

# CAUI-4 MTTFPA monitoring

Pete Anslow, Ciena

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# Introduction

The impact on MTTFPA performance of a DFE in the CAUI-4 C2C receiver was analysed in [anslow\\_03\\_0913](#).

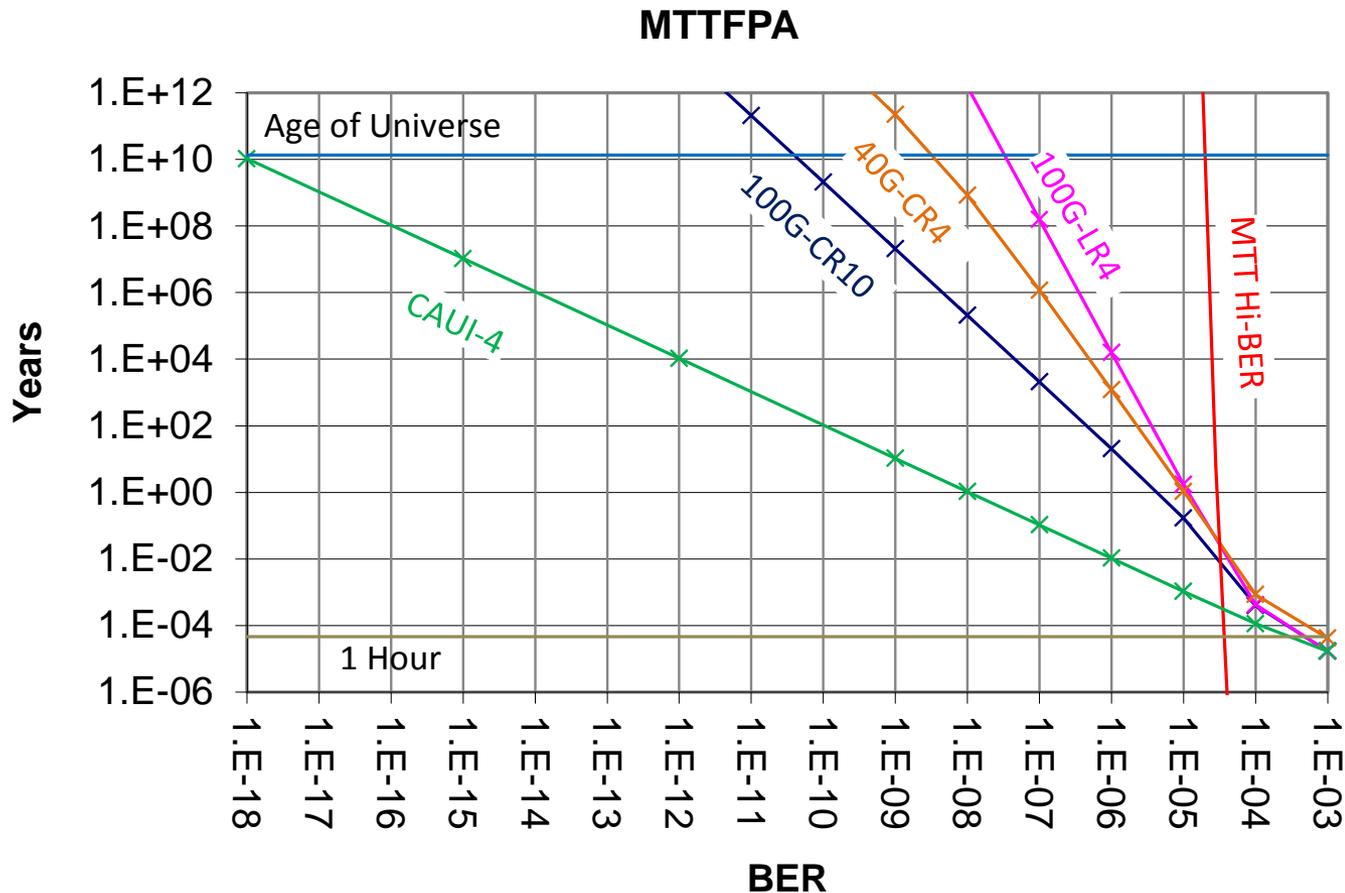
A possible BIP based error monitoring solution was proposed in [ran\\_01\\_101413\\_CAUI](#) and [Multi-lane\\_BIP](#) and this was discussed in the CAUI-4 Ad Hoc.

This presentation looks at:

- What BER would be required to give an MTTFPA of the age of the universe with no restriction on the DFE – Slides 3 and 4
- How often the proposed hi\_bip\_mismatch mechanism would trigger vs. BER for random and burst errors – Slide 5
- Proposes and analyses an improved algorithm for hi\_bip\_mismatch using multiple 2-bit events – Slides 6 and 7
- Analyses the performance of the improved algorithm using the existing BIP counters read once a second – Slides 8 and 9

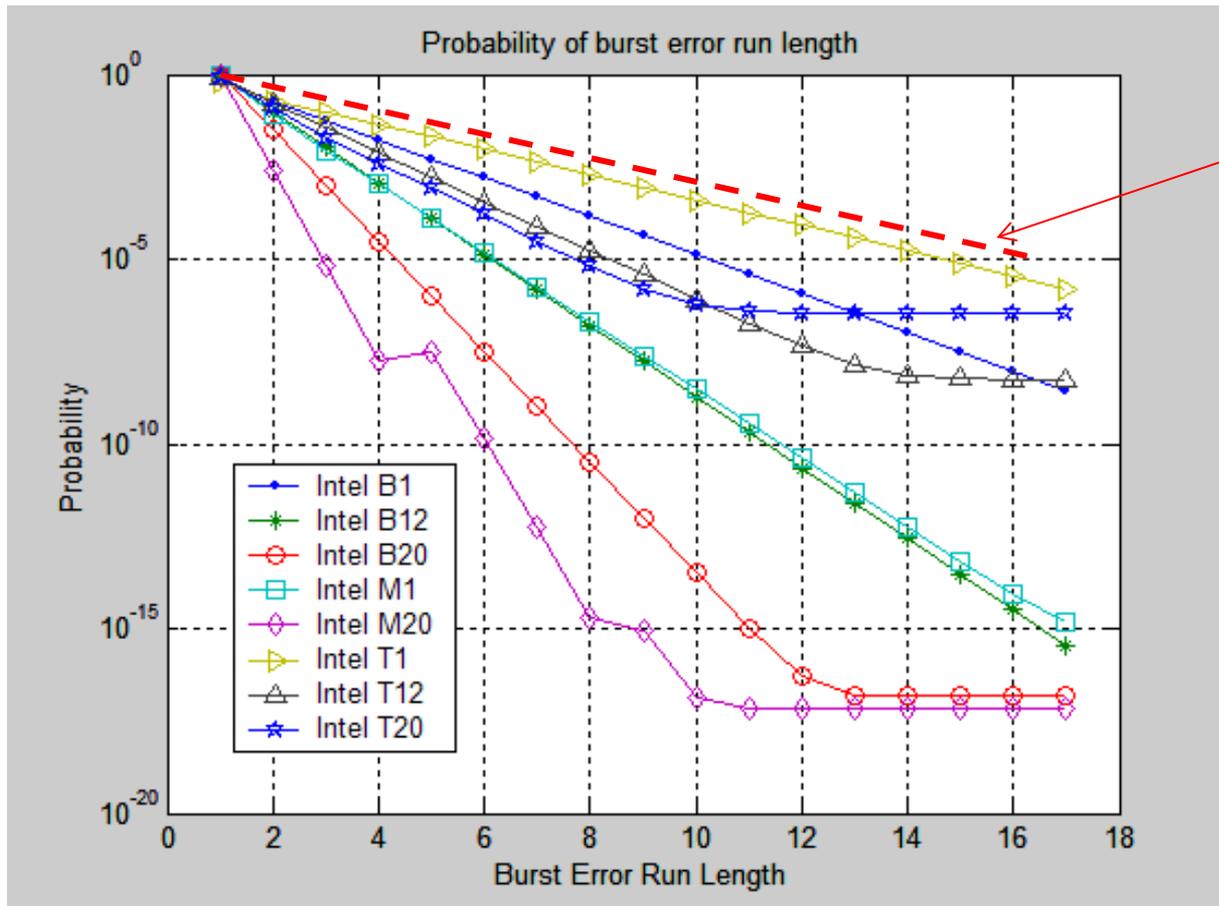
# MTTFPA for a = 0.5

- If we set the probability that a burst continues to the worst case value of  $a = 0.5$  as proposed in [cideciyan\\_01\\_0512](#) (slide 3), then we need a BER of less than  $1E-18$  for  $MTTFPA >$  age of universe. This is 3.2 errors / year.



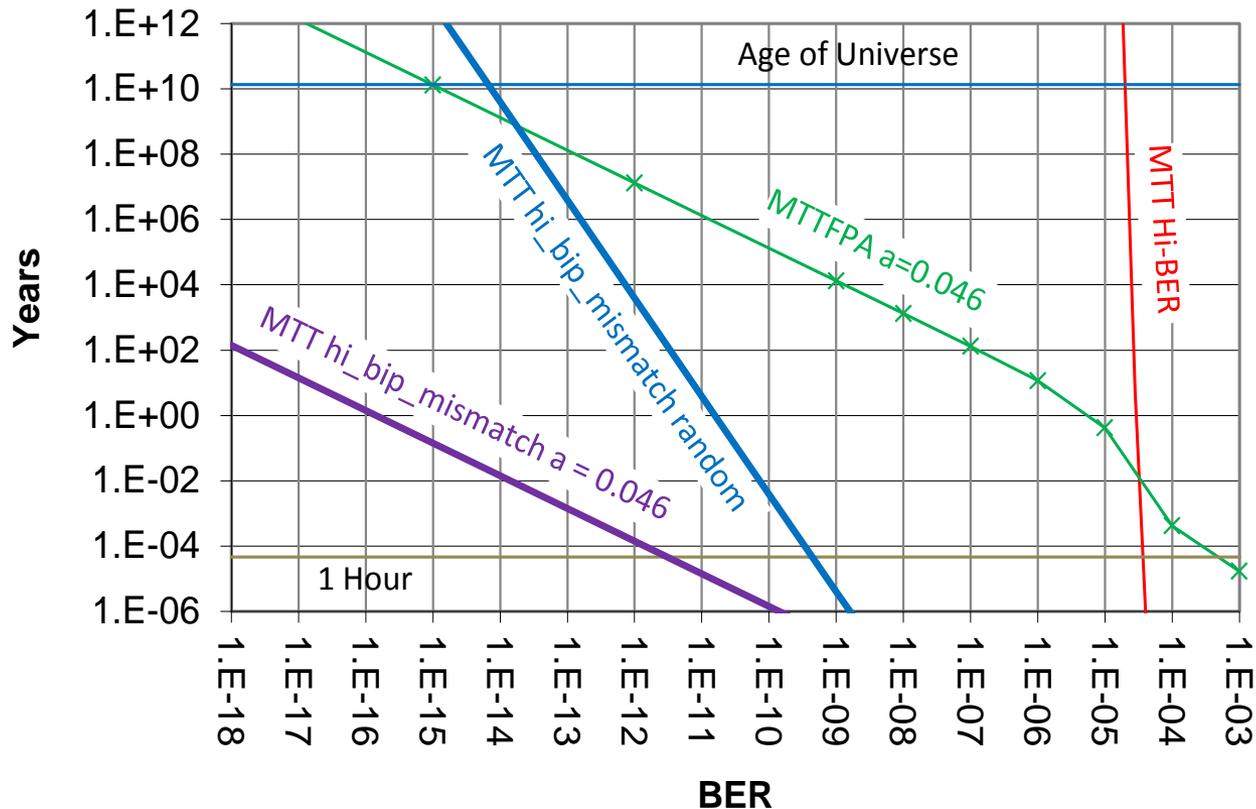
# How conservative is $a = 0.5$ ?

- Taking the worst case curves found in [liu 01 1105](#) (page 13) and superimposing the curve for  $a = 0.5$  shows that this is not very conservative if the DFE implementation is not restricted.



# Mean time to hi\_bip\_mismatch

- The mean time to high\_bip\_mismatch is plotted below for the case of random errors (blue) and for  $a = 0.046$  which gives an MTTFPA of the age of the universe (purple). For  $a = 0.046$  and a BER of  $1E-15$  the MTT high\_bip\_mismatch is 50 days, for  $1E-16$  it is 1 year.



# Discussion

The previous curves are based on asserting `hi_bip_mismatch` for a single event with 3 mis-matched BIPs as per the definition in [Multi-lane\\_BIP](#).

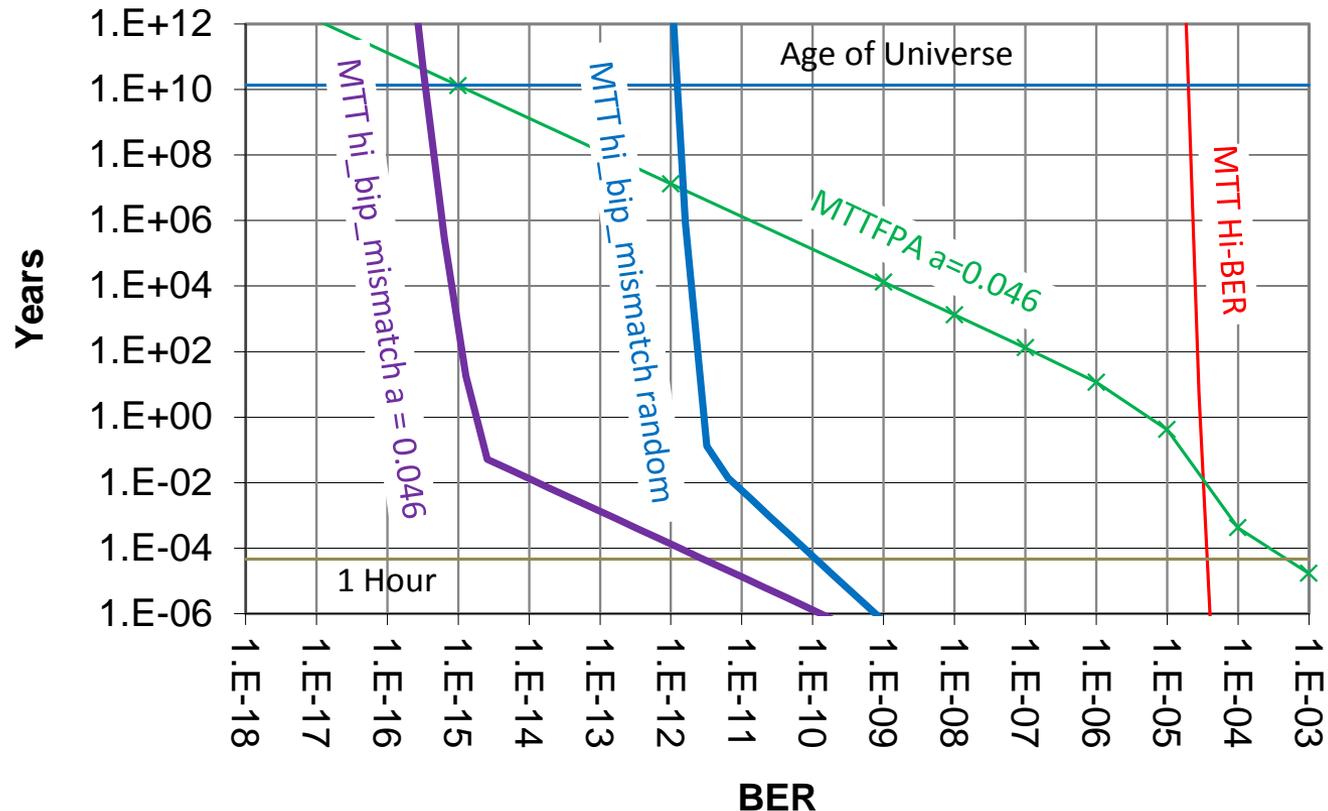
Better performance would be expected by using multiple events. As the arrival rate of 3-bit bursts for the limiting case is about one per 50 days, it seems that multiple 2-bit bursts would need to be used.

The next slide shows the performance based on asserting `hi_bip_mismatch` for greater than 20 events with 2 or more mis-matched BIPs within a 20-day window.

It is assumed that `hi_bip_mismatch` is asserted as soon as 20 events have occurred (rather than waiting for a complete 20 day window to expire).

# 20 2-bit events in 20 days

- The same curves as for slide 5 but with `hi_bip_mismatch` asserted for greater than 20 events with 2 or more mis-matched BIPs within a 20-day window. This has improved the MTT `hi_bip_mismatch` for a BER of  $1E-16$  and  $a = 0.046$  as well as for BER of  $1E-12$  random to greater than the age of the universe.



# Using existing counters

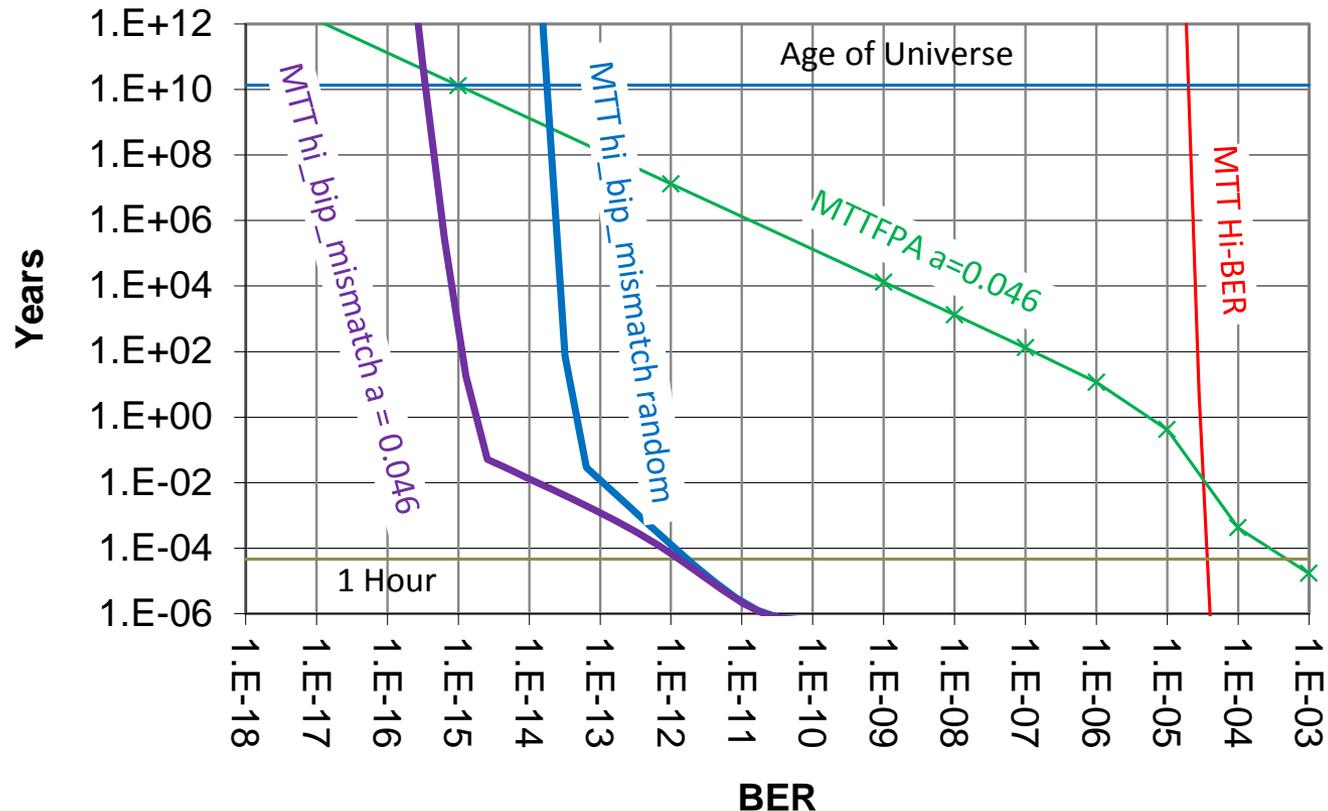
The algorithm defined in in Multi-lane\_BIP would require a hardware change in the PCS to implement. To avoid this [ran\\_01\\_101413\\_CAUI](#) raised the possibility of reading the existing BIP mismatch counters once a second and counting how often two or more lane counters have registered a mismatch.

The next slide shows the performance of this based on asserting hi\_bip\_mismatch for greater than 20 events with 2 or more lanes with errors in 1 second within a 20-day window.

It is assumed that hi\_bip\_mismatch is asserted as soon as 20 events have occurred (rather than waiting for a complete 20 day window to expire).

# 20 2-bit events in 20 days

- The same curves as for slide 7 but using the existing counters read once a second with hi\_bip\_mismatch asserted for greater than 20 events with 2 or more lanes with errors within a 20-day window. This has reduced the MTT high\_bip\_mismatch for BER of 1E-12 random to an hour.



# Performance summary

Accounting for burst errors by just improving the BER requires changing the BER requirement to  $< 1E-18$  ( $< 3.2$  errors / year). This does not seem practical.

The performance of the three hi\_bip\_mismatch algorithms studied is:

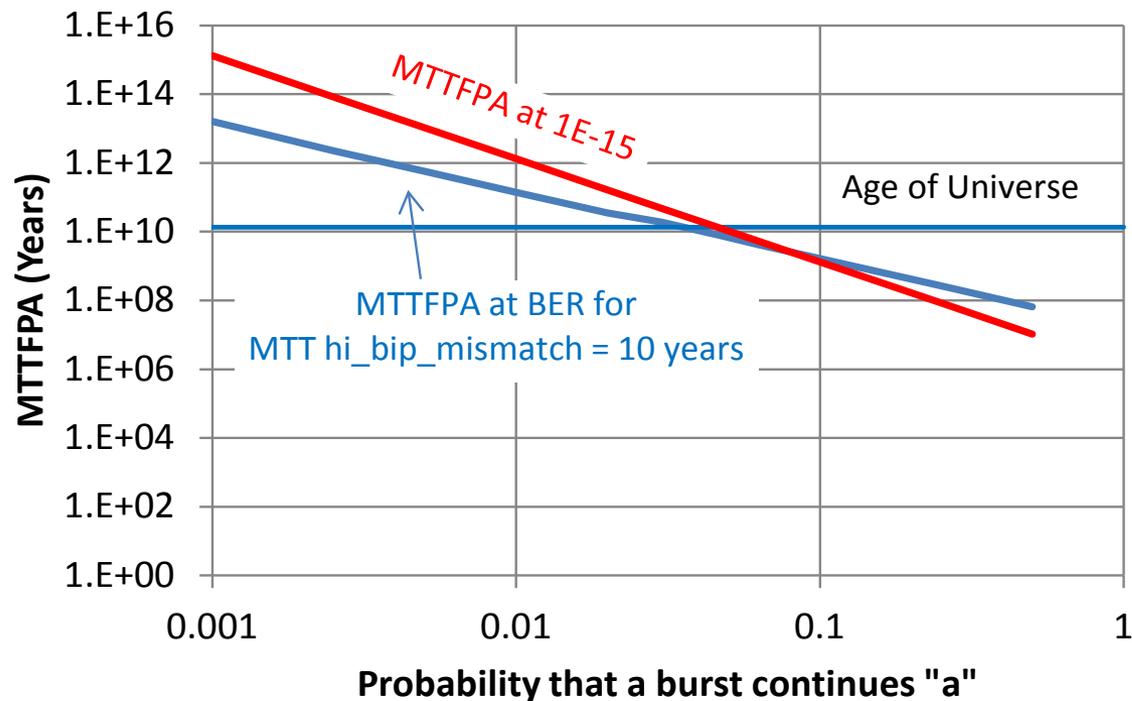
	3E-16 a=0.046	1E-15 a=0.046	1E-12 Random	1E-11 Random
3 mismatched BIPs	6 months	50 days	10,000 yr	4 years
20 x (2 mismatched BIPs) in 20 days	> AOU	300 years	>AOU	2 days
20 x (2 lanes with Err in 1 sec) in 20 days	> AOU	300 years	1 hour	1 min

The only one of these that might be acceptable is 20 x (2 mismatched BIPs) in 20 days.

Whether this algorithm is successful in protecting against false packet acceptance is also dependent on whether the relationship between the probability of 2-bit bursts and 4-bit bursts is as modelled or not.

# 20 2-bit events in 20 days

- For the 20 x (2 mismatched BIPs) in 20 days the blue curve below plots the MTTFPA at whatever BER gives a MTT hi\_bip\_mismatch of 10 years. This isn't much of an improvement over the red curve which is the MTTFPA at a BER of 1E-15.



# Conclusion

None of the algorithms analysed in this presentation meet the goals of:

- Triggering reasonably frequently whenever the MTTFPA is below the age of the universe and the CAUI-4 BER is  $< 1E-15$
- Not triggering for a 100GBASE-LR4 link operating at a BER of  $1E-12$

To meet the requirement of triggering reasonably frequently whenever the MTTFPA is below the age of the universe with  $a = 0.5$  the threshold would have to be set at something in the region of  $4 \times$  (2 mismatched BIPs) in 365 days which isn't very practical and would trigger about once a month for LR4 at  $1E-12$ .

Still to do:

Investigate the possibility of using an algorithm based on the ratio between 1-bit events and 2-bit events.

Thanks!