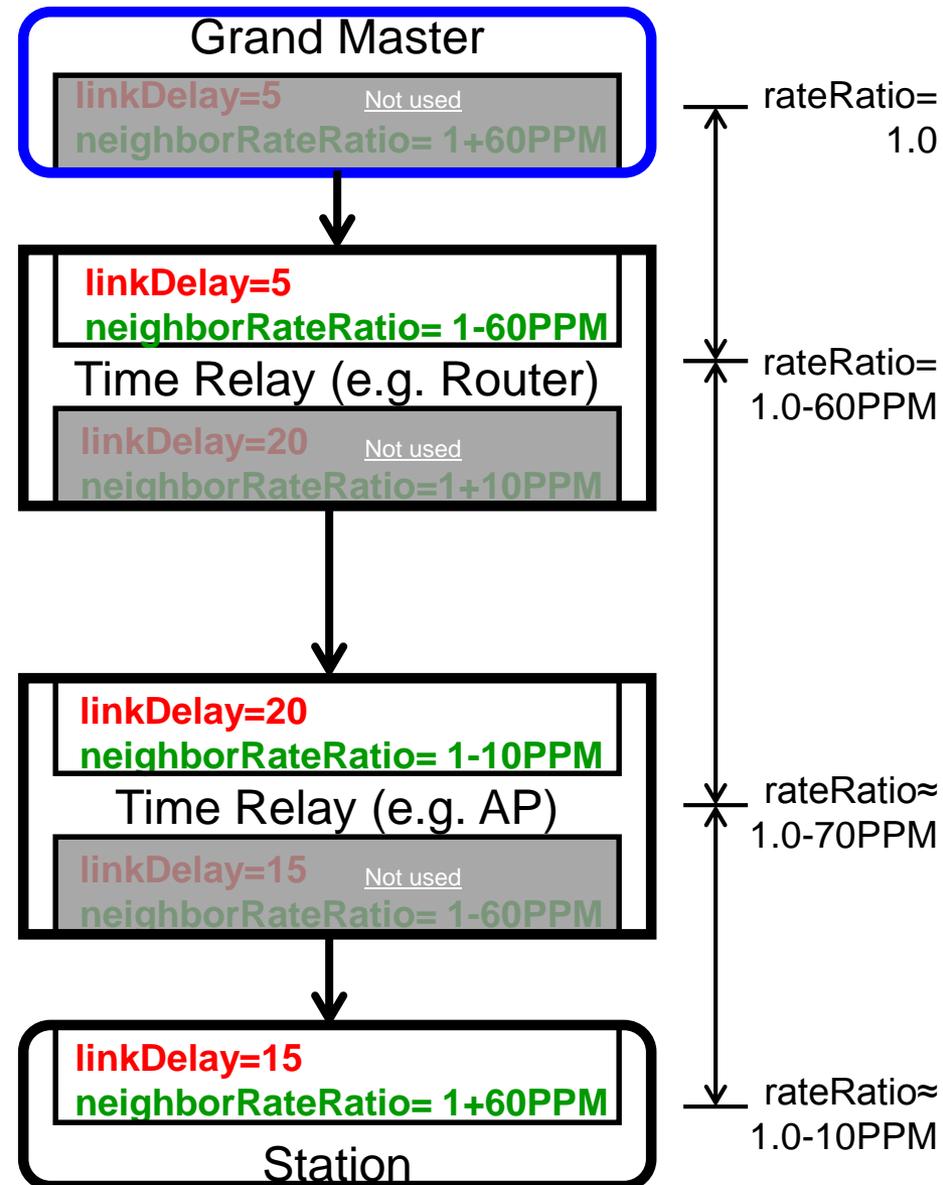
...using the approximation:

```
rateRatio +=  
    (1-neighborRateRatio)
```

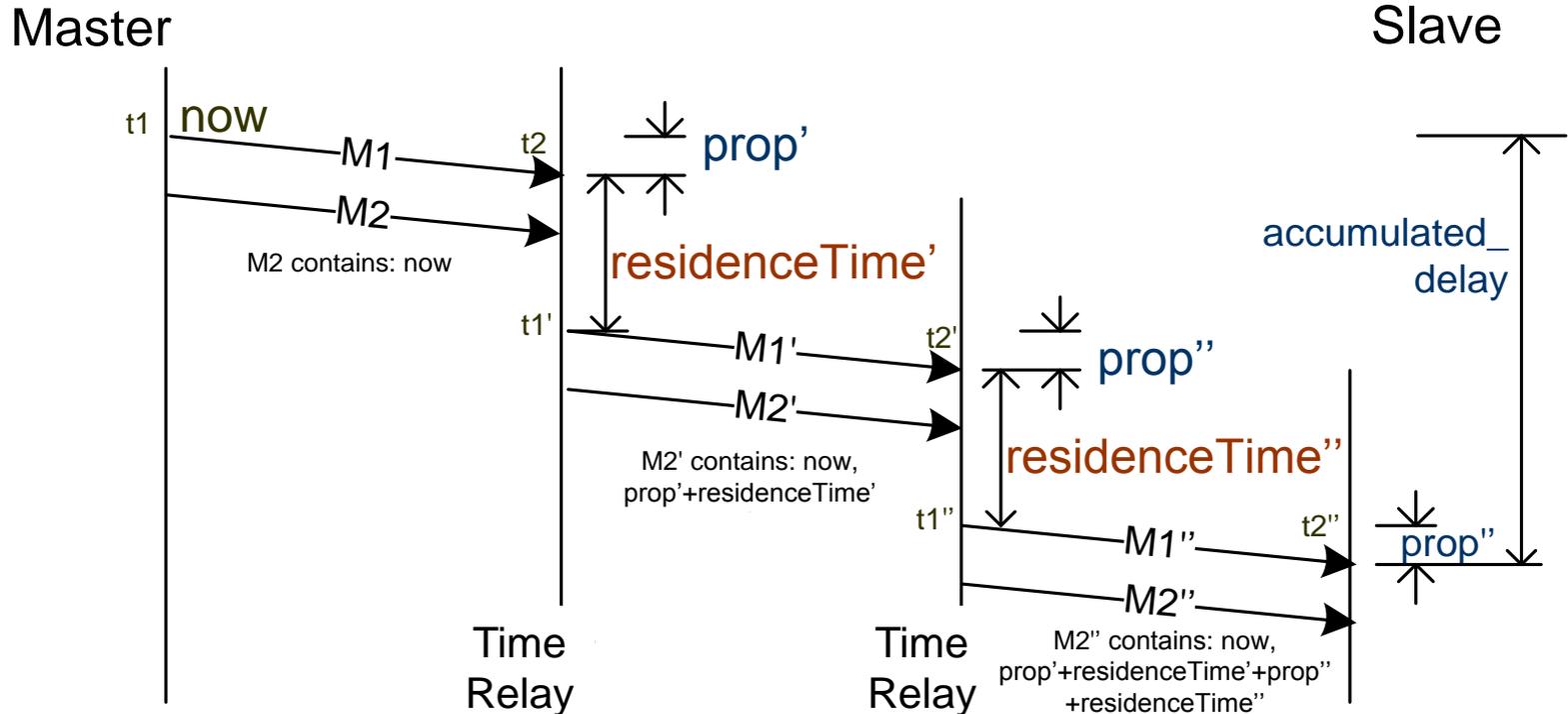
[Initial rate ratio is 1.0 at the GM]

Alternative is to do syntonization per  
relay, but cascaded PLLs are bad

Added benefit: When changing GM,  
endpoints stabilize quickly because  
neighbor parameters are already  
measured



# End-to-end time synchronization across the LAN



**Grand Master initiates timing message, shown as M1, every Sync Interval**  
**At each hop, propagates *now* and measures the actual delay of M1:**

$$\text{delay} = \text{prop} + \text{residenceTime} * \text{rateRatio}$$

**And carries the *accumulated\_delay* in another message, M2**

- Using two messages eliminates real-time processing requirements

**Slaves compute: *currentTime at  $t_2''$*  = *now* plus the *accumulated\_delay***

Note: Message interval on each link may be different

# Measurements summarized

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**Accurate local egress and ingress timestamps captured in the MAC/PHY**

**Timestamps from special time-measurement frame used to compute:**

- **Link delay (per slave port)**
- **Next-neighbor rate ratio (PPM offset to link partner)**

**When the GM sends “now”, the following are accumulated down the tree:**

- **Residence time (per transmitted time measurement frame)**
- **rateRatio (PPM relative to the GM)**

**These measurements described in 802.1AS:**

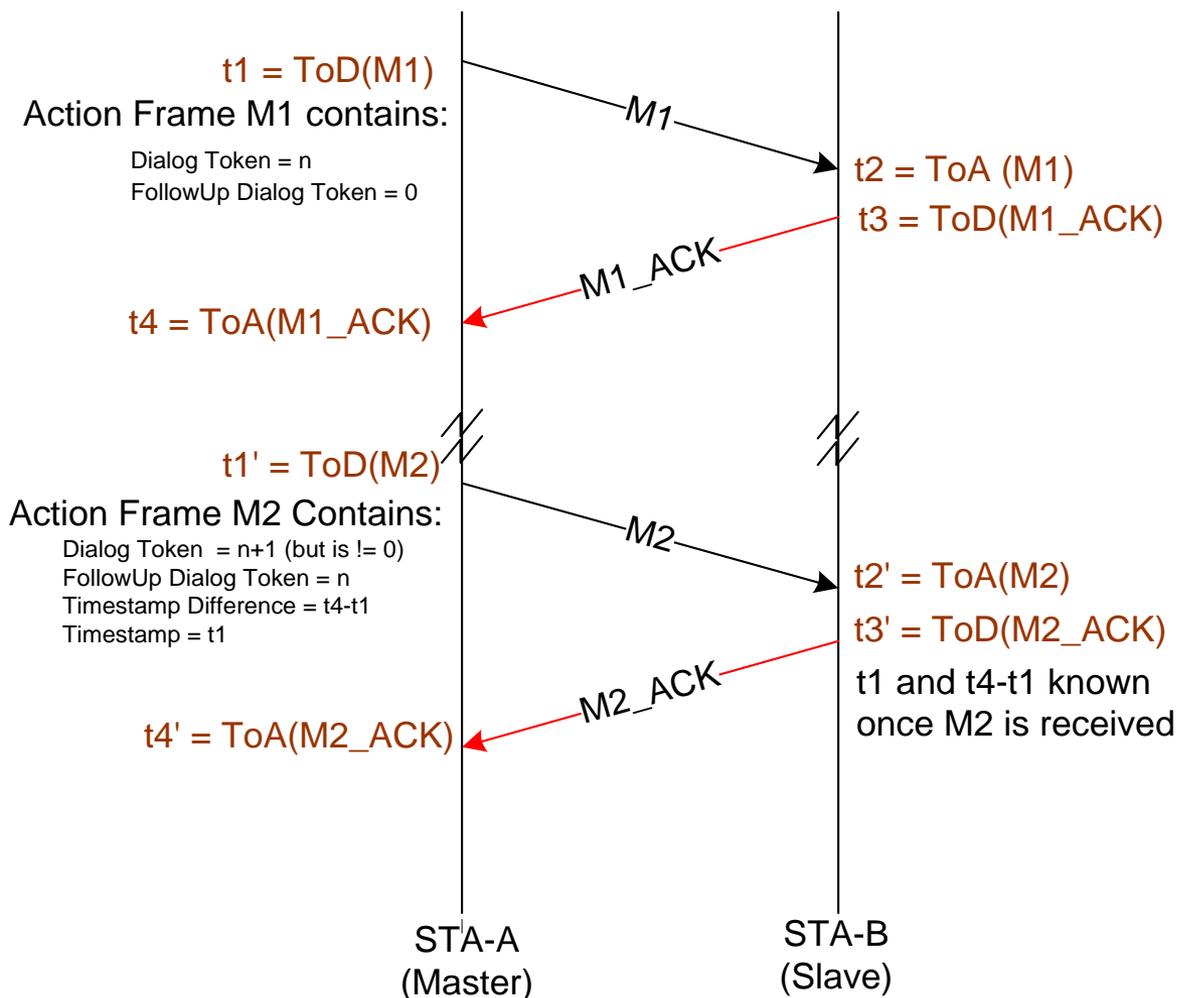
- Clause 11: IEEE 802.3 (Ethernet)
- Clause 12: IEEE 802.11 (Wi-Fi)
- Annex E: MoCA, ITU-T G.hn

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  - Media-dependent
    - 802.11 links
    - 802.3 links
    - Other CSN links
- **Time \*inside\* the system**

# 802.1AS over 802.11 links (Using the 802.11 TIMINGMSMT protocol)



## First exchange:

- takes a measurement

## Subsequent exchange:

- takes a measurement
- also passes timestamps from prior measurement

## Free-running counter used for timestamps

## Allows us to compute:

$$\text{neighborRateRatio} = \frac{(t1' - t1)}{(t2' - t2)}$$

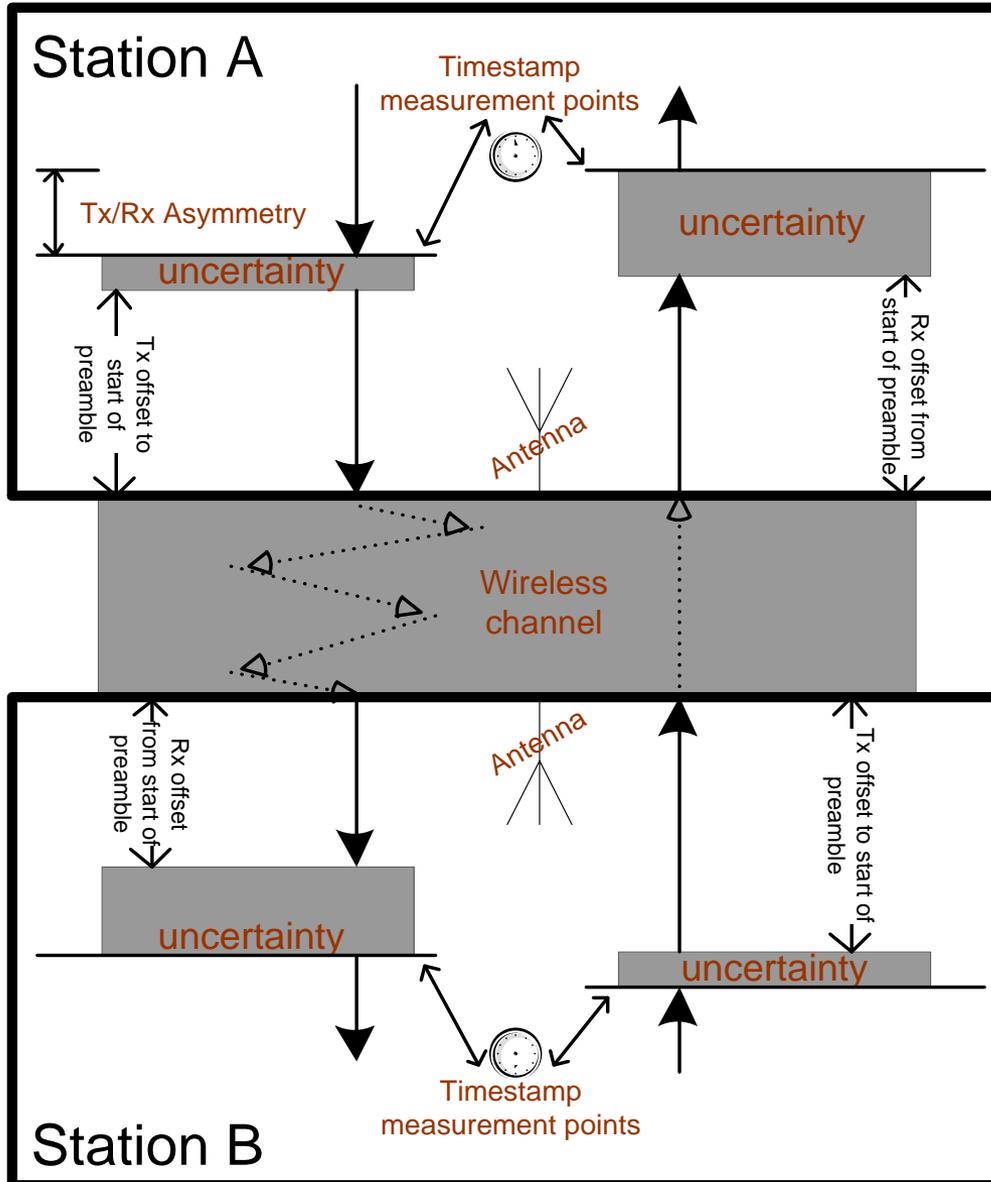
$$\text{linkDelay} = \frac{[(t4 - t1) - (t3 - t2)]}{2}$$

$$\text{timeOffset} = \frac{[(t2 - t1) - (t4 - t3)]}{2}$$

[note: rateRatio is also applied]

NOTE: M1 and M2 have exactly the same format— they're TIMINGMSMT Private Action Frames (and Unicast, BTW)

# Accounting for round-trip 802.11 path asymmetry



The Antenna is the *reference* for Rx and Tx timestamp measurements

Each system knows the delay between the antenna and where the actual Rx and Tx timestamps are captured

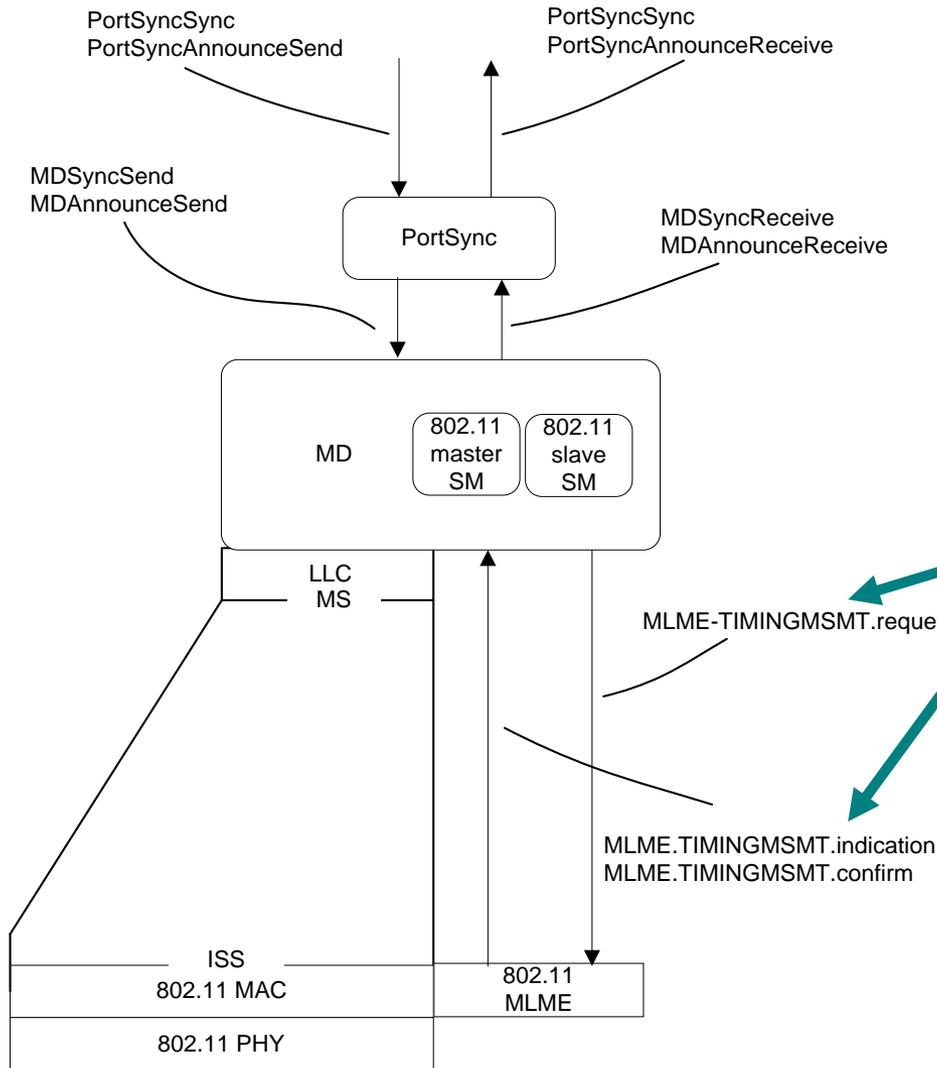
Each timestamp includes some uncertainty

The channel also introduces path asymmetry (and additional uncertainty)

In practice, the difference between the Rx offset and the Tx offset in each system is sufficient information

The rest appears as fixed channel delay and channel uncertainty

# 802.11 links in the 802.1AS architecture



**PortSync is per-port,  
media-independent**  
**MD is per-port media  
dependent**

**802.11-2012 defines these  
MLME.TIMINGMSMT  
primitives**

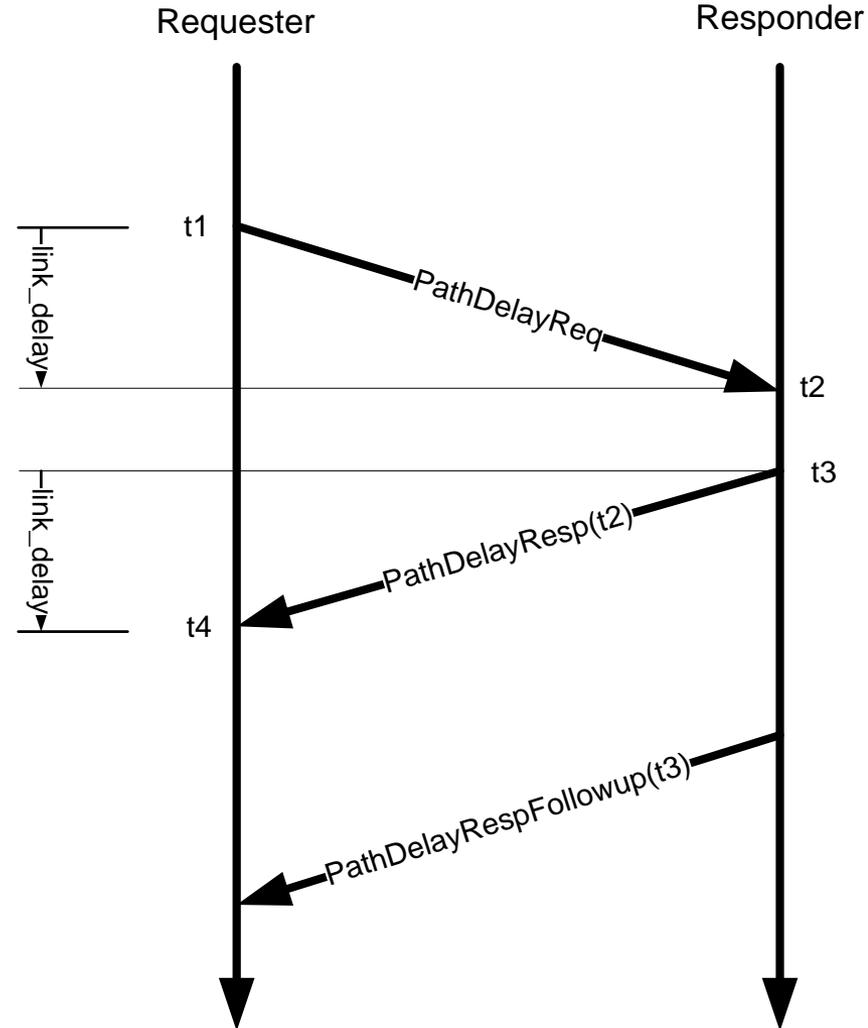
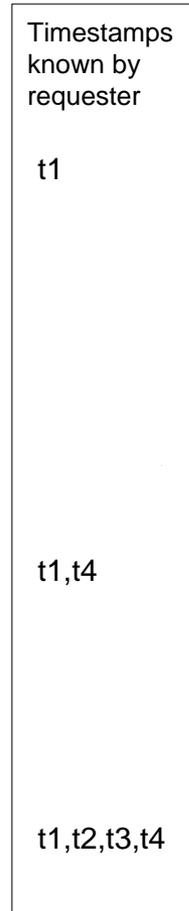
# 802.3 protocol (step 1 of 2)

## Measure link delay:

- 1. Requester schedules PathDelayReq for transmission**  
Using the master's free-running clock
- 2. As it passes out the PHY, t1 is captured**  
Using the master's free-running clock
- 3. Time t2 captured as passes from PHY to MAC**  
Using the slave's free-running clock
- 4. Responder schedules PathDelayResp for transmission, sends t2**  
Using the slave's free-running clock
- 5. Timestamps t3 and t4 captured**  
Using local free-running clocks
- 6. PathDelayRespFollowup carries t3 to requester**

If link delay is fixed & symmetric:  
$$\text{link\_delay} = [(t4-t1) - (t3-t2)] / 2$$

NeighborRateRatio is computed using these and previous timestamps



# 802.3 protocol (step 2 of 2)

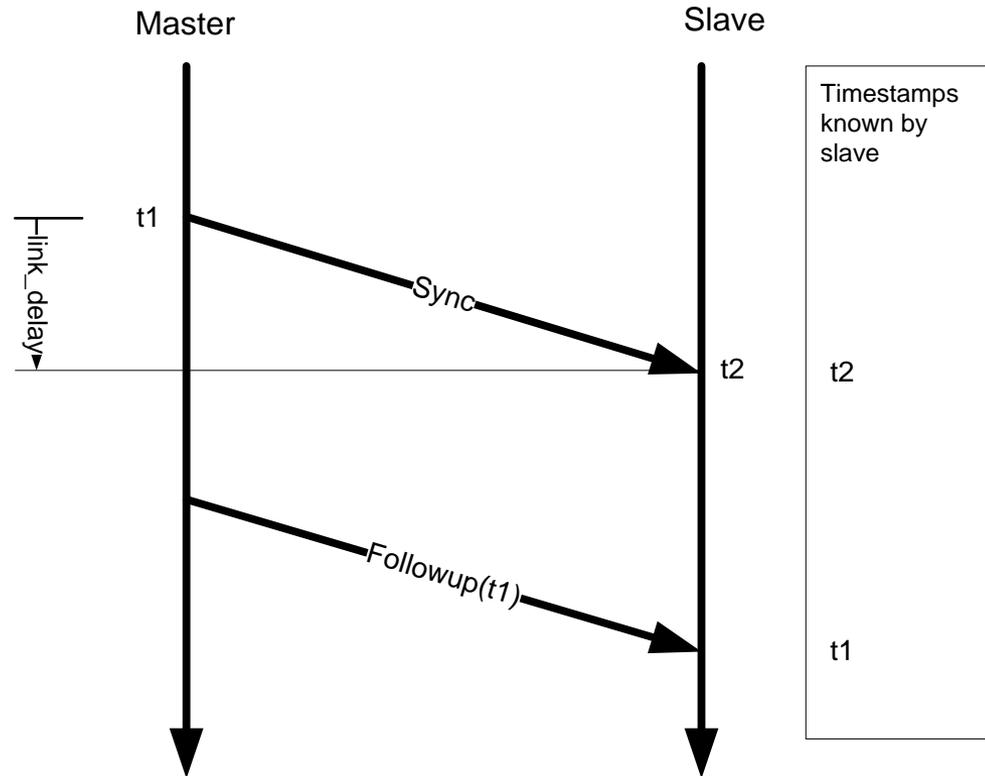
## Synchronize clocks

1. **Master schedules Sync for Tx**
2. **As it passes out the PHY, t1 captured**  
Using master's free-running clock
3. **Time t2 captured when it arrives**  
Using the slave clock
4. **FOLLOWUP carries t1 to slave**

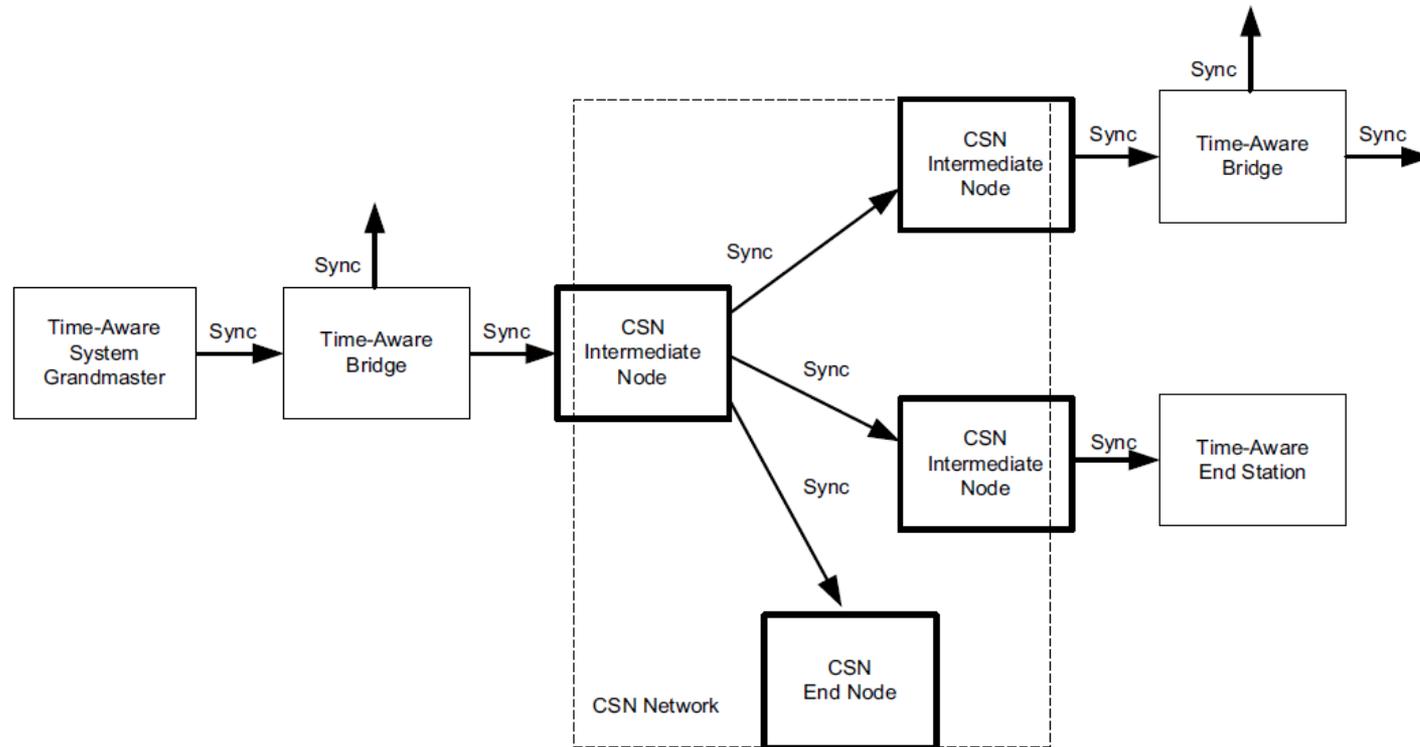
If link delay is fixed & symmetric:

Slave's clock offset  
 $= t2 - t1 - \text{link\_delay}$

**Note: Network infrastructure devices do this too, and communicate the 'residence time' per Sync in the FOLLOWUP frame**



# Coordinated Shared Network (CSN) protocol



- **CSNs use the same protocol as does 802.3 (Ethernet)**
- **Examples of CSNs:**
  - MoCA : Coax networking specification
  - ITU-T G.hn: Coax and power-line networking standard
- **This is described more in Annex E of 802.1AS**

# Message Rates

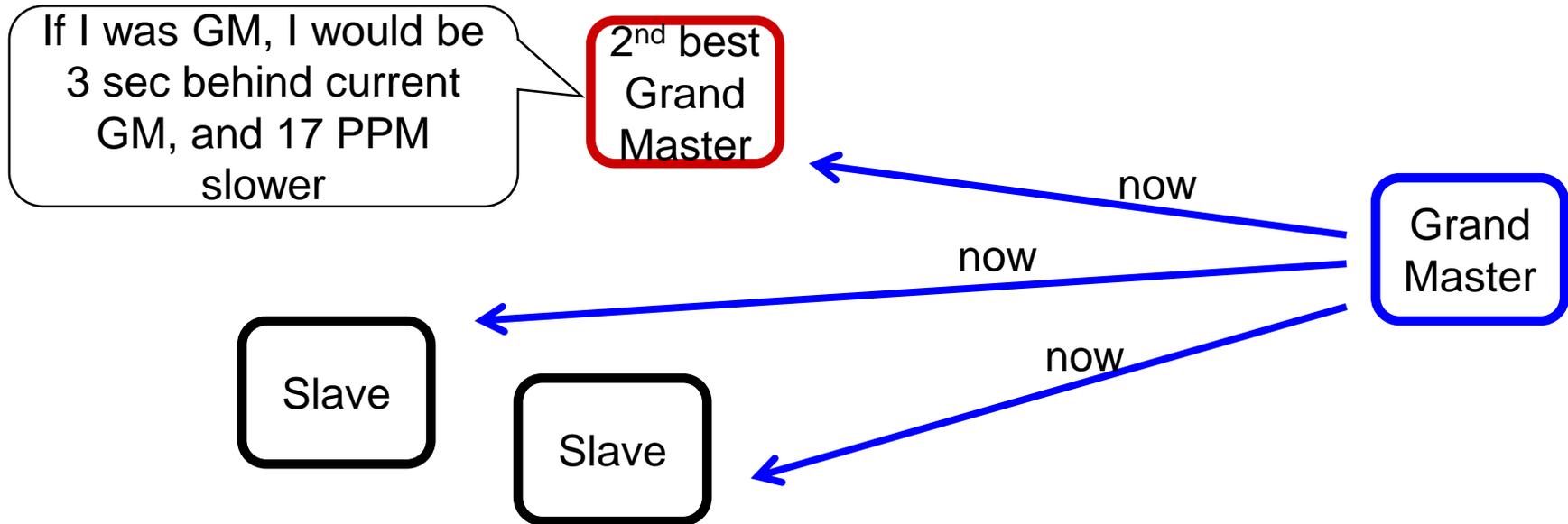
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## Default Message Intervals defined by the standard

- **Announce sent once per second per Port**
- **Timing Information is sent eight times per second**
  - For 802.11: Eight Action Frames per second (encapsulated Followup)
  - For 802.3 : Eight Sync + Eight Followup frames per second
- **Link delay measured once per second**
  - For 802.3: Three PDelay\* frames in each direction per link per second

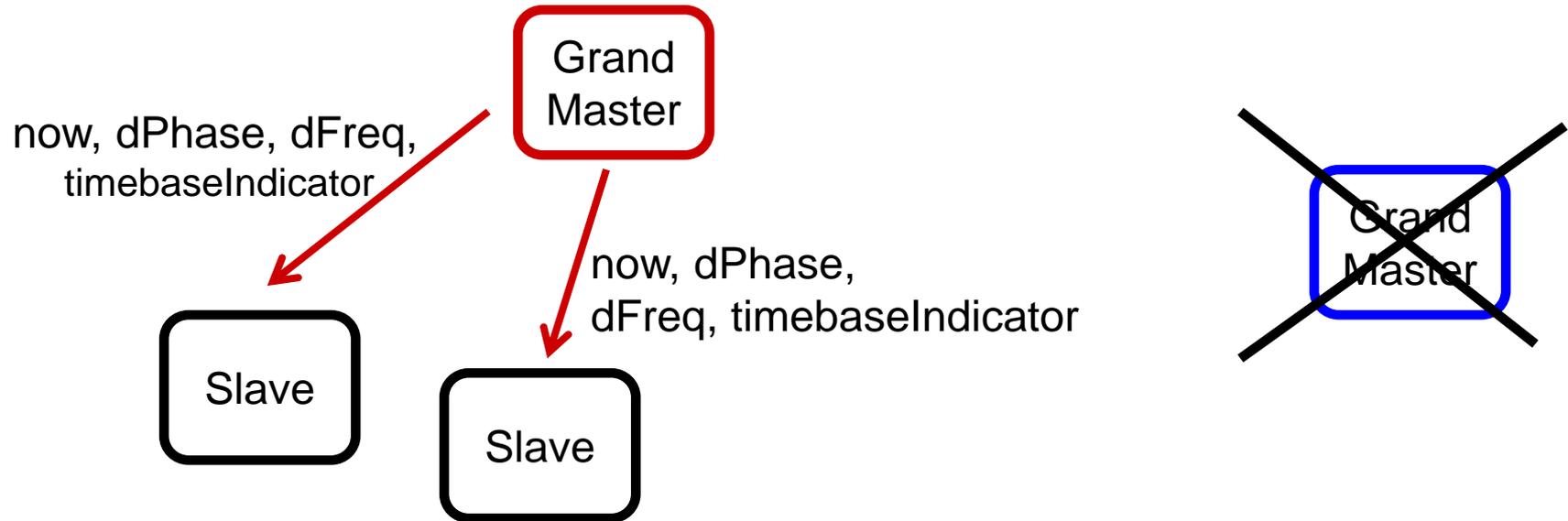
**A different rate can be requested via special SIGNALING messages**

# Handling of discontinuities



- **All potential GMs know their own time offset and intrinsic frequency offset (rateRatio) from the current GM**
- **Time offset may be close to zero (unless it knows UTC-traceable time, e.g., from GPS)**

# Handling of discontinuities – new GM



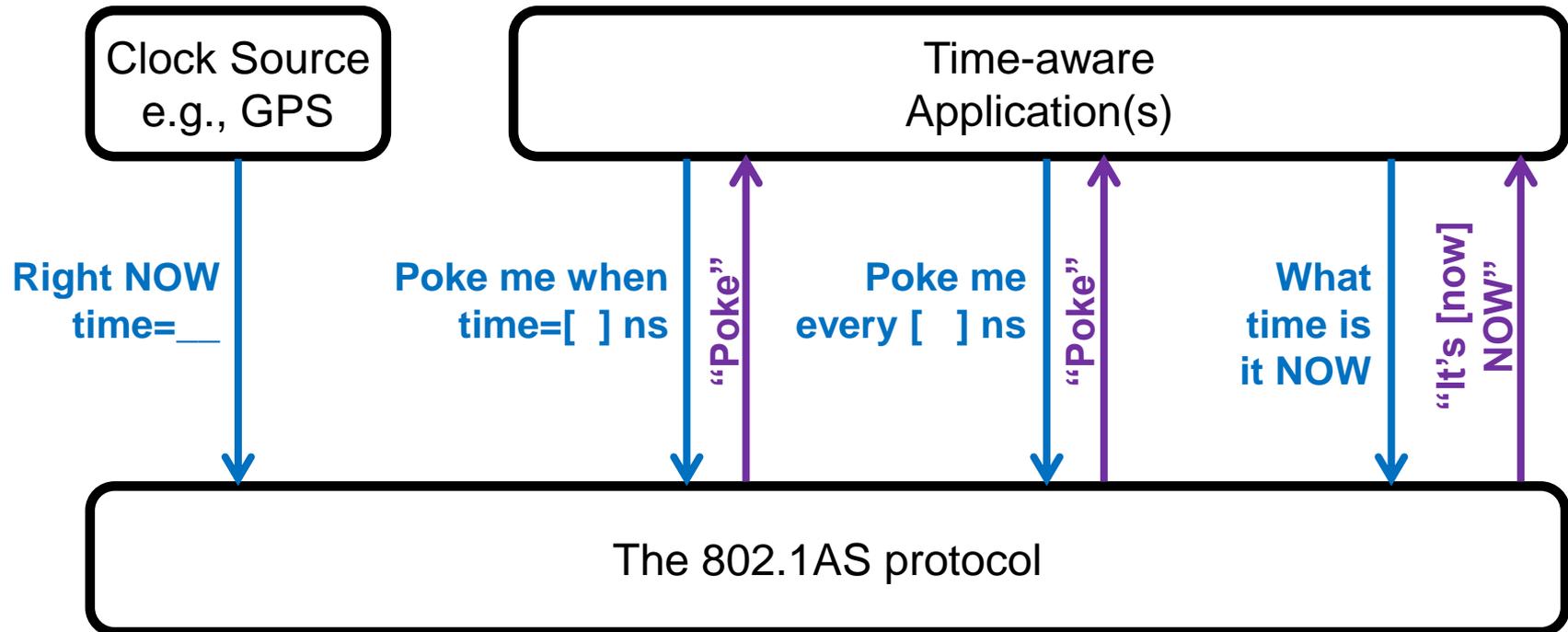
- **The new GM communicates extra info along with the time**
  - TimebaseIndicator (increments after each discontinuity)
  - dPhase: Last phase change (relative to the previous GM)
  - dFreq: Last frequency change (relative to the previous GM)
- **A GM also does this if it experiences a step change**
  - E.g. GPS regains lock or user manually sets the PTP time in the GM

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# Four abstract application interfaces



In 802.1AS these are defined as...

- ClockSourceTime()
- ClockTargetTriggerGenerate()
- ClockTargetClockGenerator()
- ClockTargetEventCapture() ...respectively

# Perspective on real APIs

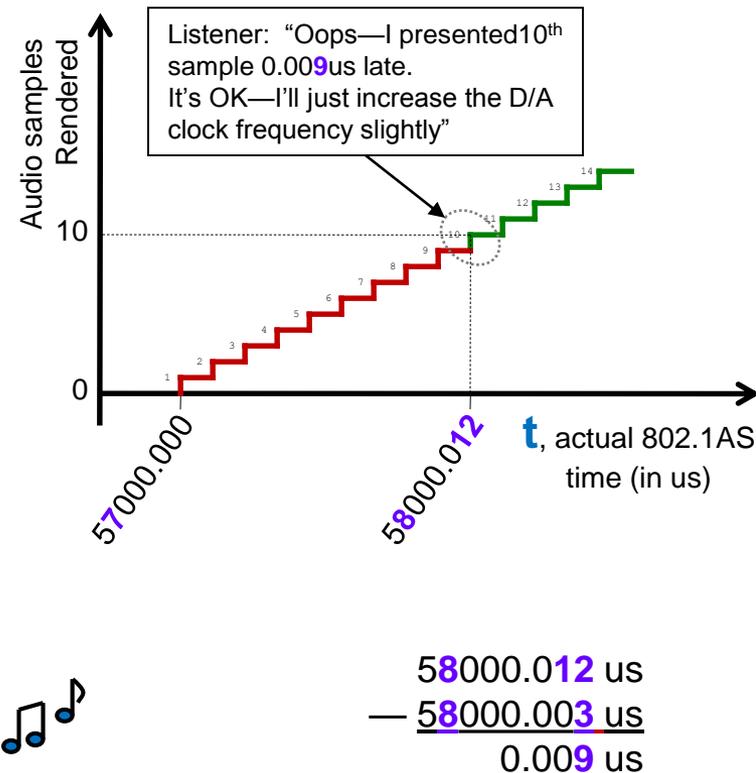
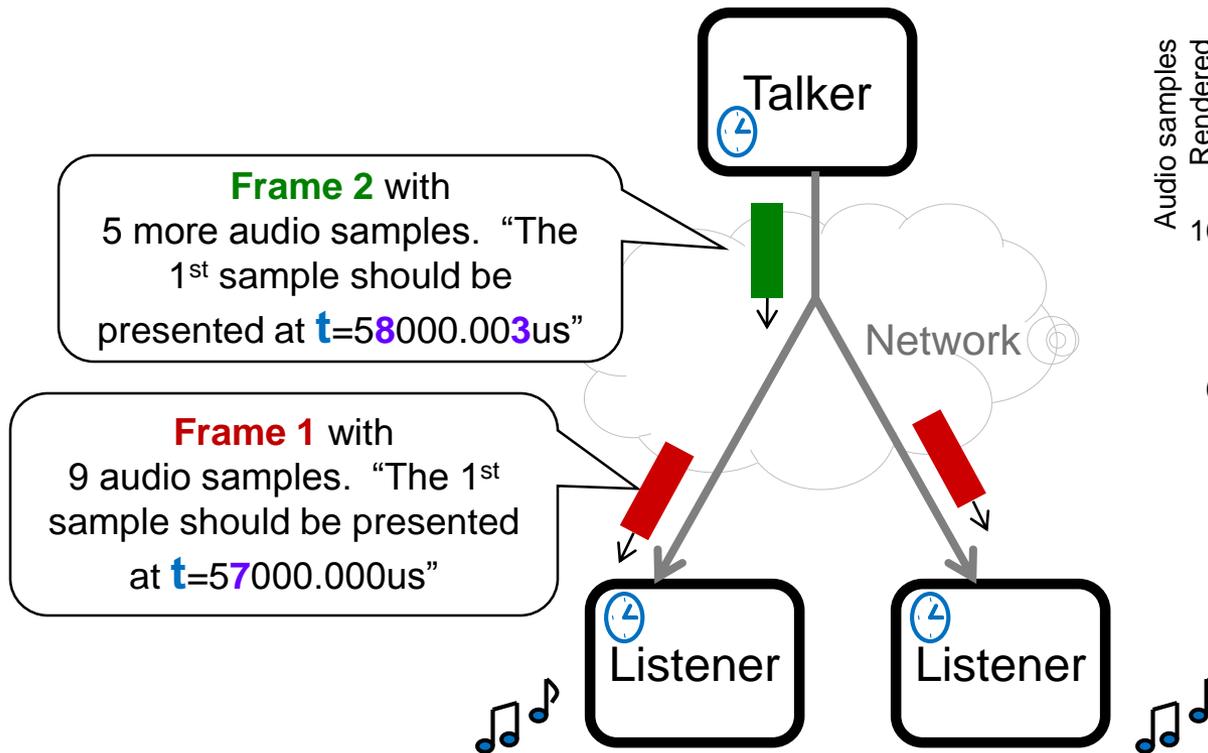
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**“Now”** is not an ideal software semantic for relating network time to System time within a time-aware systems

The preferred fundamental time measurement semantic can be expressed (poetically) as:

*At the instant  
when PTP time was X  
[the other clock] was Y*

# Regenerating the Talker's audio clock



## Notes:

- Time ( $t$ ) is 802.1AS time, known to all talkers & listeners
- Nominal audio sample rate known beforehand, e.g. 10KHz, 100us/sample above
- 802.1AS time is "seconds since Jan 1, 1970 TAI" with precision of 1 nanosecond
- This allows an ARBITRARY number of independent media clocks simultaneously
- The Grand Master need not be the Talker

# Relating 802.1AS time to media clock

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- **“Presentation Time” added to header of media packets**
  - The Presentation Time is an 802.1AS time
  - Multiple methods for carrying Presentation Time
    - Layer-2: IEEE Std 1722™-2011 (based on IEC-61883)
    - Layer-3/IP: See IETF draft avtcore-clksrc (RTP Extension)
    - Layer-3/TCP/HTTP: TBD
- **Media is buffered in the renderer until the Presentation Time**
  - In some cases, thanks to M\*RP and FQTSS (802.1Qat/Qav), Listener buffer size can be minimized
    - But this is out of scope of 802.1AS

# Future work

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- **As of July 2014:**

- An amendment (soon, a revision) of 802.1AS, is under development to add “Enhancements and performance improvements”. See

<http://www.ieee802.org/1/pages/802.1ASbt.html>

- More detailed status as of June 2014 is here:

<http://www.ieee802.org/1/files/public/docs2014/as-mjt-update-for-WSTS-0614-v01.pdf>

# References

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802 Stds Available Free of Charge  
<http://standards.ieee.org/about/get/>

- **IEEE Std. 1588™-2008**
  - The Definitive Precision Time Protocol (PTP) Definition
- **IEEE Std. 802.1AS™-2011**
  - Clause 12: Media-dependent layer specification for IEEE 802.11 links
  - Also, Technical and Editorial Corrections are published in
    - 802.1AS-2011 Corrigendum 1-2013
- **IEEE Std. 802.11™-2012**
  - Clause 6.3.57: Timing Measurement
  - Clause 10.23.5: Timing Measurement Procedure
- **API definition for use of 802.1AS (and related standards)**
  - [http://www.avnu.org/files/static\\_page\\_files/C5E0B5F8-1D09-3519-ADB32F1F88E6C057/AVnu\\_SWAPIs\\_v1.0.pdf](http://www.avnu.org/files/static_page_files/C5E0B5F8-1D09-3519-ADB32F1F88E6C057/AVnu_SWAPIs_v1.0.pdf)