

Aggregation of Micro-Streams into one Common Stream

Feng Chen, Franz-Josef Goetz, Marcel Kiessling, Jürgen Schmitt

IEEE 802.1 TSN Call
Dec. 10, 2018

Control Loops within Industrial Machines

A typical industrial machine is build up of a huge amount of different physical actuators and sensors. They are connected to so called IO-Devices by different technologies, from electric wire to small busses (e.g. IO Link, PROFIBUS, CAN...). The IO-Devices in turn send/receive data in **a wide range of different rates** (e.g. typically between 1 kHz and 60 kHz) and **in a wide range of amount of real time data** according to the requirements of the control application and the type of sensors and actuators connected.

Industrial use cases:

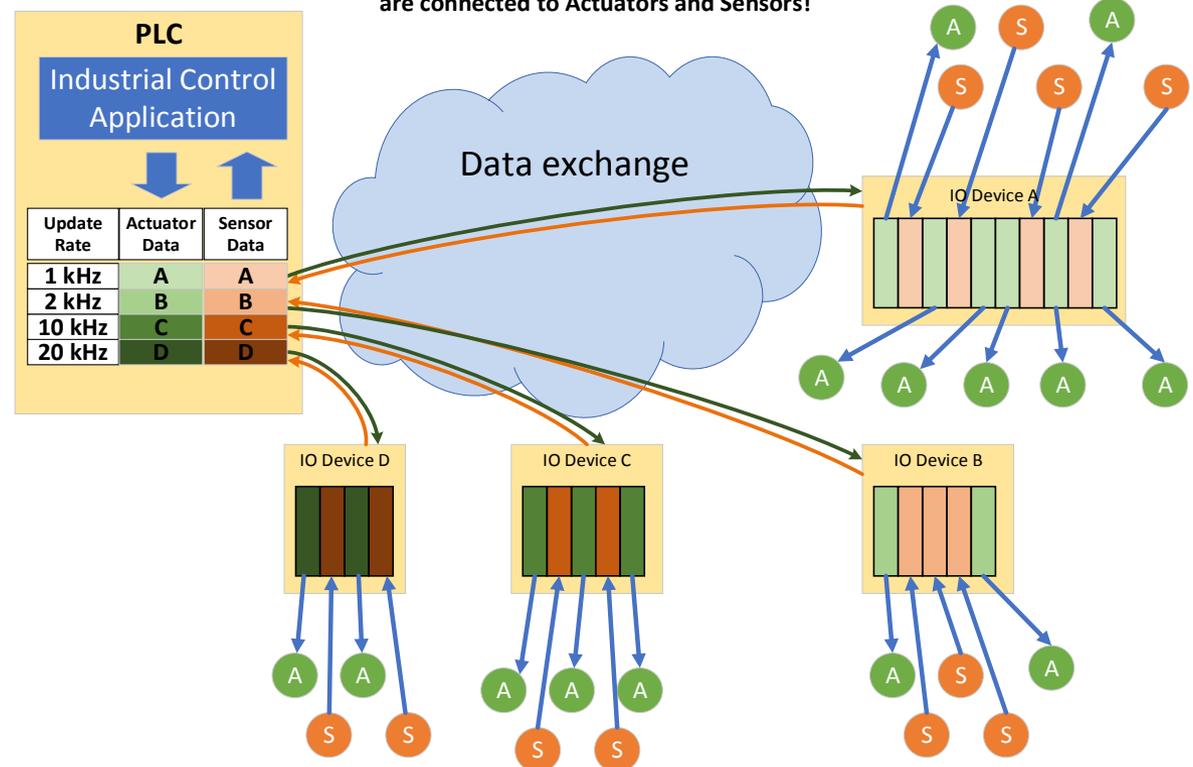
Up to several hundred IO-Devices that are connected to thousands of actuators and sensors periodically exchange their real time data with one or several PLCs with

- **low data rates and small amount of real time data,**
 - **low data rates and huge amount of real time data,**
 - **high data rates and small amount of real time data,**
 - **high data rates and huge amount of real time data,**
- or with a mixture of all of them.

In contrast, Audio / Video applications typically have high data rates with huge amount of data.

A typical Industrial Automation Use Case for Connectivity

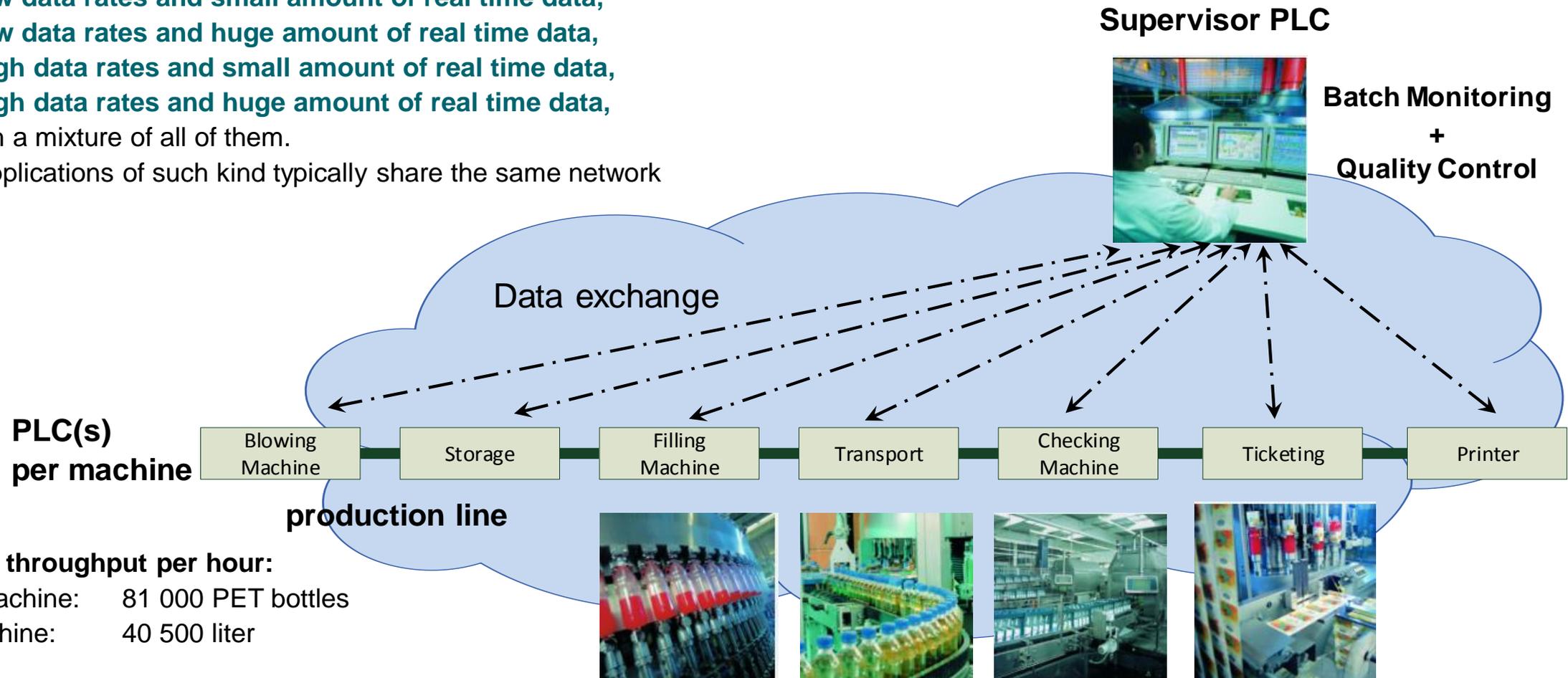
Programmable Logic Control (PLC) exchange periodically real time data with Input/Output(IO)-Devices which are connected to Actuators and Sensors!



Supervisor PLC <-> PLCs Use Case: Quality Control at Real Time in a Bottling Plant

Industrial use cases:

- ❑ Up to tens of PLCs periodically exchange real-time data with one supervisor PLC with
 - low data rates and small amount of real time data,
 - low data rates and huge amount of real time data,
 - high data rates and small amount of real time data,
 - high data rates and huge amount of real time data,or with a mixture of all of them.
- ❑ Many applications of such kind typically share the same network



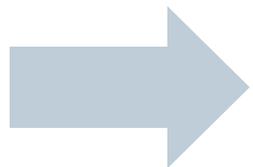
Example for throughput per hour:

- Blowing Machine: 81 000 PET bottles
- Filling Machine: 40 500 liter

Example for Stream Reservation based on MSRP without Micro-Stream (μ Stream) aggregation

Example:

- 1 PLC \leftrightarrow 50 IO-Devices (bidirectional, \sim 50 μ Streams per direction)
 - IO-Device with real time data rate of 1 kHz (μ Stream transmission rate)
 - Max e2e latency: 1ms
 - Max hop count: 16
 - ➔ Max per hop latency: 62,5 μ s
- Stream reservation based on MSRP
 - a SR-Class with class measurement interval 62,5 μ s \sim 16 kHz to fulfill the max e2e latency requirement



If making each μ Stream an individual Stream using the existing MSRP mechanisms without μ Stream-Aggregation, it will result in an overprovisioning factor of **16** for reservation of such 50 μ Streams per direction.

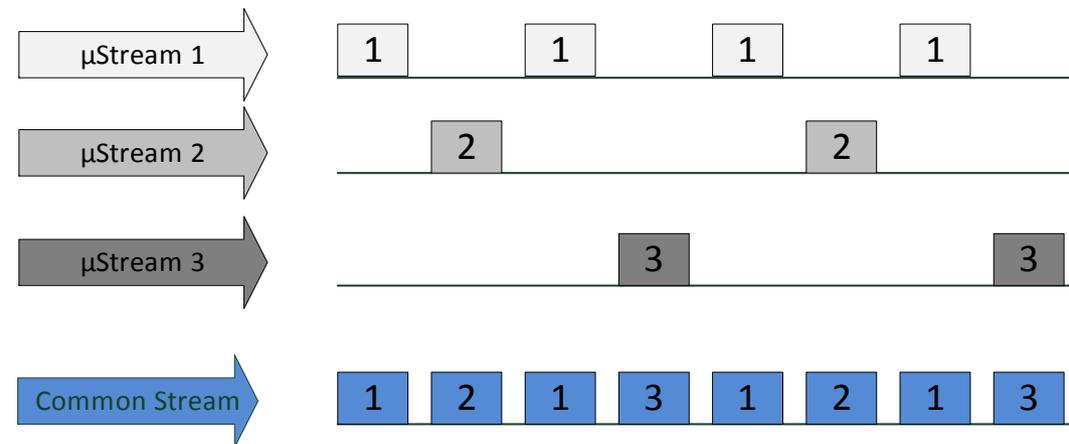
Proposal: μ Stream Aggregation using Interleaving

Definition

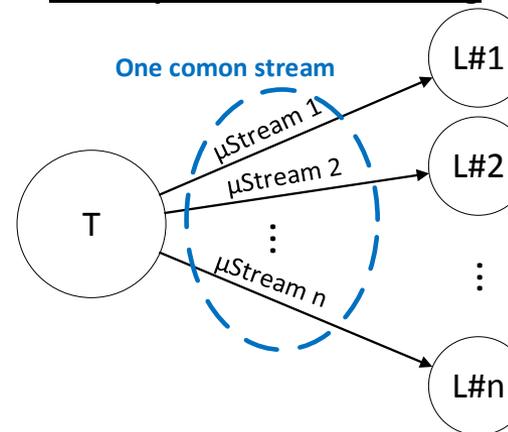
- **μ Stream Aggregation** combines multiple μ Streams into one common Stream
 - forwarded in the network using one multicast destination MAC address
 - reserved in one stream reservation process by using a common Traffic Specification (TSpec)

- **μ Stream Interleaving is a method to combine multiple μ Streams into one common Stream using multiplexing.**
 - **Talker μ Stream-Interleaving** (for 1-to-n communication)
A single Talker is responsible for organizing its μ Streams for multiple Listeners into a common Stream.
 - **Listener μ Stream-Interleaving** (for n-to-1 communication)
A single Listener is responsible for organizing the μ Streams that are transmitted by multiple talkers to this Listener into a common Stream.

μ Stream Aggregation using Interleaving

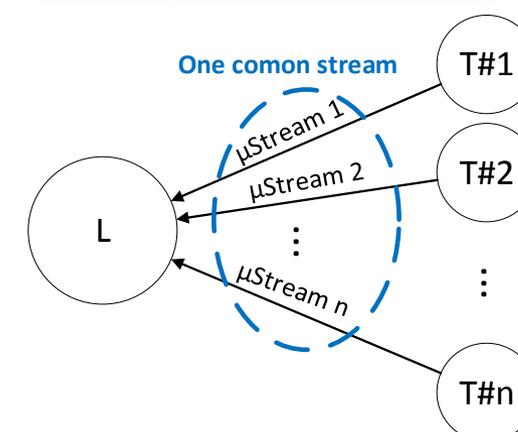


Talker μ Stream-Interleaving



One Talker \longrightarrow n Listeners

Listener μ Stream-Interleaving



One Listener \longleftarrow n Talkers

Local Computation of Interleaved Schedule for μ Stream Aggregation

An **Interleaved Schedule (IS)** for a common stream

- is locally computed either at the Talker or at the Listener, who knows the traffic specifications of all aggregated μ Streams, but not required to have the knowledge of the network topology.
- specifies a repeating time schedule that allocates the time slots to transmission or reception of all aggregated μ Streams in a certain way, e.g. distribute total bandwidth of all μ Streams among slots as evenly as possible.

□ For **Talker μ Stream-Interleaving** in 1-to-n communication, an **IS**

- is computed locally at the **Talker**, who intends to transmit n μ Streams to multiple Listeners.
- schedules interleaved μ Stream transmission at the Talker**
- is used by the Talk, not necessarily known to the Listeners

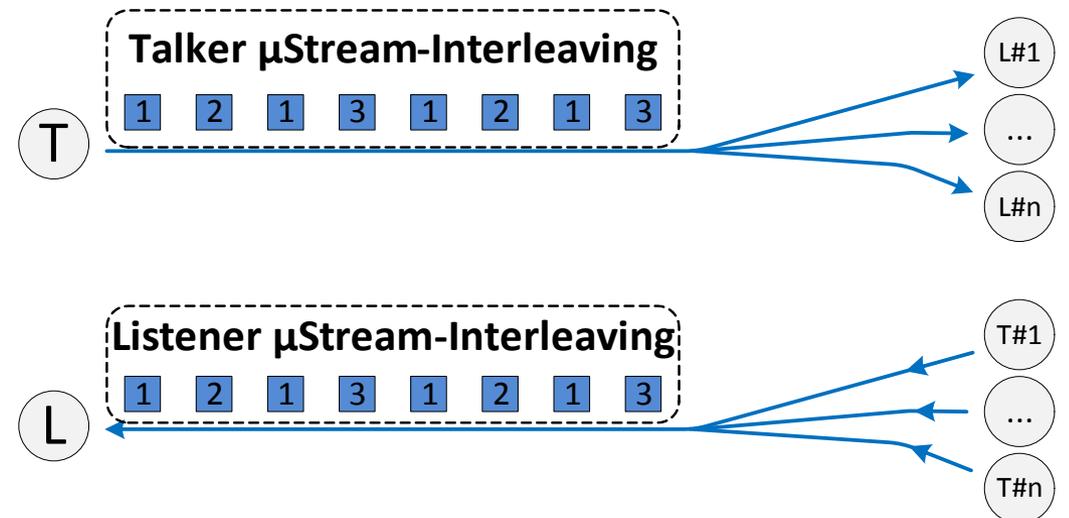
□ For **Listener μ Stream-Interleaving** in n-to-1 communication, an **IS**

- is computed locally at the **Listener**, who intends to receive n μ Streams from multiple Talkers.
- schedules interleaved μ Stream reception at the Listener**
- must be propagated to and converted for use by each Talker**

Example of an Interleaved Schedule

Traffic specification of μ Stream i given by application: $TS_i(M_i, N_i, L_i)$
 M_i : max frame size; N_i : number of frames; L_i : trans. interval

	Slot 1	Slot 2	Slot 3	Slot 4	Slot 5	Slot 6	...	Slot m
μ Stream 1	$M_1 \times N_1$	-	$M_1 \times N_1$	-	$M_1 \times N_1$	-	-	$M_1 \times N_1$
μ Stream 2	-	$M_2 \times N_2$	-	-	-	$M_2 \times N_2$	-	-
...
μ Stream n	-	-	-	$M_n \times N_n$	-	-	-	$M_n \times N_n$



Stream Reservation for Talker μ Stream-Interleaving

Target application of Talker μ Stream-Interleaving:

- one talker to n listeners communication, e.g. PLC to I/O devices, supervisor PLC to PLCs

Assumption:

- The talker has the information of all the μ Streams to be aggregated, such as μ Stream TSpecs, application-level μ Stream identification, etc.

Workflow:

- The Talker computes the IS (see previous slide) and derives the TSpec for use in the reservation of the common stream (see figure right above)
- The Talker initiates the reservation process using TSpec of the common stream, which follows the conventional reservation procedures with „Talker-Advertise / Listener-Join“

=> Talker μ Stream-Interleaving is transparent to the network and can be applied with the existing reservation method.

- Upon successful reservation, the Talker starts stream transmission according to the locally computed IS and using the same stream DA for all aggregated μ Streams.
- Each Listener performs local filtering of received μ Streams.

$$N_{s1} = \sum_{i=1}^n N_i \text{ in slot 1}$$

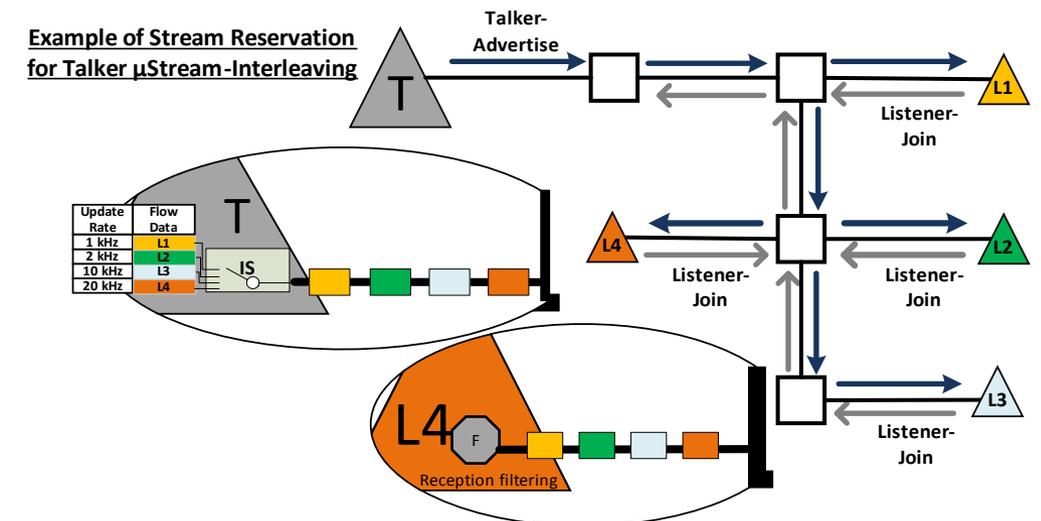
	Slot 1	Slot 2	Slot 3	Slot 4	Slot 5	Slot 6	...	Slot m
μ Stream 1	$M_1 \times N_1$	-	$M_1 \times N_1$	-	$M_1 \times N_1$	-	-	$M_1 \times N_1$
μ Stream 2	-	$M_2 \times N_2$	-	-	-	$M_2 \times N_2$	-	-
...
μ Stream n	-	-	-	$M_n \times N_n$	-	-	-	$M_n \times N_n$

Derive TSpec of the common stream: $TS(M, N, L)$ from the TSpecs of n aggregated μ Streams $TS_i(M_i, N_i, L_i)$ and the IS

$$M = \max\{M_1, M_2, M_3, \dots, M_n\} \rightarrow \text{max. frame size}$$

$$N = \max\{N_{s1}, N_{s2}, N_{s3}, \dots, N_{sm}\} \rightarrow \text{max. number of frames}$$

$$L = \text{slot length}$$



Stream Reservation for Listener μ Stream-Interleaving

Target application of Listener μ Stream-Interleaving:

- n talkers to one listener communication, e.g. I/O devices to PLC, PLCs to supervisor PLC

Assumption:

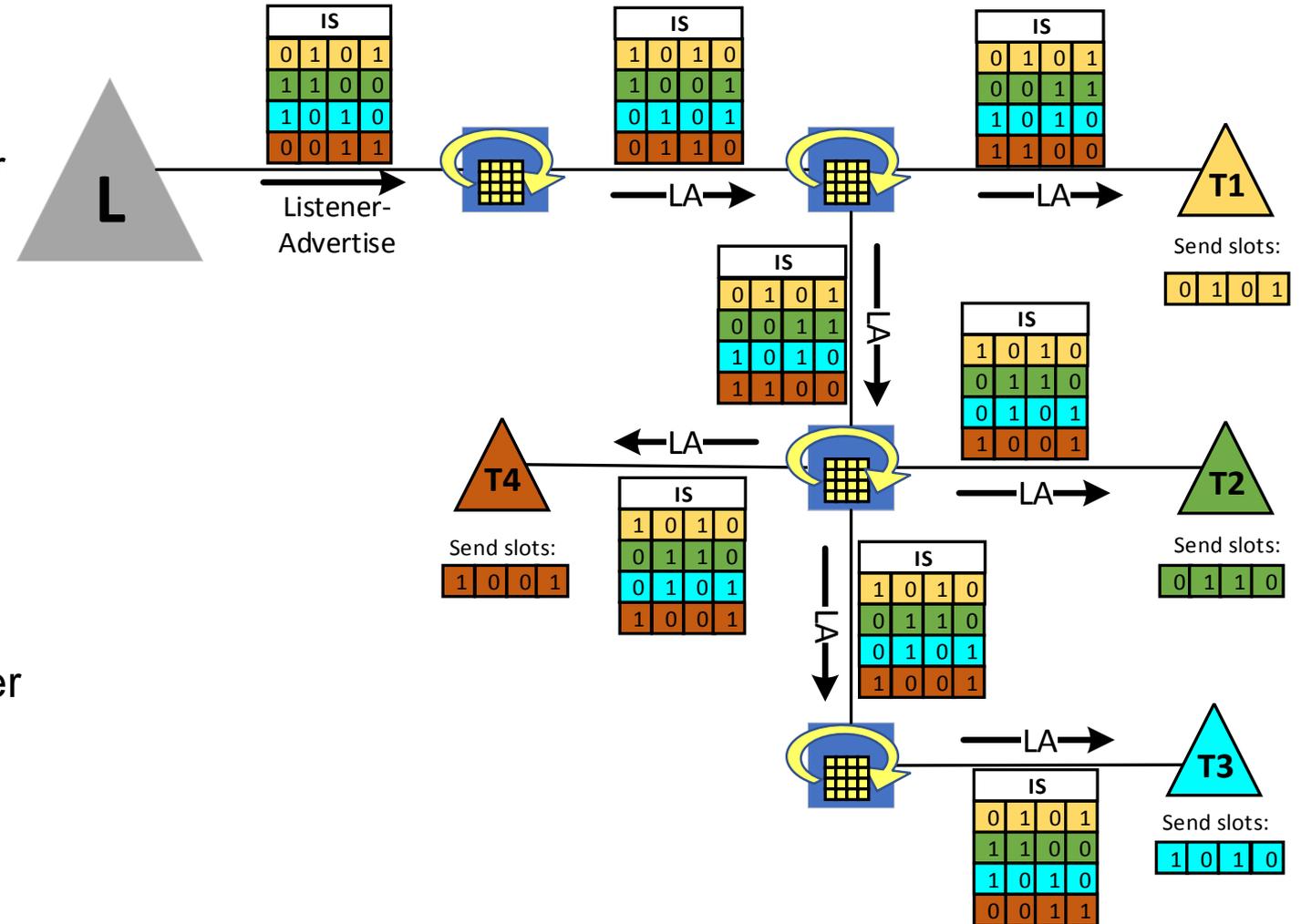
- The Listener has the information of all the μ Streams to be aggregated, such as μ Stream TSpecs, application-level μ Stream identification.
- The **Cyclic Queuing and Forwarding** (CQF) is applied in the network for transmission of the aggregated μ Streams, while the CQF cycle time is made equal to the time slot length of the IS

Workflow:

- The Listener computes the IS and derives the TSpec for use in the reservation of the common stream (*see previous slide*)
- The Listener initiates the reservation process, which follows a reverse reservation procedure (in contrast to the conventional one), called **„Listener-Advertise (LA) / Talker-Join (TJ)“**
 - The LA carries the TSpec of the common stream and other stream information as in the legacy Talker-Advertise.
 - In addition, the LA carries the IS computed by the Listener together with the timing information necessary for the Talkers to derive the beginning of each scheduling cycle.
 - The IS is adjusted at each hop during its propagation along each path from the Listener to each Talker, which rotates the IS once in that each row in the IS is shifted back (to earlier time) by exactly one time slot. (This is feasible due to the use of CQF)
- Each Talker transmits its μ Stream according to the corresponding row in the received IS.

Example of Interleaved Schedule Rotation in Reservation for Listener μ Stream-Interleaving

- ❑ An essential step in Listener μ Stream-interleaving is to convert a receiving schedule locally computed by the Listener to a sending schedule for use by each Talker within a distributed stream reservation process.
- ❑ Using CQF on the data plane and adding support for a new reservation μ Stream with “Listener-Advertise / Talker-Join” to the reservation protocol are two keys to applying μ Stream aggregation with listener μ Stream-interleaving for n-to-1 communication.



□ Advantages of μ Stream-Aggregation to a common Stream (focus of this presentation)

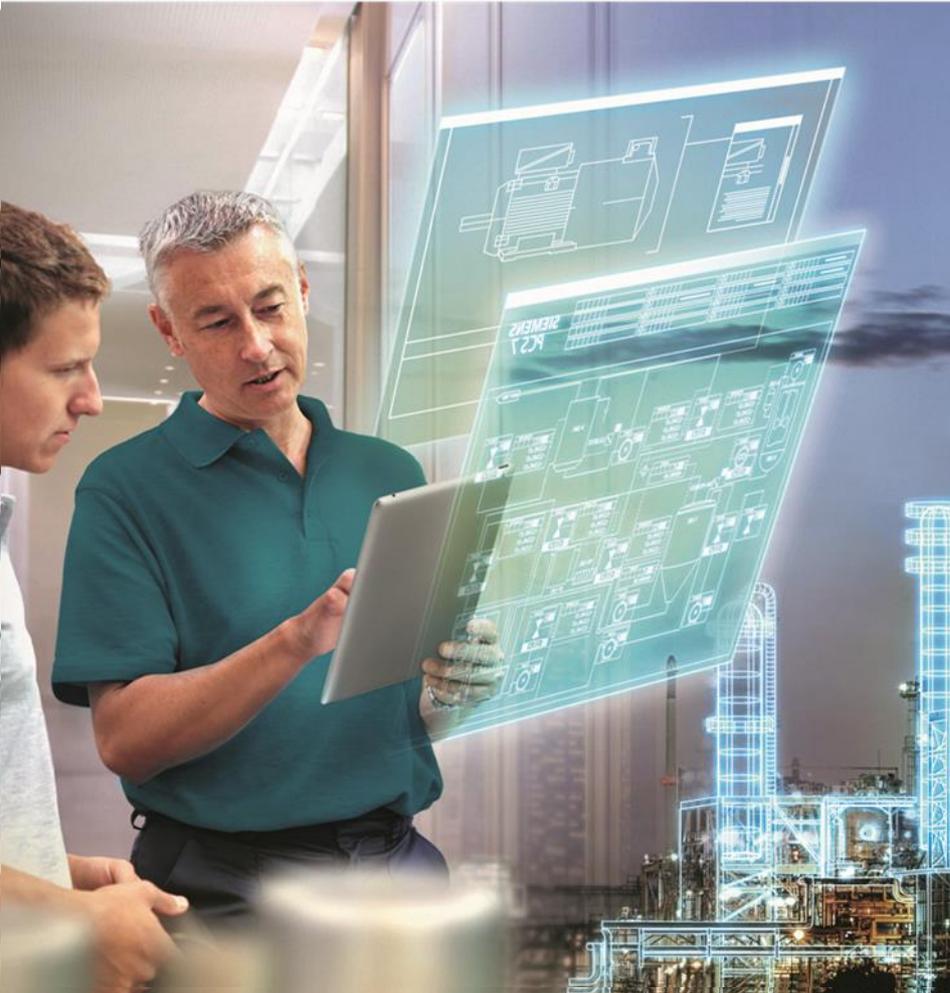
- Reduce bandwidth overprovisioning
- Reduce the control plane overhead, e.g. the number of reservation data and Stream Das
- Talker / Listener μ Stream-Interleaving is independent from topology
- Locally computable by Talker / Listener, no need for a central controller
- Streams / common Streams remain independent from each other
- Dynamic given by a reservation protocol is still available to Streams / common Streams

□ Outlook for aggregated Streams (upcoming)

- Support for a large number of streams is required by industrial backbone networks.

Thank You!

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Discussion