#### IEEE 802.16 Presentation Submission Template (Rev. 9)

Document Number:

IEEE \$802.16maint-08/220

Date Submitted:

2008-05-13

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#### Venue:

IEEE 802.16-08/018, "IEEE 802.16 Working Group Letter Ballot Recirc #26c: Announcement"

Base Contribution:

IEEE C802.16maint-08/220

#### Purpose:

Review and discuss in support for the adoption of the proposal contained in C80216maint-08/220

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#### Contents

□ Summary

□ Fragmentation conditions example

Case in point

> Authentication and key exchange during network entry

HARQ considerations

Limiting scope of ARQ on primary management connection

### Summary

- MAC layer fragmentation supported for:
  - Basic, primary, secondary management connections
  - Transport connections.
- MAC layer ARQ (retransmissions of fragments) is supported for:
  - Secondary management connection
  - Transport connections
- ARQ is not supported for primary management connections
  - Transmission failure recover only by full message retransmission

#### Summary (cont.)

- In poor radio conditions, such as at the cell edge, multiple factors contribute to an increase in latency
  - More messages will be fragmented due to low modulation order
    - Message transmission is spread over multiple frames
  - Transmission errors cause retransmissions of entire MAC messages
  - MAC message retransmission is timer-based
    - Retransmissions are not triggered at the time of a transmission failure
    - Retransmissions triggered after a timer has expired and a response to the message has not been received
  - If ARQ were supported for primary management connection:
    - Retransmissions would be triggered at the time of a transmission failure
    - Only those blocks of data that fail would be retransmitted
  - This is particularly relevant for network entry, which often occurs at the cell edge and is time critical
- Avoid using ARQ for every management message
  - ARQ costs overhead for
  - Use ARQ only for fragmented SDUs

#### Proposal: Add support for ARQ for the primary management connection

Fragmentation conditions example

- Assumptions for MS uplink transmissions near cell edge:
  - DL:UL split of 23:24
  - MS limited to no better than QPSK, rate ½ coding
  - The MS is not able to transmit 2 UL subchannels concurrently
  - Allocation sizes would then be limited to 384 bits
  - Each fragment contains a GMH, FSH, 32-bit CRC
- 288 bits are available for payload
- Messages larger than 288 bits would be fragmented

### Case in point

- Authentication and key exchange during network entry
  - The PKM-REQ/RSP messages are used to encapsulate messages from multiple authentication protocols
  - For device certification, an X.509 certificate is used.
    - For PKMv1 an X.509 certificate may be sent in the Auth Request message
    - For PKMv2, an X.509 certificate may be sent in an EAP Response
  - The X.509 device certificate is around 1000 bytes
    - Usually a "certificate chain" will be transmitted
    - Includes multiple certificates so the network can trace the MS's authorization to its root certification authority
  - EAP-TLS supports fragmentation but NWG has set the MTU for EAP-TLS to be 1400 bytes
  - Messages carrying this certificate will be split into multiple fragments, particularly when transmitted near the cell edge

#### Case in point (cont.)

- Retransmission of PKM messages is dependent on the protocol being encapsulating
  - EAP is required by NWG for device authentication using X.509 certificate
  - EAP supports message retransmission but the EAP specification recommends to avoid it, if possible
    - EAP is a "lockstep" protocol
      - Only one message may be outstanding at any time
      - This means a Response to a Request message must be received before a new Request may be sent
      - It is inefficient when a message is fragmented above EAP (EAP-TLS)
    - Receiver must send Response to all retransmitted Requests
      - If receiver sends a Response but the transmitter's Request timer expires before the transmitter receives it
      - Transmitter will retransmit the Request
      - The receiver must send a Response to this retransmitted Request also
- Without ARQ, latency for the authentication and key exchange phase of network entry will be significantly increased

### HARQ considerations

- HARQ is supported for the primary management connection
- Use ARQ in conjunction with HARQ.

## Overview

- In order to examine the performance of PKM message delivery with HARQ for celledge mobiles, system-level simulations were performed
- The obtained results are:
  - Users in the worst radio conditions (worst 1%)
    - 10% of the time can't deliver a 3000 byte message over the air in less than 3 seconds.
  - Users in worst 5% radio conditions
    - 10% of the time can't deliver a 3000 byte message over the air in less than 1.5 seconds
  - Users in worst 10 and 20% radio conditions,
    - 10% of the time can't deliver a 3000 byte message over the air in less than 750 msecs.

- Figure 1 shows the resulting geometry distributions for each of the four simulations
- Figure 2 shows:
  - Users in the worst radio conditions (worst 1%)
    - 10% of the time can't deliver a 3000 byte message over the air in less than 3 seconds.
  - Users in worst 5% radio conditions
    - 10% of the time can't deliver a 3000 byte message over the air in less than 1.2 seconds
  - Users in worst 10 and 20% radio conditions,
    - 10% of the time can't deliver a 3000 byte message over the air in less than 750 msecs.
- Figure 3 shows the statistics for the number of fragments required for PKM message delivery

Figure 1: Resulting geometry distributions for each of the four simulations



| Worst 1%:  | -5.3 dB threshold   |
|------------|---------------------|
| Worst 5%:  | -3.563 dB threshold |
| Worst 10%: | -2.589 dB threshold |
| Worst 20%: | -1.24 dB threshold  |

Figure 2: Time required for a PKM packet to be delivered measured in seconds



| Worst 1%:  | -5.3 dB threshold   |
|------------|---------------------|
| Worst 5%:  | -3.563 dB threshold |
| Worst 10%: | -2.589 dB threshold |
| Worst 20%: | -1.24 dB threshold  |

Figure 3: Number of fragments for delivery of PKM message



#### Limiting use of ARQ

- In 802.16e, if ARQ is enabled for a connection:
  - It is used for every SDU on the connection
  - The fragmentation subheader is added even if the SDU is not fragmented
    - Includes an 11-bit sequence number
  - The blocks of every message would require acknowledgement
- It is best to limit the use of ARQ to avoid unnecessary overhead

Limiting the scope of ARQ (cont.)

- In order to detect lost blocks, the BSN must be sequential
- It is proposed that:
  - ARQ is used only when a MAC management SDU is fragmented
  - BSNs are only assigned to the blocks of SDUs that use a fragmentation/packing header
  - If the fragmentation/packing header is not included in a MPDU:
    - The receiver knows that ARQ is not used for the SDU
  - If the fragmentation/packing header is included but the fragmentation state=00 (no fragmentation):
    - The receiver knows that ARQ is not used for this SDU
    - However, in this case, BSNs are assigned to the blocks of the SDU
    - Only the blocks associated with ARQ (fragmented SDUs) are ack/nacked.

Limiting the scope of ARQ (cont.)

- In order to detect lost blocks, the BSN must be sequential
- It is proposed that:
  - ARQ is used only when a MAC management SDU is fragmented
  - BSNs are only assigned to the blocks of SDUs that use a fragmentation/packing header
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  - If the fragmentation/packing header is included but the fragmentation state=00 (no fragmentation):
    - The receiver knows that ARQ is not used for this SDU
    - However, in this case, BSNs are assigned to the blocks of the SDU
    - Only the blocks associated with ARQ (fragmented SDUs) are ack/nacked

- Alternatively, ARQ for all primary management connection SDUs may be enabled using a new TLV:
  - If ARQ\_ALL\_PMC\_SDUS is set to '1':
    - Behavior is the same as for transport connections
    - Blocks of all SDUs are assigned a BSN and the fragmentation or packing header is used for all PDUs

• Backup slides for simulations

 The simulation assumptions are based on the 802.16m Evaluation Methodology Document (IEEE 80216m-07\_037r2). The primary simulation assumptions are summarized in Tables 2-5.

| Layout Model       |                                   |
|--------------------|-----------------------------------|
| Network Topology:  | 19cell, 3sectors/cell, wraparound |
| BS-BS Distance:    | 1.5 km                            |
| Center Frequency:  | 2.5 GHz                           |
| Channel Bandwidth: | 10 MHz                            |
| Frequency Reuse:   | 1                                 |

| Table 2. System-Level Simulation Layout Assumptions | Table 2. | System-Level Simulation Layout Assumptions |
|---|----------|--|
|---|----------|--|

#### Table 3. System-Level Simulation BS & MS Assumptions

| Base Station Model        |  |  |
|---------------------------|--|--|
| Max TX Power Per Sector:  | 46 dBm                                 |  |
| BS Height:                | 32 m                                   |  |
| Sector Antenna Pattern:   | 3 dB beamwidth of 70°; 20 dB F/B Ratio |  |
| Sector Gain:              | 17 dBi                                 |  |
| Cable Loss:               | 2 dB                                   |  |
| Penetration Loss:         | 10 dB                                  |  |
| Number of RX Antennas:    | 2                                      |  |
| MS Model                  |  |  |
| MS Height:                | 1.5 m                                  |  |
| MS Noise Figure:          | 7 dB                                   |  |
| MS Antenna Pattern:       | Omni-directional                       |  |
| MS Antenna Gain:          | 0 dBi                                  |  |
| Number of MS TX antennas: | 1                                      |  |

| Propagation Model          |  |  |
|----------------------------|--|--|
| Pathloss Model:            | Loss (dB) = 130.62+37.6*log10(R, km)                   |  |
| Lognormal Shadow Fading:   | μ=0 dB, σ=8 dB   |  |
| Shadow Fading Correlation: | 100% inter-sector, 50% inter-BS, 50 m corr. distance   |  |
| Channel Model:             | Modified ITU Ped B, 3 km/hr (60% of users)             |  |
|                            | Modified ITU Ped A, 30 km/hr (30% of users)            |  |
|                            | Modified ITU Ped A, 120 km/hr (10% of users)           |  |
| Temporal Correlation:      | Jakes Spectrum   |  |
| Spatial Correlation:       | Correlated antennas at BS (4 $\lambda$ spacing, 3° AS) |  |

| PHY Assumptions              |   |  |
|------------------------------|---|--|
| Frame Duration:              | 5 ms  |  |
| UL Permutation:              | PUSC, reuse 1                               |  |
| # of Symbols in UL subframe: | 24 (first 3 reserved for control signaling) |  |
| UL Transmission Scheme:      | 1x2 MRC                                     |  |
| PHY Abstraction:             | RBIR  |  |
| Channel Estimation:          | Modeled                                     |  |
| HARQ Type:                   | Chase Combining (Maximum of 4 TX)           |  |
| Primary System Load:         | UL VoIP (100 users/sector)                  |  |

# Simulation Methodology

- The general simulation methodology was as follows:
  - 1. Randomly drop 100 users/sector throughout the network and establish UL traffic connections for conducting UL VoIP traffic.
  - 2. Generate 1 user/cell at 1%, 5%, 10%, 20% worst cell geometry and establish UL primary connections for conducting PKM message delivery.
    - In simulation 1, each test mobile (i.e. PKM message mobile) was re-dropped if their geometry not worse than or equal to the 1% worst cell geometry.
    - In simulation 2, each test mobile (i.e. PKM message mobile) was re-dropped if their geometry not worse than or equal to the 5% worst cell geometry.
    - In simulation 3, each test mobile (i.e. PKM message mobile) was re-dropped if their geometry not worse than or equal to the 10% worst cell geometry.
    - In simulation 4, each test mobile (i.e. PKM message mobile) was re-dropped if their geometry not worse than or equal to the 20% worst cell geometry.
  - 3. At time t=0 seconds, begin UL VoIP traffic. Allow to warm-up until time t=2 seconds.
  - 4. At t=2 seconds, a PKM packet is queued for each PKM mobile. The PKM packet size is 3000 bytes.
    - PKM message size determined by 1000 byte X.509 certificate size and includes the certificate chain

# Simulation Methodology (continued)

- 5. Once the PKM packet is queued, the scheduler begins making allocations for its delivery using the following methodology:
  - 1. Priority is given to UL VoIP traffic. The UL VoIP traffic is allowed to consume all bandwidth at peak demand times.
  - 2. PKM messages are given secondary priority. The scheduler will make allocations for PKM message delivery only if bandwidth is available to make the allocation.
  - 3. PKM transmissions are performed using UL open-loop power control.
  - 4. PKM allocations are scheduled based on the UL available bandwidth, the UL noise+interference measurement, and the estimated average UL propagation loss

# Simulation Methodology (continued)

- 5. The scheduler assumed the following values when making allocations for the PKM messages:
  - 1. An UL receive antenna gain of 3 dB due to the 2 transmit antennas and MRC
  - 2. A 3 dB mobile station boost above and beyond the UL open-loop power control table
    - This is the "Relative Power Offset for UL Burst Containing MAC Management Message".
  - 3. A HARQ gain of 6 dB.
    - This is the "Relative Power Offset For UL HARQ burst"
- 6. Because of the associated MAC overhead (48 bits for GMH, 16 bits for FSH, 16 bits for CRC), the minimum allowable allocations was for an Nep of size 96 bits.
- 7. Delay caused by errors in the downlink, such as DL Ack/Nack errors was not considered
- 8. Once all PKM messages had been delivered, the statistics related to the PKM messages were appended to an output file, and the simulation trial was repeated for a new set of VoIP/PKM mobiles beginning with step 1 described previously.